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CMOS 16-BIT SINGLE CHIP MICROCONTROLLER

S1C17 Family
S1C17 Core Manual

SEIKO EPSON CORPORATION

Rev. 1.2

– Contents –

1 Summary

The S1C17 Core is a Seiko Epson original 16-bit RISC-type processor.
It features low power consumption, high-speed operation with a maximum 60 MHz to 90 MHz clock, large address space up to 16M bytes addressable, main instructions executable in one clock cycle, and a small sized design. The S1C17 Core is suitable for embedded applications that do not need a lot of data processing power like the S1C33 Cores the high-end processors, such as controllers and sequencers for which an eight-bit CPU is commonly used.
The S1C17 Core incorporates a coprocessor interface allowing implementation of additional computing features.

Furthermore, Seiko Epson provides a software development environment similar to the S1C33 Family that includes an IDE work bench, a C compiler, a serial ICE and a debugger, for supporting the developer to develop application software.

1.1 Features

Processor type

- Seiko Epson original 16-bit RISC processor
- 0.35–0.15 µm low power CMOS process technology

Operating-clock frequency

- 90 MHz maximum (depending on the processor model and process technology)

Instruction set

- Code length: 16-bit fixed length
- Number of instructions: 111 basic instructions (184 including variations)
- Execution cycle: Main instructions executed in one cycle
- Extended immediate instructions: Immediate extended up to 24 bits
- Compact and fast instruction set optimized for development in C language

Register set

- Eight 24-bit general-purpose registers
- Two 24-bit special registers
- One 8-bit special register

Memory space and bus

- Up to 16M bytes of memory space (24-bit address)
- Harvard architecture using separated instruction bus (16 bits) and data bus (32 bits)

Interrupts

- Reset, NMI, and 32 external interrupts supported
- Address misaligned interrupt
- Debug interrupt
- Direct branching from vector table to interrupt handler routine
- Programmable software interrupts with a vector number specified (all vector numbers specifiable)

Power saving

- HALT (`halt` instruction)
- SLEEP (`slp` instruction)

Coprocessor interface

- ALU instructions can be enhanced

2 Registers

The S1C17 Core contains eight general-purpose registers and three special registers.

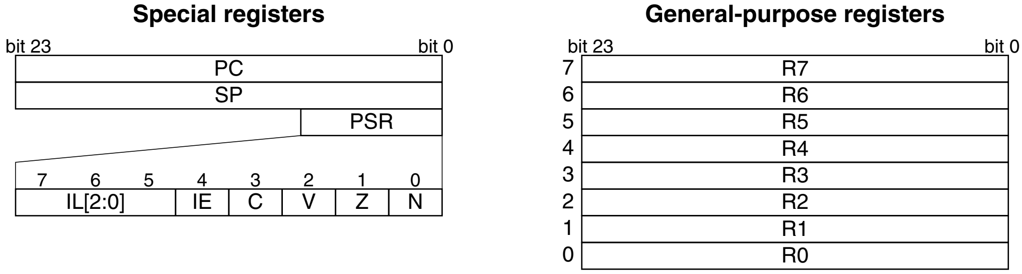


Figure 2.1 Registers

2.1 General-Purpose Registers (R0–R7)

Symbol	Register name	Size	R/W	Initial value
R0–R7	General-Purpose Register	24 bits	R/W	0x000000

The eight registers R0–R7 (r0–r7) are 24-bit general-purpose registers that can be used for data manipulation, data transfer, memory addressing, or other general purposes. The contents of all of these registers are handled as 24-bit data or addresses. 8- or 16-bit data can be sign- or zero-extended to a 24-bit quantity when it is loaded into one of these registers using a load instruction or a conversion instruction. When these registers are used for address references, 24-bit memory space can be accessed directly.
At initial reset, the contents of the general-purpose registers are set to 0.

2.2 Program Counter (PC)

Symbol	Register name	Size	R/W	Initial value
PC	Program Counter	24 bits	R	(Reset vector)

The Program Counter (hereinafter referred to as the "PC") is a 24-bit counter for holding the address of an instruction to be executed. More specifically, the PC value indicates the address of the next instruction to be executed.
As the instructions in the S1C17 Core are fixed at 16 bits in length, the LSB (bit 0) of the PC is always 0.
Although the S1C17 Core allows the PC to be referenced in a program, the user cannot alter it. Note, however, that the value actually loaded into the register when a `ld.a %rd,%pc` instruction (can be executed as a delayed slot instruction) is executed is the "PC value for the `ld` instruction + 2."
At an initial reset, the reset vector (address) written at the top of vector table indicated by TTBR is loaded into the PC, and the processor starts executing a program from the address indicated by the PC.

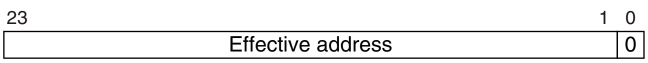


Figure 2.2.1 Program Counter (PC)

2.3 Processor Status Register (PSR)

Symbol	Register name	Size	R/W	Initial value
PSR	Processor Status Register	8 bits	R/W	0x00

The Processor Status Register (hereinafter referred to as the "PSR") is an 8-bit register for storing the internal status of the processor.
The PSR stores the internal status of the processor when the status has been changed by instruction execution. It is referenced in arithmetic operations or branch instructions, and therefore constitutes an important internal status in program composition. The PSR does not allow the program to directly alter its contents except for the IE bit.
As the PSR affects program execution, whenever an interrupt occurs, the PSR is saved to the stack, except for debug interrupts, to maintain the PSR value. The IE flag (bit 4) in it is cleared to 0. The `reti` instruction is used to return from interrupt handling, and the PSR value is restored from the stack at the same time.

	7	6	5	4	3	2	1	0
PSR	IL[2:0]			IE	C	V	Z	N
Initial value	0	0	0	0	0	0	0	0

Figure 2.3.1 Processor Status Register (PSR)

IL[2:0] (bits 7–5): Interrupt Level

These bits indicate the priority levels of the processor interrupts. Maskable interrupt requests are accepted only when their priority levels are higher than that set in the IL bit field. When an interrupt request is accepted, the IL bit field is set to the priority level of that interrupt, and all interrupt requests generated thereafter with the same or lower priority levels are masked, unless the IL bit field is set to a different level or the interrupt handler routine is terminated by the `reti` instruction.

IE (bit 4): Interrupt Enable

This bit controls maskable external interrupts by accepting or disabling them. When IE bit = 1, the processor enables maskable external interrupts. When IE bit = 0, the processor disables maskable external interrupts.
When an interrupt is accepted, the PSR is saved to the stack and this bit is cleared to 0. However, the PSR is not saved to the stack for debug interrupts, nor is this bit cleared to 0.

C (bit 3): Carry

This bit indicates a carry or borrow. More specifically, this bit is set to 1 when, in an add or subtract instruction in which the result of operation is handled as an unsigned 16-bit or 24-bit integer, the execution of the instruction resulted in exceeding the range of values representable by an unsigned 16-bit or 24-bit integer, or is reset to 0 when the result is within the range of said values.
Furthermore, the C flag will be set or reset by executing an shift instruction.
The C flag is set under the following conditions:

(1) When an addition executed by a 16-bit integer addition instruction (except a case of conditional execution) results in a value greater than the maximum value 0xffff representable by an unsigned 16-bit integer

(2) When a subtraction executed by a 16-bit integer subtraction instruction (except a case of conditional execution) results in a value smaller than the minimum value 0x0000 representable by an unsigned 16-bit integer

(3) When a comparison (subtraction) executed by a 16-bit integer comparison instruction (except a case of conditional execution) results in a value smaller than the minimum value 0x0000 representable by an unsigned 16-bit integer

(4) When a comparison (subtraction) executed by a 24-bit integer comparison instruction (except a case of conditional execution) results in a value smaller than the minimum value 0x000000 representable by an unsigned 24-bit integer

(5) When a shift operation of the register in which bit 0 is 1 is executed using a right logical shift instruction

(6) When a shift operation of the register in which bit 15 is 1 is executed using a left logical shift instruction

(7) When a shift operation of the register in which bit 0 is 1 is executed using a right arithmetic shift instruction

V (bit 2): OVerflow

This bit indicates that an overflow or underflow occurred in an arithmetic operation. More specifically, this bit is set to 1 when, in an add or subtract instruction in which the result of operation is handled as a signed 16-bit integer, the execution of the instruction resulted in an overflow or underflow, or is reset to 0 when the result of the add or subtract operation is within the range of values representable by a signed 16-bit integer. This flag is also reset to 0 by executing a logical operation instruction.

Note that 16-bit arithmetic operation instructions can set the V flag, but 24-bit arithmetic operation instructions cannot.

The V flag is set under the following conditions:

(1) When negative integers are added together, the operation produced a 0 (positive) in the sign bit (most significant bit of the result)

(2) When positive integers are added together, the operation resulted in a 1 (negative) in the sign bit (most significant bit of the result)

(3) When a negative integer is subtracted from a positive integer, the operation resulted in producing a 1 (negative) in the sign bit (most significant bit of the result)

(4) When a positive integer is subtracted from a negative integer, the operation resulted in producing a 0 (positive) in the sign bit (most significant bit of the result)

Z (bit 1): Zero

This bit indicates that an operation resulted in 0. More specifically, this bit is set to 1 when the execution of a logical operation, arithmetic operation, or shift instruction resulted in 0, or is otherwise reset to 0.

Note that 16-bit arithmetic operation and comparison instructions can set the Z flag, but 24-bit addition and subtraction instructions cannot.

N (bit 0): Negative

This bit indicates a sign. More specifically, the most significant bit (bit 15) of the result of a logical operation, arithmetic operation, or shift instruction is copied to this N flag.

Note that 16-bit arithmetic operation instructions can set the N flag, but 24-bit arithmetic operation instructions cannot.

2.4 Stack Pointer (SP)

Symbol	Register name	Size	R/W	Initial value
SP	Stack Pointer	24 bits	R/W	0x000000

The Stack Pointer (hereinafter referred to as the "SP") is a 24-bit register for holding the start address of the stack. The stack is an area locatable at any place in the system RAM, the start address of which is set in the SP during the initialization process. The 2 low-order bits of the SP are fixed to 0 and cannot be accessed for writing. Therefore, the addresses specifiable by the SP are those that lie on 32-bit boundaries.

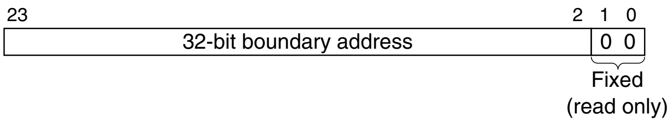


Figure 2.4.1 Stack Pointer (SP)

2.4.1 About the Stack Area

The size of an area usable as the stack is limited according to the RAM size available for the system and the size of the area occupied by ordinary RAM data. Care must be taken to prevent the stack and data area from overlapping. Furthermore, as the SP becomes 0x000000 when it is initialized upon reset, "last stack address + 4, with 2 low-order bits = 0" must be written to the SP in the beginning part of the initialization routine. A load instruction may be used to write this address. If an interrupt occurs before the stack is set up, it is possible that the PC or PSR will be saved to an indeterminate location, and normal operation of a program cannot be guaranteed. To prevent such a problem, NMIs (nonmaskable interrupts) that cannot be controlled in software are masked out in hardware until the SP is initialized.

2.4.2 SP Operation at Subroutine Call/Return

A subroutine call instruction, `call` or `calla`, uses four bytes of the stack. The `call`/`calla` instruction saves the contents of the PC (return address) onto the stack before branching to a subroutine. The saved address is restored into the PC by the `ret` instruction, and the program is returned to the address next to that of the `call`/`calla` instruction.

SP operation by the `call`/`calla` instruction

(1) SP = SP - 4

(2) PC + 2 → [SP]

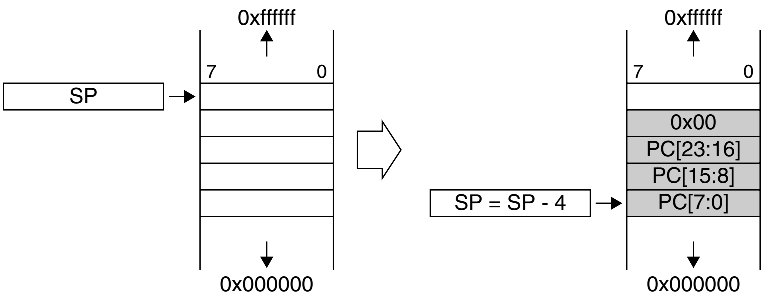


Figure 2.4.2.1 SP and Stack (1)

SP operation by the `ret` instruction

(1) [SP] → PC
(2) SP = SP + 4

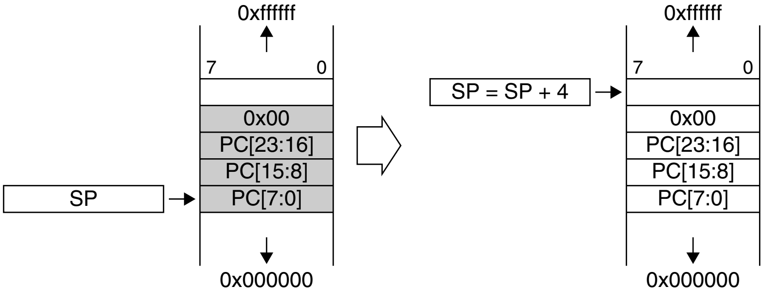


Figure 2.4.2.2 SP and Stack (2)

2.4.3 SP Operation when an Interrupt Occurs

If an interrupt or a software interrupt resulting from the `int`/`intl` instruction occurs, the processor enters an interrupt handling process.
The processor saves the contents of the PC and PSR into the stack indicated by the SP before branching to the relevant interrupt handler routine. This is to save the contents of the two registers before they are altered by interrupt handling. The PC and PSR data is saved into the stack as shown in the diagram below.
For returning from the handler routine, the `reti` instruction is used to restore the contents of the PC and PSR from the stack. In the `reti` instruction, the PC and PSR are read out of the stack, and the SP address is altered as shown in the diagram below.

SP operation when an interrupt occurred

(1) SP = SP - 4
(2) PC + 2 → [SP]
(3) PSR → [SP + 3]

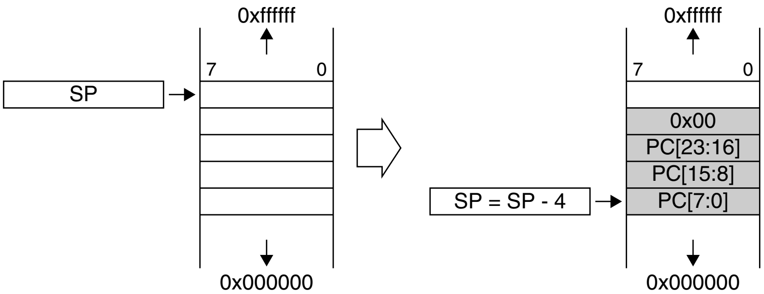


Figure 2.4.3.1 SP and Stack (3)

SP operation when the `reti` instruction is executed

(1) [SP] → PC
(2) [SP+ 3] → PSR
(3) SP = SP + 4

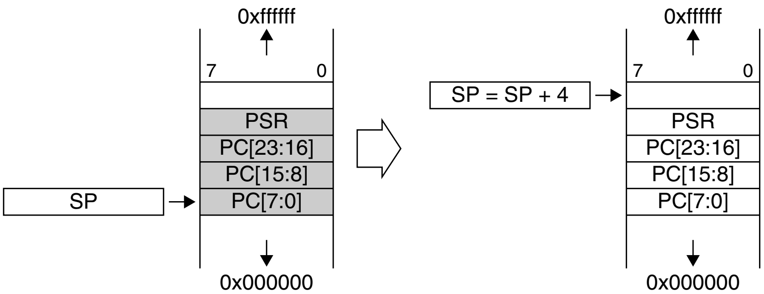


Figure 2.4.3.2 SP and Stack (4)

2.4.4 Saving/Restoring Register Data Using a Load Instruction

The S1C17 Core provides load instructions to save and restore register data to/from the stack instead of push/pop instructions.

Saving register data into the stack

Example: `ld.a -[%sp],%r0`

(1) SP = SP - 4

(2) R0 → [SP]

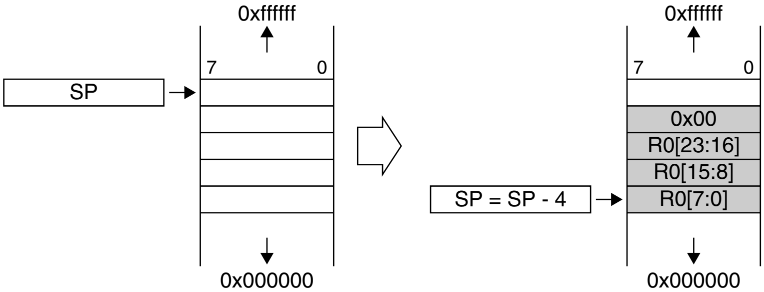


Figure 2.4.4.1 SP and Stack (5)

Restoring register data from the stack

Example: `ld.a %r0,[%sp]+`

(1) [SP] → R0

(2) SP = SP + 4

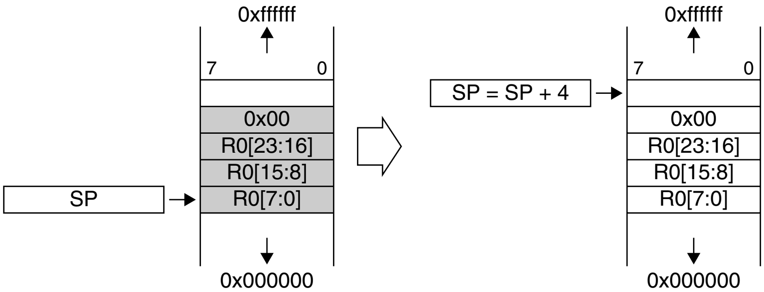


Figure 2.4.4.2 SP and Stack (6)

In addition to the instructions shown above, some other load instructions have been provided for operating the stack. Refer to Chapter 7, “Details of Instructions,” for more information on those instructions.

2.5 Register Notation and Register Numbers

The following describes the register notation and register numbers in the S1C17 Core instruction set.

2.5.1 General-Purpose Registers

In the instruction code, a general-purpose register is specified using a 3-bit field, with the register number entered in that field. In the mnemonic, a register is specified by prefixing the register name with "%."

%rs *rs* is a metasymbol indicating the general-purpose register that holds the source data to be operated on or transferred. The register is actually written as `%r0`, `%r1`, ... or `%r7`.

%rd *rd* is a metasymbol indicating the general-purpose register that is the destination in which the result of operation is to be stored or data is to be loaded. The register is actually written as `%r0`, `%r1`, ... or `%r7`.

%rb *rb* is a metasymbol indicating the general-purpose register that holds the base address of memory to be accessed. In this case, the general-purpose registers serve as an index register. The register is actually written as `[%r0]`, `[%r1]`, ... or `[%r7]`, with each register name enclosed in brackets "`[]`" to denote register indirect addressing.
In register indirect addressing, the post-increment/decrement and pre-decrement functions provided for continuous memory addresses can be used.

Post-increment function
Example: `ld %rd,[%rb]+ ; (1)ld %rd,[%rb] (2)%rb = %rb + 2`
The base address is incremented by an amount equal to the accessed size after the memory has been accessed.

Post-decrement function
Example: `ld.a %rd,[%rb]- ; (1)ld.a %rd,[%rb] (2)%rb = %rb - 4`
The base address is decremented by an amount equal to the accessed size after the memory has been accessed.

Pre-decrement function
Example: `ld.b -[%rb],%rs ; (1)%rb = %rb - 1 (2)ld.b [%rb],%rs`
The base address is decremented by an amount equal to the access size before accessing the memory.

Also any desired value can be specified as the address increment/decrement value using the `ext` instruction.

rb is also used as a symbol indicating the register that contains the jump address for the call or jump instructions. In this case, the brackets "`[]`" are unnecessary, and the register is written as `%r0`, `%r1`, ... or `%r7`.

The bit field that specifies a register in the instruction code contains the code corresponding to a given register number. The relationship between the general-purpose registers and the register numbers is listed in the table below.

Table 2.5.1.1 General-Purpose Registers

General-purpose register	Register number	Register notation
R0	0	%r0
R1	1	%r1
R2	2	%r2
R3	3	%r3
R4	4	%r4
R5	5	%r5
R6	6	%r6
R7	7	%r7

2.5.2 Special Registers

The special registers that can be directly specified in the S1C17 Core instructions are the SP (Stack Pointer) and PC (Program Counter) only. The register is actually written as `%sp`, `[%sp]`, `-[%sp]`, `[%sp]+`, `[%sp]-`, `[%sp+imm7]`, or `%pc`.

3 Data Formats

3.1 Data Formats Handled in Operations Between Registers

The S1C17 Core can handle 8-, 16-, and 24-bit data in register operations. In this manual, data sizes are expressed as follows:

8-bit data **Byte**, **B**, or **b**
16-bit data **Word**, **W**, or **w**
24-bit data **Address data**, **A**, **a**

Data sizes can be selected only in data transfer (load instruction) between one general-purpose register and another. In an 8-bit data transfer with a general-purpose register as the destination, the data is sign- or zero-extended to 16 bits before being loaded into the register. Whether the data will be sign- or zero-extended is determined by the load instruction used.
In a 16-bit or 8-bit data transfer using a general-purpose register as the source, the data to be transferred is stored in the low-order 16 bits or the low-order 8 bits of the source register.
The data transfer sizes and types are described below.

3.1.1 Unsigned 8-Bit Transfer (Register → Register)

Example: `ld.ub` *%rd*, *%rs*

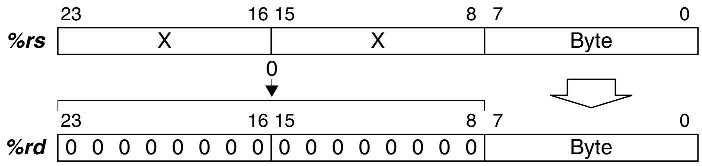


Figure 3.1.1.1 Unsigned 8-Bit Transfer (Register → Register)

Bits 23–8 in the destination register are set to 0x0000.

3.1.2 Signed 8-Bit Transfer (Register → Register)

Example: `ld.b` *%rd*, *%rs*

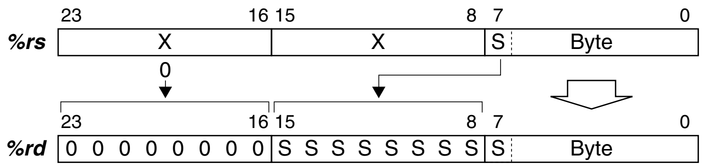


Figure 3.1.2.1 Signed 8-Bit Transfer (Register → Register)

Bits 15–8 in the destination register are sign-extended and bits 23–16 are set to 0x00.

3.1.3 16-Bit Transfer (Register → Register)

Example: `ld` *%rd*, *%rs*

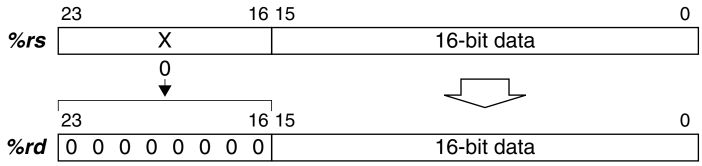


Figure 3.1.3.1 16-Bit Transfer (Register → Register)

Bits 23–16 in the destination register are set to 0x00.

3.1.4 24-Bit Transfer (Register → Register)

Example: `ld.a` *%rd*, *%rs*

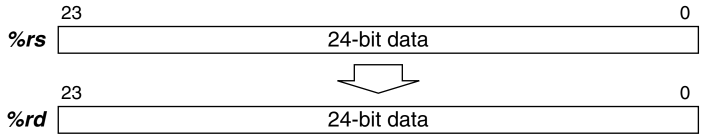


Figure 3.1.4.1 24-Bit Transfer (Register → Register)

3.2 Data Formats Handled in Operations Between Memory and a Register

The S1C17 Core can handle 8-, 16-, and 32-bit data in memory operations. In this manual, data sizes are expressed as follows:

8-bit data **Byte**, **B**, or **b**
16-bit data **Word**, **W**, or **w**
32-bit data **Address data**, **A**, **a**

Data sizes can be selected only in data transfer (load instruction) between memory and a general-purpose register.
In an 8-bit data transfer with a general-purpose register as the destination, the data is sign- or zero-extended to 16 bits before being loaded into the register. Whether the data will be sign- or zero-extended is determined by the load instruction used.
In a 16-bit or 8-bit data transfer using a general-purpose register as the source, the data to be transferred is stored in the low-order 16 bits or the low-order 8 bits of the source register.
Memory is accessed in little endian format one byte, 16 bits, or 32 bits at a time.
If memory is to be accessed in 16-bit or 32-bit units, the specified base address must be on a 16-bit boundary (least significant address bit = 0) or 32-bit boundary (2 low-order address bits = 00), respectively. Unless this condition is satisfied, an address-misaligned interrupt is generated.

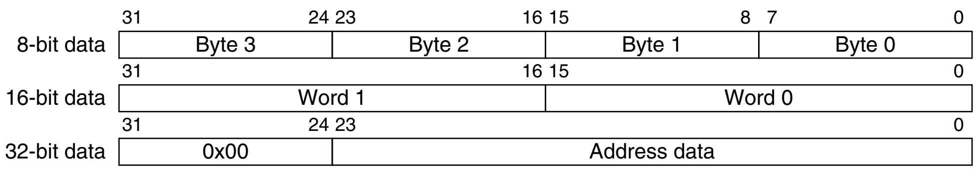


Figure 3.2.1 Data Format (Little Endian)

∗ Handling the eight high-order bits during 32-bit accesses

During writing, the eight high-order bits are written as 0. During reading from a memory, the eight high-order bits are ignored. However, the eight high-order bits are effective as the PSR value only in the stack operation when an interrupt occurs.

The data transfer sizes and types are described below.

3.2.1 Unsigned 8-Bit Transfer (Memory → Register)

Example: `ld.ub   %rd,[%rb]`

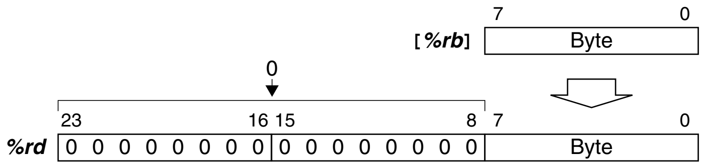


Figure 3.2.1.1 Unsigned 8-Bit Transfer (Memory → Register)

Bits 23–8 in the destination register are set to 0x0000.

3.2.2 Signed 8-Bit Transfer (Memory → Register)

Example: `ld.b   %rd,[%rb]`

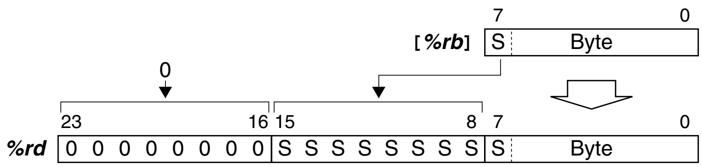


Figure 3.2.2.1 Signed 8-Bit Transfer (Memory → Register)

Bits 15–8 in the destination register are sign-extended and bits 23–16 are set to 0x00.

3.2.3 8-Bit Transfer (Register → Memory)

Example: `ld.b   [%rb],%rs`

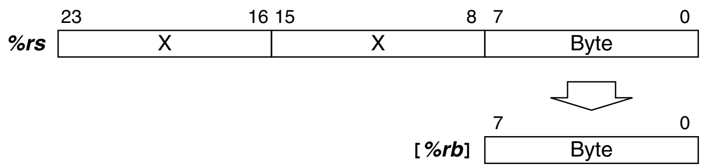


Figure 3.2.3.1 8-Bit Transfer (Register → Memory)

3.2.4 16-Bit Transfer (Memory → Register)

Example: `ld   %rd,[%rb]`

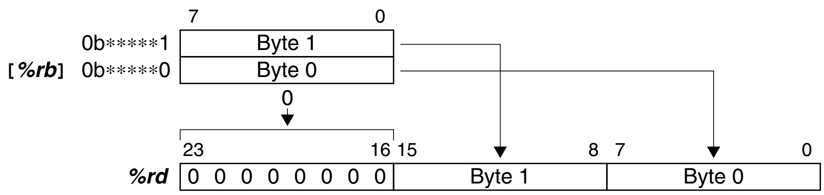


Figure 3.2.4.1 16-Bit Transfer (Memory → Register)

Bits 23–16 in the destination register are set to 0x00.

3.2.5 16-Bit Transfer (Register → Memory)

Example: `ld [%rb],%rs`

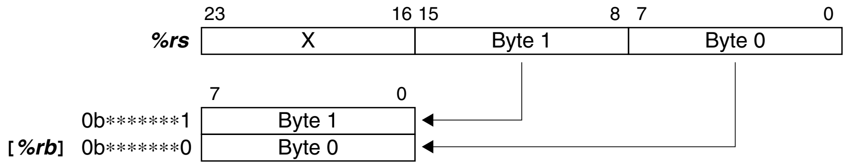


Figure 3.2.5.1 16-Bit Transfer (Register → Memory)

3.2.6 32-Bit Transfer (Memory → Register)

Example: `ld.a %rd,[%rb]`

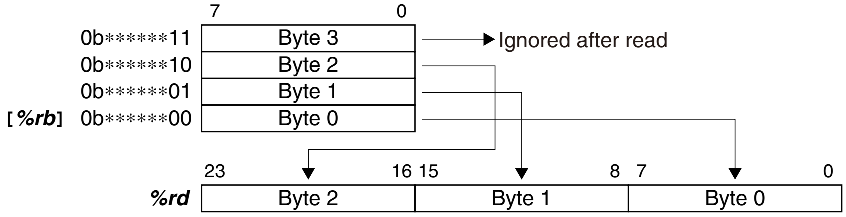


Figure 3.2.6.1 32-Bit Transfer (Memory → Register)

3.2.7 32-Bit Transfer (Register → Memory)

Example: `ld.a [%rb],%rs`

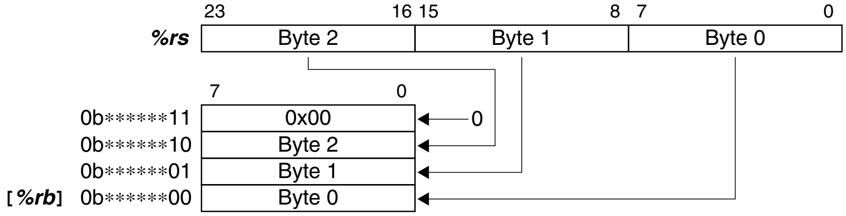


Figure 3.2.7.1 32-Bit Transfer (Register → Memory)

4 Address Map

4.1 Address Space

The S1C17 Core supports a 24-bit address allowing linear use of address space up to 16M bytes. Addresses 0xfffc00 to 0xffffff are reserved as an I/O area for the core. In addition to this area, a 64-byte area located in the user RAM is required for debugging.
Figure 4.1.1 shows the address space of the S1C17 Core.

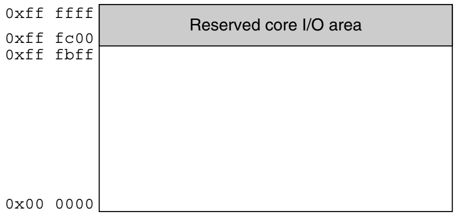


Figure 4.1.1 Address Space of the S1C17 Core

The boot address and debug RAM address depend on the specifications of each the S1C17 Series models. Refer to the Technical Manual of each model.

4.2 Processor Information in the Core I/O Area

The reserved core I/O area contains the processor information described below.

4.2.1 Vector Table Base Register (TTBR, 0xffff80)

Register name	Address	Bit	Name	Function	Setting	Init.	R/W	Remarks
Vector table base register	**FFFF80 (L)**	D31–24	–	Unused (fixed at 0)	0x0	0x0	R	
		D23 | D0	**TTBR23 | TTBR0**	Vector table base address TTBR[7:0] is fixed at 0x0.	0x0–0xFFFB00 (256 byte units)	*	R	Initial value is set by the TTBR pins of the C17 macro.

This is a read-only register that contains the vector table base address.
The vector table contains the vectors to the interrupt handler routines (handler routine start address) that will be read by the S1C17 Core to execute the handler when an interrupt occurs. The boot address from which the program starts running after a reset must be written to the top of the vector table.

Refer to the Technical Manual of each model for the address stored in this register.

4.2.2 Processor ID Register (IDIR, 0xffff84)

Register name	Address	Bit	Name	Function	Setting	Init.	R/W	Remarks
Processor ID register	**FFFF84 (B)**	D7 | D0	**IDIR7 | IDIR0**	Processor ID 0x10: S1C17 Core	0x10	0x10	R	

This is a read-only register that contains the ID code to represent a processor model. The S1C17 Core's ID code is 0x10.

4.2.3 Debug RAM Base Register (DBRAM, 0xffff90)

Register name	Address	Bit	Name	Function	Setting	Init.	R/W	Remarks
Debug RAM base register	**FFFF90 (L)**	D31–24	–	Unused (fixed at 0)	0x0	0x0	R	
		D23 | D0	**DBRAM23 | DBRAM0**	Debug RAM base address DBRAM[5:0] is fixed at 0x0.	0x0–0xFFFDC0 (64 byte units)	*	R	Initial value is set in the C17 RTL-define DBRAM_BASE.

This is a read-only register that contains the start address of a work area (64 bytes) for debugging.

Refer to the Technical Manual of each model for the address stored in this register.

* In addition to the above registers, the reserved core I/O area contains some registers for debugging. For the debug registers, refer to Section 6.5, "Debug Circuit."

5 Instruction Set

The S1C17 Core instruction codes are all fixed to 16 bits in length which, combined with pipelined processing, allows most important instructions to be executed in one cycle. For details, refer to the description of each instruction in the latter sections of this manual.

5.1 List of Instructions

Table 5.1.1 S1C17 Instructions List

Classification	Mnemonic		Function
Data transfer	ld.b	%rd,%rs	General-purpose register (byte) → general-purpose register (sign-extended)
		%rd,[%rb] %rd,[%rb]+ %rd,[%rb]- %rd,-[%rb]	Memory (byte) → general-purpose register (sign-extended) Memory address post-increment, post-decrement, and pre-decrement functions can be used.
		%rd,[%sp+imm7]	Stack (byte) → general-purpose register (sign-extended)
		%rd,[imm7]	Memory (byte) → general-purpose register (sign-extended)
		[%rb],%rs [%rb]+,%rs [%rb]-,%rs -[%rb],%rs	General-purpose register (byte) → memory Memory address post-increment, post-decrement, and pre-decrement functions can be used.
		[%sp+imm7],%rs	General-purpose register (byte) → stack
		[imm7],%rs	General-purpose register (byte) → memory
	ld.ub	%rd,%rs	General-purpose register (byte) → general-purpose register (zero-extended)
		%rd,[%rb] %rd,[%rb]+ %rd,[%rb]- %rd,-[%rb]	Memory (byte) → general-purpose register (zero-extended) Memory address post-increment, post-decrement, and pre-decrement functions can be used.
		%rd,[%sp+imm7]	Stack (byte) → general-purpose register (zero-extended)
		%rd,[imm7]	Memory (byte) → general-purpose register (zero-extended)
	ld	%rd,%rs	General-purpose register (16 bits) → general-purpose register
		%rd,sign7	Immediate → general-purpose register (sign-extended)
		%rd,[%rb] %rd,[%rb]+ %rd,[%rb]- %rd,-[%rb]	Memory (16 bits) → general-purpose register Memory address post-increment, post-decrement, and pre-decrement functions can be used.
		%rd,[%sp+imm7]	Stack (16 bits) → general-purpose register
		%rd,[imm7]	Memory (16 bits) → general-purpose register
		[%rb],%rs [%rb]+,%rs [%rb]-,%rs -[%rb],%rs	General-purpose register (16 bits) → memory Memory address post-increment, post-decrement, and pre-decrement functions can be used.
		[%sp+imm7],%rs	General-purpose register (16 bits) → stack
		[imm7],%rs	General-purpose register (16 bits) → memory
	ld.a	%rd,%rs	General-purpose register (24 bits) → general-purpose register
		%rd,imm7	Immediate → general-purpose register (zero-extended)
		%rd,[%rb] %rd,[%rb]+ %rd,[%rb]- %rd,-[%rb]	Memory (32 bits) → general-purpose register * Memory address post-increment, post-decrement, and pre-decrement functions can be used.
		%rd,[%sp+imm7]	Stack (32 bits) → general-purpose register *
		%rd,[imm7]	Memory (32 bits) → general-purpose register *
		[%rb],%rs [%rb]+,%rs [%rb]-,%rs -[%rb],%rs	General-purpose register (32 bits, zero-extended) → memory * Memory address post-increment, post-decrement, and pre-decrement functions can be used.
		[%sp+imm7],%rs	General-purpose register (32 bits, zero-extended) → stack *
		[imm7],%rs	General-purpose register (32 bits, zero-extended) → memory *
		%rd,%sp	SP → general-purpose register
		%rd,%pc	PC → general-purpose register
		%rd,[%sp] %rd,[%sp]+ %rd,[%sp]- %rd,-[%sp]	Stack (32 bits) → general-purpose register * Stack pointer post-increment, post-decrement, and pre-decrement functions can be used.

Classification	Mnemonic		Function
Data transfer	ld.a	[%sp],%rs [%sp]+,%rs [%sp]-,%rs -[%sp],%rs	General-purpose register (32 bits, zero-extended) → stack * Stack pointer post-increment, post-decrement, and pre-decrement functions can be used.
		%sp,%rs	General-purpose register (24 bits) → SP
		%sp,imm7	Immediate → SP
Integer arithmetic operation	add add/c add/nc	%rd,%rs	16-bit addition between general-purpose registers Supports conditional execution (/c: executed if C = 1, /nc: executed if C = 0).
	add	%rd,imm7	16-bit addition of general-purpose register and immediate
	add.a add.a/c add.a/nc	%rd,%rs	24-bit addition between general-purpose registers Supports conditional execution (/c: executed if C = 1, /nc: executed if C = 0).
	add.a	%sp,%rs	24-bit addition of SP and general-purpose register
		%rd,imm7	24-bit addition of general-purpose register and immediate
		%sp,imm7	24-bit addition of SP and immediate
	adc adc/c adc/nc	%rd,%rs	16-bit addition with carry between general-purpose registers Supports conditional execution (/c: executed if C = 1, /nc: executed if C = 0).
	adc	%rd,imm7	16-bit addition of general-purpose register and immediate with carry
	sub sub/c sub/nc	%rd,%rs	16-bit subtraction between general-purpose registers Supports conditional execution (/c: executed if C = 1, /nc: executed if C = 0).
	sub	%rd,imm7	16-bit subtraction of general-purpose register and immediate
	sub.a sub.a/c sub.a/nc	%rd,%rs	24-bit subtraction between general-purpose registers Supports conditional execution (/c: executed if C = 1, /nc: executed if C = 0).
	sub.a	%sp,%rs	24-bit subtraction of SP and general-purpose register
		%rd,imm7	24-bit subtraction of general-purpose register and immediate
		%sp,imm7	24-bit subtraction of SP and immediate
	sbc sbc/c sbc/nc	%rd,%rs	16-bit subtraction with carry between general-purpose registers Supports conditional execution (/c: executed if C = 1, /nc: executed if C = 0).
	sbc	%rd,imm7	16-bit subtraction of general-purpose register and immediate with carry
	cmp cmp/c cmp/nc	%rd,%rs	16-bit comparison between general-purpose registers Supports conditional execution (/c: executed if C = 1, /nc: executed if C = 0).
	cmp	%rd,sign7	16-bit comparison of general-purpose register and immediate
	cmp.a cmp.a/c cmp.a/nc	%rd,%rs	24-bit comparison between general-purpose registers Supports conditional execution (/c: executed if C = 1, /nc: executed if C = 0).
	cmp.a	%rd,imm7	24-bit comparison of general-purpose register and immediate
	cmc cmc/c cmc/nc	%rd,%rs	16-bit comparison with carry between general-purpose registers Supports conditional execution (/c: executed if C = 1, /nc: executed if C = 0).
	cmc	%rd,sign7	16-bit comparison of general-purpose register and immediate with carry
Logical operation	and and/c and/nc	%rd,%rs	Logical AND between general-purpose registers Supports conditional execution (/c: executed if C = 1, /nc: executed if C = 0).
	and	%rd,sign7	Logical AND of general-purpose register and immediate
	or or/c or/nc	%rd,%rs	Logical OR between general-purpose registers Supports conditional execution (/c: executed if C = 1, /nc: executed if C = 0).
	or	%rd,sign7	Logical OR of general-purpose register and immediate
	xor xor/c xor/nc	%rd,%rs	Exclusive OR between general-purpose registers Supports conditional execution (/c: executed if C = 1, /nc: executed if C = 0).
	xor	%rd,sign7	Exclusive OR of general-purpose register and immediate
	not not/c not/nc	%rd,%rs	Logical inversion between general-purpose registers (1's complement) Supports conditional execution (/c: executed if C = 1, /nc: executed if C = 0).
	not	%rd,sign7	Logical inversion of general-purpose register and immediate (1's complement)

Classification	Mnemonic		Function
Shift and swap	sr	%rd,%rs	Logical shift to the right with the number of bits specified by the register
		%rd,imm7	Logical shift to the right with the number of bits specified by immediate
	sa	%rd,%rs	Arithmetic shift to the right with the number of bits specified by the register
		%rd,imm7	Arithmetic shift to the right with the number of bits specified by immediate
	sl	%rd,%rs	Logical shift to the left with the number of bits specified by the register
		%rd,imm7	Logical shift to the left with the number of bits specified by immediate
	swap	%rd,%rs	Bytewise swap on byte boundary in 16 bits
Immediate extension	ext	imm13	Extend operand in the following instruction
Conversion	cv.ab	%rd,%rs	Convert signed 8-bit data into 24 bits
	cv.as	%rd,%rs	Convert signed 16-bit data into 24 bits
	cv.al	%rd,%rs	Convert 32-bit data into 24 bits
	cv.la	%rd,%rs	Converts 24-bit data into 32 bits
	cv.ls	%rd,%rs	Converts 16-bit data into 32 bits
Branch	jpr jpr.d	sign10 %rb	PC relative jump Delayed branching possible
	jpa ipa.d	imm7 %rb	Absolute jump Delayed branching possible
	jrgt jrgt.d	sign7	PC relative conditional jump Branch condition: !Z & !(N ^ V) Delayed branching possible
	jrge jrge.d	sign7	PC relative conditional jump Branch condition: !(N ^ V) Delayed branching possible
	jrlt jrlt.d	sign7	PC relative conditional jump Branch condition: N ^ V Delayed branching possible
	jrle jrle.d	sign7	PC relative conditional jump Branch condition: Z | N ^ V Delayed branching possible
	jrugt jrugt.d	sign7	PC relative conditional jump Branch condition: !Z & !C Delayed branching possible
	jruge jruge.d	sign7	PC relative conditional jump Branch condition: !C Delayed branching possible
	jrult jrult.d	sign7	PC relative conditional jump Branch condition: C Delayed branching possible
	jrule jrule.d	sign7	PC relative conditional jump Branch condition: Z | C Delayed branching possible
	jreq jreq.d	sign7	PC relative conditional jump Branch condition: Z Delayed branching possible
	jrne jrne.d	sign7	PC relative conditional jump Branch condition: !Z Delayed branching possible
	call call.d	sign10 %rb	PC relative subroutine call Delayed call possible
	calla calla.d	imm7 %rb	Absolute subroutine call Delayed call possible
	ret ret.d		Return from subroutine Delayed return possible
	int	imm5	Software interrupt
	intl	imm5,imm3	Software interrupt with interrupt level setting
	reti reti.d		Return from interrupt handling Delayed call possible
	brk		Debug interrupt
	retd		Return from debug processing
System control	nop		No operation
	halt		HALT mode
	slp		SLEEP mode
	ei		Enable interrupts
	di		Disable interrupts
Coprocessor control	ld.cw	%rd,%rs %rd,imm7	Transfer data to coprocessor
	ld.ca	%rd,%rs %rd,imm7	Transfer data to coprocessor and get results and flag statuses
	ld.cf	%rd,%rs %rd,imm7	Transfer data to coprocessor and get flag statuses

* The `ld.a` instruction accesses memories in 32-bit length. During data transfer from a register to a memory, the 32-bit data in which the eight high-order bits are set to 0 is written to the memory. During reading from a memory, the eight high-order bits of the read data are ignored.

The symbols in the above table each have the meanings specified below.

Table 5.1.2 Symbol Meanings

Symbol	Description
`%rs`	General-purpose register, source
`%rd`	General-purpose register, destination
`[%rb]`	Memory addressed by general-purpose register
`[%rb]+`	Memory addressed by general-purpose register with address post-incremented
`[%rb]-`	Memory addressed by general-purpose register with address post-decremented
`-[%rb]`	Memory addressed by general-purpose register with address pre-decremented
`%sp`	Stack pointer
`[%sp],[%sp+imm7]`	Stack
`[%sp]+`	Stack with address post-incremented
`[%sp]-`	Stack with address post-decremented
`-[%sp]`	Stack with address pre-decremented
imm3,imm5,imm7,imm13	Unsigned immediate (numerals indicating bit length)
sign7,sign10	Signed immediate (numerals indicating bit length)

5.2 Addressing Modes (without `ext` extension)

The instruction set of the S1C17 Core has seven discrete addressing modes, as described below. The processor determines the addressing mode according to the operand in each instruction before it accesses data.
(1) Immediate addressing
(2) Register direct addressing
(3) Register indirect addressing
(4) Register indirect addressing with post-increment/post-decrement/pre-decrement
(5) Register indirect addressing with displacement
(6) Signed PC relative addressing
(7) PC absolute addressing

5.2.1 Immediate Addressing

The immediate included in the instruction code that is indicated as *immX* (unsigned immediate) or *signX* (signed immediate) is used as the source data. The immediate size specifiable in each instruction is indicated by a numeral in the symbol (e.g., *imm7* = unsigned 7 bits; *sign7* = signed 7 bits). For signed immediates such as *sign7*, the most significant bit is the sign bit, which is extended to 16 or 24 bits when the instruction is executed.

Example:
```
ld   %r0,0x70         ; Load 16-bit data
```

Before execution r0 = 0x*XXXXXX*
After execution r0 = 0x00fff0

The immediate *sign7* can represent values in the range of +63 to -64 (0b0111111 to 0b1000000).

Except in the case of shift-related instructions, immediate data can be extended to a maximum of 24 bits by a combined use of the operand value and the `ext` instruction.

Example:
```
ext    imm13   (1)
ext    imm13   (2)
ld.a   %r0,imm7       ; Load 24-bit data
```

r0 after execution

	23–20	19–7	6–0
r0	*imm13*(3:0) (1)	*imm13* (2)	*imm7*

5.2.2 Register Direct Addressing

The content of a specified register is used directly as the source data. Furthermore, if this addressing mode is specified as the destination for an instruction that loads the result in a register, the result is loaded in this specified register. The instructions that have the following symbols as the operand are executed in this addressing mode.

%rs *rs* is a metasymbol indicating the general-purpose register that holds the source data to be operated on or transferred. The register is actually written as `%r0`, `%r1`, ... or `%r7`.

%rd *rd* is a metasymbol indicating the general-purpose register that is the destination for the result of operation. The register is actually written as `%r0`, `%r1`, ... or `%r7`. Depending on the instruction, it will also be used as the source data.

Special register names are written as follows:
Stack pointer `%sp`
Program counter `%pc`

The register names are always prefixed by "`%`" to discriminate them from symbol names, label names, and the like.

5.2.3 Register Indirect Addressing

In this mode, memory is accessed indirectly by specifying a general-purpose register or the stack pointer that holds the address needed. This addressing mode is used only for load instructions that have `[%rb]` or `[%sp]` as the operand. Actually, this general-purpose register is written as `[%r0]`, `[%r1]`, ... `[%r7]`, or `[%sp]`, with the register name enclosed in brackets "`[]`."

The processor refers to the content of a specified register as the base address, and transfers data in the format that is determined by the type of load instruction.

Examples: Memory → Register

```
ld.b   %r0,[%r1]          ; Load 8-bit data
ld     %r0,[%r1]          ; Load 16-bit data
ld.a   %r0,[%r1]          ; Load 24-bit data
```

Register → Memory

```
ld.b   [%r1],%r0          ; Store 8-bit data
ld     [%r1],%r0          ; Store 16-bit data
ld.a   [%r1],%r0          ; Store 24-bit data
```

In this example, the address indicated by r1 is the memory address from or to which data is to be transferred.

In 16-bit and 24-bit transfers, the base address that is set in a register must be on a 16-bit boundary (least significant address bit = 0) or 32-bit boundary (2 low-order address bits = 0), respectively. Otherwise, an address-misaligned interrupt will be generated.

5.2.4 Register Indirect Addressing with Post-increment/decrement or Pre-decrement

As in register indirect addressing, the memory location to be accessed is specified indirectly by a general-purpose register or the stack pointer. In this addressing mode, the base address held in a specified register is incremented/decremented by an amount equal to the transferred data size before or after a data transfer. In this way, data can be read from or written to continuous addresses in memory only by setting the start address once at the beginning.

* Increment/decrement size (without `ext`)

Byte transfer (`ld.b`, `ld.ub`): *rb* → *rb* + 1, *rb* → *rb* - 1
16-bit transfer (`ld`): *rb* → *rb* + 2, *rb* → *rb* - 2
24-bit transfer (`ld.a`): *rb* → *rb* + 4, *rb* → *rb* - 4

Register indirect addressing with post-increment

When a data transfer finishes, the base address is incremented.
This addressing mode is specified by enclosing the register name in brackets "`[]`," which is then suffixed by "+."
The register name is actually written as `[%r0]+`, `[%r1]+`, ... `[%r7]+`, or `[%sp]+`.

Register indirect addressing with post-decrement

When a data transfer finishes, the base address is decremented.
This addressing mode is specified by enclosing the register name in brackets "`[]`," which is then suffixed by "`-`."
The register name is actually written as `[%r0]-`, `[%r1]-`, ... `[%r7]-`, or `[%sp]-`.

Register indirect addressing with pre-decrement

The base address is decremented before a data transfer starts.
This addressing mode is specified by enclosing the register name in brackets "`[]`," which is prefixed by "`-`."
The register name is actually written as `-[%r0]`, `-[%r1]`, ... `-[%r7]`, or `-[%sp]`.

5.2.5 Register Indirect Addressing with Displacement

In this mode, memory is accessed beginning with the address that is derived by adding a specified immediate (displacement) to the register content. Unless `ext` instructions are used, this addressing mode can only be used for load instructions that have `[%sp+`*`imm7`*`]` as the operand.

Example: `ld.b %r0,[%sp+0x10]`

The byte data at the address derived by adding 0x10 to the content of the current SP is loaded into the R0 register.

If `ext` instructions described in Section 5.3 are used, ordinary register indirect addressing (`[%`*`rb`*`]`) becomes a special addressing mode in which the immediate specified by the `ext` instruction constitutes the displacement.

Example: `ext` *`imm13`*

`ld.b %`*`rd`*`,[%`*`rb`*`]` The memory address to be accessed is "`%`*`rb`*`+`*`imm13`*."

5.2.6 Signed PC Relative Addressing

This addressing mode is used for the `jpr`, `jr*`, and `call` instructions that have a signed 7- or 10-bit immediate (*sign7*/*sign10*) or `%`*`rb`* in their operand. When these instructions are executed, the program branches to the address derived by twice adding the *sign7*/*sign10* value (16-bit boundary) or the *rb* register value to the current PC.

Example: PC + 0 `jrne 0x04` The program branches to the PC + 8 address when the `jrne` branch condition holds true.

: : (PC + 0) + 0x04 ∗ 2 → PC + 8

:

PC + 8

5.2.7 PC Absolute Addressing

This addressing mode is used for the `jpa`, and `calla` instructions that have an unsigned 7-bit immediate (*imm7*) or `%`*`rb`* in their operand. When these instructions are executed, the program directly branches to the address specified with the *imm7* or *rb* register value by loading the value to the PC. Also this addressing mode is used for the `int` and `intl` instructions that execute interrupt handler routines.

Example: `int 0x03` Executes the interrupt handler of vector No. 3 (TTBR + 0xc).

5.3 Addressing Modes with `ext`

The immediate specifiable in 16-bit, fixed-length instruction code is specified in a bit field of a 7- or 10-bit length, depending on the instruction used. The `ext` instructions are used to extend the size of this immediate.
The `ext` instructions are used in combination with data transfer, arithmetic/logic, or branch instructions, and is placed directly before the instruction whose immediate needs to be extended. The instruction is expressed in the form `ext` *`imm13`*, in which the immediate size extendable by one `ext` instruction is 13 bits and up to two `ext` instructions can be written in succession to extend the immediate further.
The `ext` instructions are effective only for the instructions for which the immediate extension written directly after `ext` is possible, and have no effect for all other instructions. When three or more `ext` instructions have been described sequentially, the last two are effective and others are ignored.
When an instruction, which does not support the extension in the `ext` instruction, follows an `ext`, the `ext` instruction will be executed as a `nop` instruction.

5.3.1 Extension of Immediate Addressing

Extension of *imm7*

The *imm7* immediate is extended to a 16-, 20-, or 24-bit immediate.

Extending to a 16-bit immediate

To extend the immediate to 16-bit quantity, enter one `ext` instruction directly before the target instruction.

Example:
```
ext   imm13
add   %rd,imm7     ; = add   %rd,imm16
```

Extended immediate

15 – 7	6 – 0
imm13(8:0)	*imm7*

Extending to a 20-bit immediate

To extend the immediate to 20-bit quantity, enter one `ext` instruction directly before the target instruction.

Example:
```
ext      imm13
add.a    %rd,imm7   ; = add.a   %rd,imm20
```

Extended immediate

23 – 20	19 – 7	6 – 0
0 0 0 0	*imm13*	*imm7*

Bits 23–20 are filled with 0 (zero-extension).

Extending to a 24-bit immediate

To extend the immediate to 24-bit quantity, enter two `ext` instructions directly before the target instruction.

Example:
```
ext   imm13          (1)
ext   imm13          (2)
ld    %rd,[imm7]   ; = ld    %rd,[imm24]
```

Extended immediate

23 – 20	19 – 7	6 – 0
imm13(3:0) (1)	*imm13* (2)	*imm7*

Extension of *sign7*

The *sign7* immediate is extended to a 16-bit immediate.

Extending to a 16-bit immediate

To extend the immediate to 16-bit quantity, enter one `ext` instruction directly before the target instruction.

Example:
```
ext   imm13
ld    %rd,sign7
```

Extended immediate

15	14–7	6–0
S	*imm13*(7:0)	*sign7*

Bit 8 of the *imm13* in the `ext` instruction is the sign, with the immediate extended to become signed 16-bit data. The most significant bit in *sign7* is handled as the MSB data of 7-bit data, and not as the sign.

5.3.2 Extension of Register Direct Addressing

Extending register-to-register operation instructions

Register-to-register operation instructions are extended by one or two `ext` instructions. Unlike data transfer instructions, these instructions add or subtract the content of the *rs* register and the immediate specified by an `ext` instruction according to the arithmetic operation to be performed. They then store the result in the *rd* register. The content of the *rd* register does not affect the arithmetic operation performed. An example of how to extend for an add operation is shown below.

Extending to *rs* + *imm13* (for 16-bit and 24-bit operation instructions)

To extend to *rs* + *imm13*, enter one `ext` instruction directly before the target instruction.

Example:
```
ext     imm13
add.a   %rd,%rs
```

If not extended, *rd* = *rd* + *rs*

When extended by one `ext` instruction, *rd* = *rs* + *imm13*

	23–13	12–0
rs	Data	
+		
Immediate	0 0 0 0 0 0 0 0 0 0 0	*imm13*
↓		
rd	Data + *imm13*	

Extending to *rs* + *imm16* (for 16-bit operation instructions)

To extend to *rs* + *imm16*, enter two `ext` instructions directly before the target instruction.

Example:
```
ext   imm13     (1)
ext   imm13     (2)
add   %rd,%rs
```

If not extended, *rd* = *rd* + *rs*

When extended by two `ext` instructions, *rd* = *rs* + *imm16*

	23–16	15–13	12–0
rs	X X X X X X X X	Data	
+			
Immediate		(1) *imm13*(2:0)	*imm13* (2)
↓			
rd	0 0 0 0 0 0 0 0	Data + *imm16*	

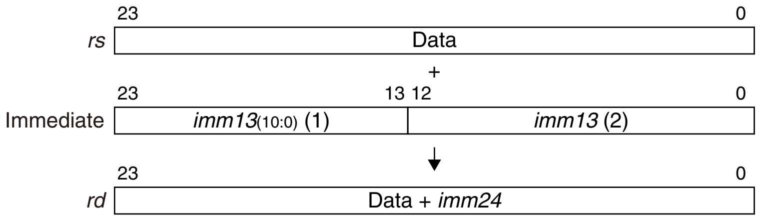
Extending to *rs* + *imm24* (24-bit operation instructions)

To extend to *rs* + *imm24*, enter two `ext` instructions directly before the target instruction.

Example:
```
ext     imm13       (1)
ext     imm13       (2)
add.a   %rd,%rs
```

If not extended, *rd* = *rd* + *rs*

When extended by two `ext` instructions, *rd* = *rs* + *imm24*



5.3.3 Extension of Register Indirect Addressing

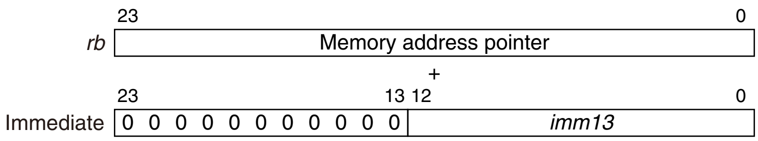
Adding displacement to `[%rb]`

Memory is accessed at the address derived by adding the immediate specified by an `ext` instruction to the address that is indirectly referenced by `[%rb]`.

Adding a 13-bit immediate

Memory is accessed at the address derived by adding the 13-bit immediate specified by *imm13* to the address specified by the *rb* register. During address calculation, *imm13* is zero-extended to 24-bit quantity.

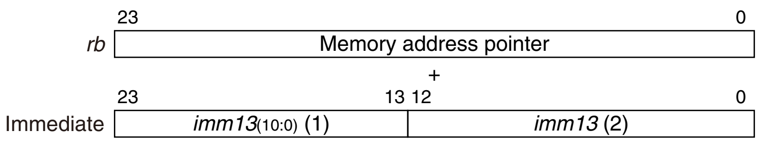
Example:
```
ext     imm13
ld.b    %rd,[%rb]    ; = ld.b   %rd,[%rb+imm13]
```



Adding a 24-bit immediate

Memory is accessed at the address derived by adding the 24-bit immediate specified by *imm24* to the address specified by the *rb* register.

Example:
```
ext     imm13        (1)
ext     imm13        (2)
ld.b    %rd,[%rb]    ; = ld.b   %rd,[%rb+imm24]
```



5.3.4 Extension of Register Indirect Addressing with Displacement

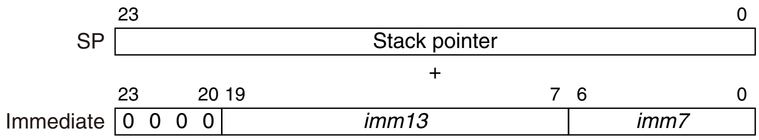
Extending `[%sp+imm7]` displacement

The immediate (*imm7*) in displacement-added register indirect addressing instructions is extended.
The extended data and the SP are added to comprise the source or destination address of transfer.

Extending to a 20-bit immediate

To extend the immediate to 20-bit quantity, enter one `ext` instruction directly before the target instruction.

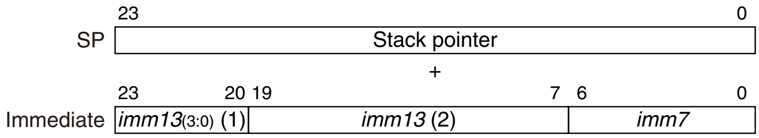
Example:
```
ext   imm13
ld    %rd,[%sp+imm7]     ; = ld   %rd,[%sp+imm20]
```



Extending to a 24-bit immediate

To extend the immediate to 24-bit quantity, enter two `ext` instructions directly before the target instruction.

Example:
```
ext   imm13 (1)
ext   imm13 (2)
ld    %rd,[%sp+imm7]     ; = ld   %rd,[%sp+imm24]
```



5.3.5 Extension of Signed PC Relative Addressing

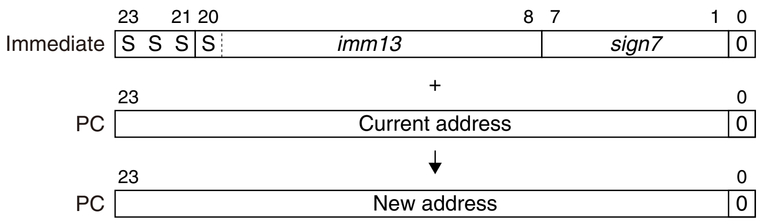
Extending the displacement of PC relative branch instructions

The *sign7* immediate in PC relative branch instructions is extended to a signed 21-bit or a signed 24-bit immediate. The *sign7* immediate in PC relative branch instructions is multiplied by 2 for conversion to a relative value for the jump address, and the derived value is then added to PC to determine the jump address. The `ext` instructions extend this relative jump address value.

Extending to a 21-bit immediate

To extend the *sign7* immediate to a 21-bit immediate, enter one `ext` instruction directly before the target instruction.

Example:
```
ext    imm13
jrgt   sign7       ; = jrgt  sign21
```

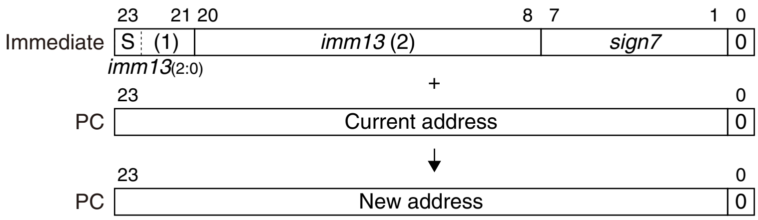


The most significant bit "S" in the immediate that has been extended by the `ext` instruction is the sign, with which bits 23–21 are extended to become signed 21-bit data. The most significant bit in *sign7* is handled as the MSB data of 7-bit data, and not as the sign.

Extending to a 24-bit immediate

To extend the *sign7* immediate to a 24-bit immediate, enter two `ext` instructions directly before the target instruction.

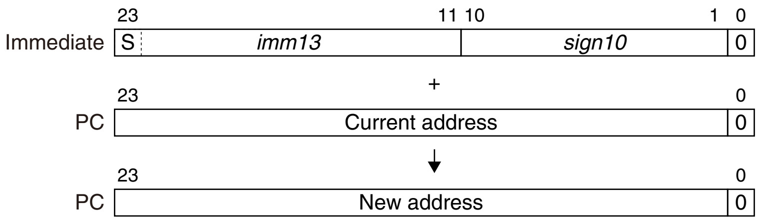
Example:
```
ext    imm13      (1)
ext    imm13      (2)
jrgt   sign7      ; = jrgt  sign24
```



The most significant bit "S" in the immediate that has been extended by `ext` instructions is the sign. Bits 12–3 in the first `ext` instruction are unused.

Also the *sign10* operand in the `jpr` and `call` instructions can be extended to 24-bit quantity using one `ext` instruction.

Example:
```
ext    imm13
call   sign10     ; = call  sign24
```



5.3.6 Extension of PC Absolute Addressing

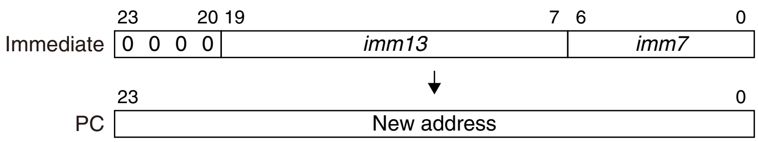
Extending the branch destination address

The *imm7* immediate is extended to a 20- or 24-bit immediate.

Extending to a 20-bit immediate

To extend the immediate to 20-bit quantity, enter one `ext` instruction directly before the target instruction.

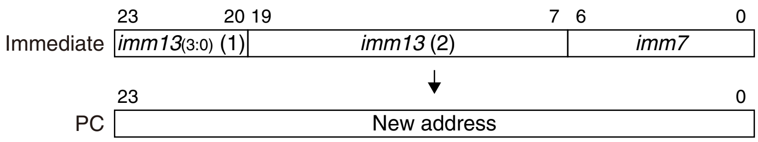
Example:
```
ext     imm13
calla   imm7      ; = calla  imm20
```



Extending to a 24-bit immediate

To extend the immediate to 24-bit quantity, enter two `ext` instructions directly before the target instruction.

Example:
```
ext   imm13       (1)
ext   imm13       (2)
jpa   imm7        ; = jpa  imm24
```



5.4 Data Transfer Instructions

The transfer instructions in the S1C17 Core support data transfer between one register and another, as well as between a register and memory. A transfer data size and data extension format can be specified in the instruction code. In mnemonics, this specification is classified as follows:

`ld.b`	Signed byte data transfer
`ld.ub`	Unsigned byte data transfer
`ld`	16-bit data transfer
`ld.a`	24/32-bit data transfer

In signed byte transfers to registers, the source data is sign-extended to 16 bits. In unsigned byte transfers, the source data is zero-extended to 16 bits.
In transfers in which data is transferred from registers, data of a specified size on the lower side of the register is the data to be transferred.

If the destination of transfer is a general-purpose register, the register content after a transfer is as follows:

Signed byte data transfer

	23–16	15–8	7–0
rd	0 0 0 0 0 0 0 0	S S S S S S S S	S Byte data

Extended with the sign in bit 7 of the byte data

Unsigned byte data transfer

	23–16	15–8	7–0
rd	0 0 0 0 0 0 0 0	0 0 0 0 0 0 0 0	Byte data

16-bit data transfer

	23–16	15–0
rd	0 0 0 0 0 0 0 0	16-bit data

24/32-bit data transfer

	23–0
rd	24-bit data

Refer to Chapter 3, “Data Formats,” for the data layout in the memory.

5.5 Logical Operation Instructions

Four discrete logical operation instructions are available for use with the S1C17 Core.

and	Logical AND
or	Logical OR
xor	Exclusive-OR
not	Logical NOT

All logical operations are performed in a specified general-purpose register (R0–R7). The source is one of two, either 16-bit data in a specified general-purpose register or immediate data (7, 13, or 16 bits).
When a logical operation is performed, the V flag (bit 2) in the PSR is cleared.

Conditional execution

The logical operation instructions for between registers (*op %rd, %rs*) allow use of the switches to specify whether the instruction will be executed or not depending on the C flag status.

Unconditional execution instructions

op %rd, %rs (*op* = and, or, xor, not)

The instruction without a switch will be always executed regardless how the C flag is set.
Example: and *%rd, %rs*

Instructions executable under C condition

op/c *%rd, %rs* (*op* = and, or, xor, not)

The instruction with the /c switch will be executed only when the C flag has been set to 1.
Example: or/c *%rd, %rs*

Instructions executable under NC condition

op/nc *%rd, %rs* (*op* = and, or, xor, not)

The instruction with the /nc switch will be executed only when the C flag has been cleared to 0.
Example: xor/nc *%rd, %rs*

5.6 Arithmetic Operation Instructions

The instruction set of the S1C17 Core supports add/subtract, and compare instructions for arithmetic operations.

add	16-bit addition
add.a	24-bit addition
adc	16-bit addition with carry
sub	16-bit subtraction
sub.a	24-bit subtraction
sbc	16-bit subtraction with borrow
cmp	16-bit comparison
cmp.a	24-bit comparison
cmc	16-bit comparison with borrow

The above arithmetic operations are performed between one general-purpose register and another (R0–R7), or between a general-purpose register and an immediate. Furthermore, the `add.a` and `sub.a` instructions can perform operations between the SP and a general-purpose register/immediate. Immediates in sizes smaller than the operation unit (16 bits or 24 bits), except for the `cmp` instruction, are zero-extended when operation is performed.
The `cmp` instruction compares two operands, and may alter a flag, depending on the comparison result. Basically, it is used to set conditions for conditional jump instructions. If an immediate smaller than operation unit in size is specified as the source, it is sign-extended when comparison is performed.

Conditional execution

The arithmetic operation instructions for between registers (*op %rd*, *%rs*) allow use of the switches to specify whether the instruction will be executed or not depending on the C flag status.

Unconditional execution instructions

op *%rd*, *%rs* (*op* = add, add.a, adc, sub, sub.a, sbc, cmp, cmp.a, cmc)

The instruction without a switch will be always executed regardless how the C flag is set.

Example: `add %rd,%rs`

Instructions executable under C condition

op/c *%rd*, *%rs* (*op* = add, add.a, adc, sub, sub.a, sbc, cmp, cmp.a, cmc)

The instruction with the `/c` switch will be executed only when the C flag has been set to 1.

Example: `sub/c %rd,%rs`

Instructions executable under NC condition

op/nc *%rd*, *%rs* (*op* = add, add.a, adc, sub, sub.a, sbc, cmp, cmp.a, cmc)

The instruction with the `/nc` switch will be executed only when the C flag has been cleared to 0.

Example: `cmp/nc %rd,%rs`

5.7 Shift and Swap Instructions

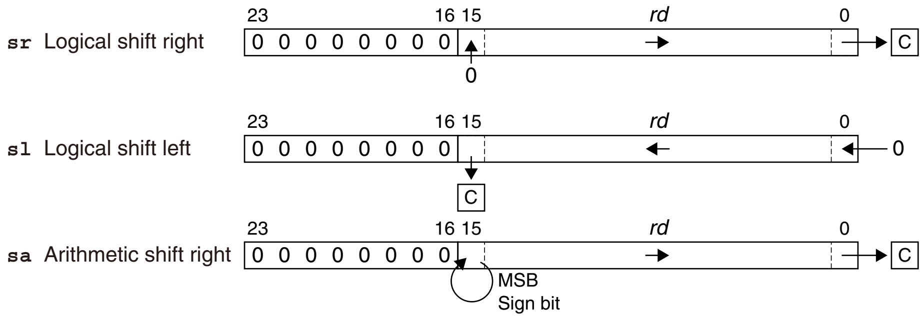
The S1C17 Core supports instructions to shift or swap the register data.

sr	Logical shift right
sl	Logical shift left (= Arithmetic shift left)
sa	Arithmetic shift right
swap	Swap upper and lower bytes

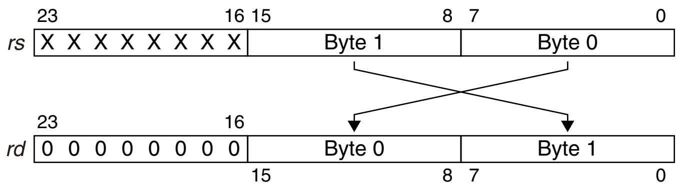
The shift operation is effective for bits 15 to 0 in the specified register and bits 23 to 16 are set to 0.
The number of bits to be shifted can be specified to 0–3 bits, 4 bits, or 8 bits using the operand *imm5* or the *rs* register.

%*rs*/*imm7* = 0–3: Shift 0 to 3 bits
%*rs*/*imm7* = 4–7: Shift 4 bits (fixed)
%*rs*/*imm7* = 8 or more: Shift 8 bits (fixed)

Example: `sr`	`%rd,1`	Bits 15–0 in `%rd` logically shifted one bit to the right
`sl`	`%rd,7`	Bits 15–0 in `%rd` logically shifted four bits to the left
`sa`	`%rd,0xf`	Bits 15–0 in `%rd` arithmetically shifted eight bits to the right



The `swap` instruction replaces the contents of general-purpose registers with each other, as shown below.



5.8 Branch and Delayed Branch Instructions

5.8.1 Types of Branch Instructions

(1) PC relative jump instructions

PC relative jump instructions include the following:

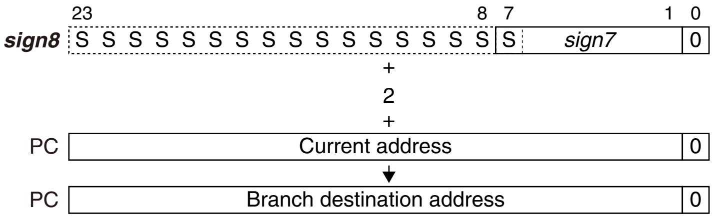
```
jr*   sign7
jpr   sign10
jpr   %rb
```

PC relative jump instructions are provided for relocatable programming, so that the program branches to the address calculated as PC + 2 (the next address of the branch instruction) + signed displacement (specified by the operand).
The number of instruction steps to the jump address is specified for *sign7/10* or *rb*. However, since the instruction length in the S1C17 Core is fixed to 16 bits, the value of *sign7/10* is doubled to become a word address in 16-bit units. Therefore, the displacement actually added to the PC is a signed 8-bit/11-bit quantity derived by doubling *sign7/10* (least significant bit always 0). When the *rb* register is used to specify the displacement, the register contents are added to the PC without doubling.
The specifiable displacement can be extended by the `ext` instruction, as shown below.

For branch instructions used singly

`jr*` *`sign7`* Functions as "`jr*` *`sign8`*" (*sign8* = {*sign7*, 0})

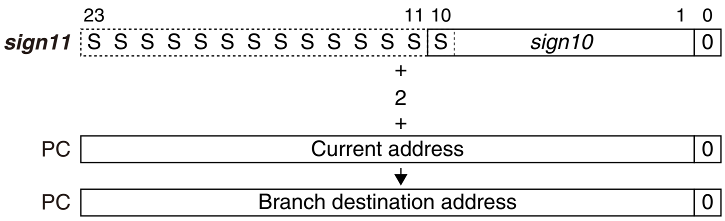
For the `jr*` instructions that are used singly, a signed 7-bit displacement (*sign7*) can be specified.



The range of addresses to which jumped is (PC - 126) to (PC + 128).

`jpr` *`sign10`* Functions as "`jpr` *`sign11`*" (*sign11* = {*sign10*, 0})

For the `jpr` instruction that is used singly, a signed 10-bit displacement (*sign10*) can be specified.

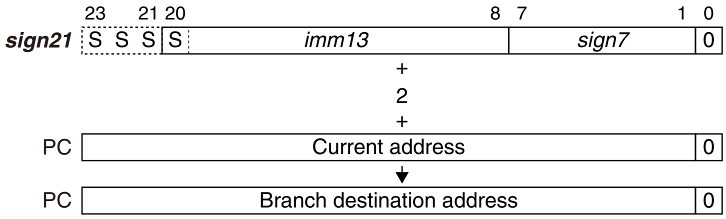


The range of addresses to which jumped is (PC - 2,046) to (PC + 2,048).

When extended by one `ext` instruction

`ext imm13`
`jr* sign7` Functions as "`jr* sign21`" (*sign21* = {*imm13*, *sign7*, 0})

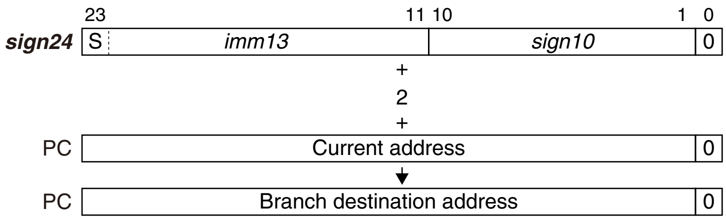
The *imm13* specified by the `ext` instruction is extended as the 13 high-order bits of *sign21*.



The range of addresses to which jumped is (PC - 1,048,574) to (PC + 1,048,576).

`ext imm13`
`jpr sign10` Functions as "`jpr sign24`" (*sign24* = {*imm13*, *sign10*, 0})

The *imm13* specified by the `ext` instruction is extended as the 13 high-order bits of *sign24*.



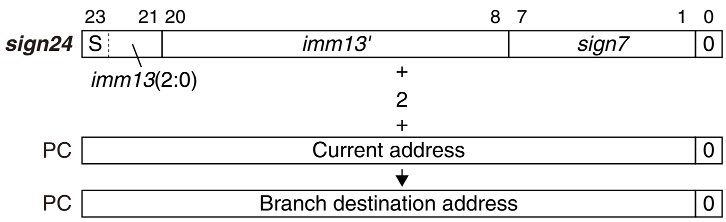
The range of addresses to which jumped is (PC - 8,388,606) to (PC + 8,388,608).

When extended by two `ext` instructions

`ext imm13`
`ext imm13'`
`jr* sign7` Functions as "`jr* sign24`"

The *imm13* specified by the first `ext` instruction is effective for only 3 bits, from bit 2 to bit 0 (with the 10 high-order bits ignored), so that *sign24* is configured as follows:

sign24 = {*imm13*(2:0), *imm13*', *sign7*, 0}



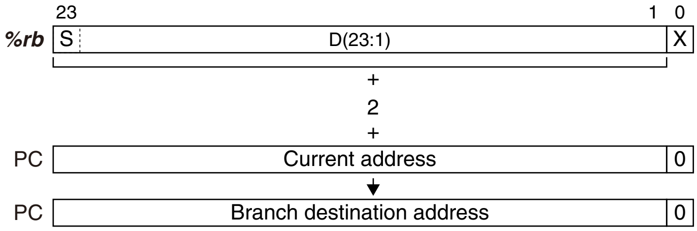
The range of addresses to which jumped is (PC - 8,388,606) to (PC + 8,388,608).

The above range of addresses to which jumped is a theoretical value, and is actually limited by the range of memory areas used.

For `jpr %rb`

`jpr %rb`

A signed 24-bit relative value is specified for *rb*.
The jump address is configured as follows:
{*rb*(23:1), 0}



The least significant bit in the *rb* register is always handled as 0.
The range of addresses to which jumped is (PC - 8,388,606) to (PC + 8,388,608).

The above range of addresses to which jumped is a theoretical value, and is actually limited by the range of memory areas used.

Branch conditions

The `jpr` instruction is an unconditional jump instruction that always cause the program to branch.
Instructions with names beginning with `jr` are conditional jump instructions for which the respective branch conditions are set by a combination of flags, so that only when the conditions are satisfied do they cause the program to branch to a specified address. The program does not branch unless the conditions are satisfied.
The conditional jump instructions basically use the result of the comparison of two values by the `cmp` instruction to determine whether to branch. For this reason, the name of each instruction includes a character that represents relative magnitude.
The types of conditional jump instructions and branch conditions are listed in Table 5.8.1.1.

Table 5.8.1.1 Conditional Jump Instructions and Branch Conditions

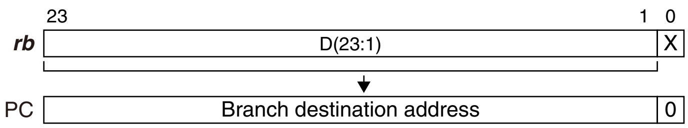
Instruction		Flag condition	Comparison of A:B	Remark
`jrgt`	Greater Than	!Z & !(N ^ V)	A > B	Used to compare signed data
`jrge`	Greater or Equal	!(N ^ V)	A ≥ B	
`jrlt`	Less Than	N ^ V	A < B	
`jrle`	Less or Equal	Z | (N ^ V)	A ≤ B	
`jrugt`	Unsigned, Greater Than	!Z & !C	A > B	Used to compare unsigned data
`jruge`	Unsigned, Greater or Equal	!C	A ≥ B	
`jrult`	Unsigned, Less Than	C	A < B	
`jrule`	Unsigned, Less or Equal	Z | C	A ≤ B	
`jreq`	Equal	Z	A = B	
`jrne`	Not Equal	!Z	A ≠ B	

Comparison of A:B made when "`cmp A,B`"

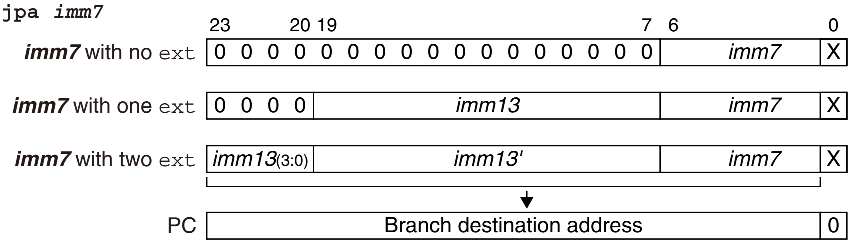
(2) Absolute jump instructions

The absolute jump instruction `jpa` causes the program to unconditionally branch to the location indicated by the content of a specified general-purpose register (*rb*) or an immediate *imm7* (can be extended to *imm20* or *imm24* using the `ext` instruction) as the absolute address. When the content of the *rb* register or the immediate is loaded into the PC, its least significant bit is always made 0.

`jpa %rb`



`jpa imm7`



(3) PC relative call instructions

The PC relative call instruction `call` *`sign10`*/*`%rb`* is a subroutine call instruction that is useful for relocatable programming, as it causes the program to unconditionally branch to a subroutine starting from an address calculated as PC + 2 (the next address of the branch instruction) + signed displacement (specified by the operand). During branching, the program saves the address of the instruction next to the `call` instruction (for delayed branching, the address of the second instruction following `call`) to the stack as the return address. When the `ret` instruction is executed at the end of the subroutine, this address is loaded into the PC, and the program returns to it from the subroutine.

Note that because the instruction length is fixed to 16 bits, the least significant bit of the displacement is always handled as 0 (*sign10* doubled, *rb* is not doubled), causing the program to branch to an even address.

As with the PC relative jump instructions, the specifiable displacement can be extended by the `ext` instruction. For details on how to extend the displacement, refer to the "(1) PC relative jump instructions."

(4) Absolute call instructions

The absolute call instruction `calla` causes the program to unconditionally call a subroutine starting from the location indicated by the content of a specified general-purpose register (*rb*) or an immediate *imm7* (can be extended to *imm20* or *imm24* using the `ext` instruction) as the absolute address. When the content of the *rb* register or the immediate is loaded into the PC, its least significant bit is always made 0. (Refer to the "(2) Absolute jump instructions.")

(5) Software interrupts

The software interrupts `int` and `intl` are the instructions that cause the software to generate an interrupt with the vector numbers specified by the operand *imm5*, by which a specified interrupt handler routine can be executed. When a software interrupt occurs, the processor saves the PSR and the instruction address next to `int`/`intl` to the stack, and reads the specified vector from the vector table in order to execute an interrupt handler routine. Therefore, to return from the interrupt handler routine, the `reti` instruction must be used, as it restores the PSR as well as the PC from the stack. For details on the software interrupt, refer to Section 6.3, "Interrupts."

(6) Return instructions

The `ret` instruction, which is a return instruction for the `call` and `calla` instructions, loads the saved return address from the stack into the PC as it terminates the subroutine. Therefore, the value of the SP when the `ret` instruction is executed must be the same as when the subroutine was executed (i.e., one that indicates the return address).

The `reti` instruction is a return instruction for the interrupt handler routine. Since the PSR is saved to the stack along with the return address in interrupt handling, the content of the PSR must be restored from the stack using the `reti` instruction. In the `reti` instruction, the PC and the PSR are read out of the stack in that order. As in the case of the `ret` instruction, the value of the SP when the `reti` instruction is executed must be the same as when the subroutine was executed.

(7) Debug interrupts

The `brk` and `retd` instructions are used to call a debug interrupt handler routine, and to return from that routine. Since these instructions are basically provided for the debug firmware, please do not use them in application programs. For details on the functionality of these instructions, refer to Section 6.5, "Debug Circuit."

5.8.2 Delayed Branch Instructions

The S1C17 Core uses pipelined instruction processing, in which instructions are executed while other instructions are being fetched. In a branch instruction, because the instruction that follows it has already been fetched when it is executed, the execution cycles of the branch instruction can be reduced by one cycle by executing the prefetched instruction before the program branches. This is referred to as a delayed branch function, and the instruction executed before branching (i.e., the instruction at the address next to the branch instruction) is referred to as a delayed slot instruction.
The delayed branch function can be used in the instructions listed below, which in mnemonics is identified by the extension "`.d`" added to the branch instruction name.

Delayed branch instructions

`jrgt.d` **`jrge.d`** **`jrlt.d`** **`jrle.d`** **`jrugt.d`** **`jruge.d`** **`jrult.d`**
`jrule.d` **`jreq.d`** **`jrne.d`** **`call.d`** **`calla.d`** **`jpr.d`** **`jpa.d`**
`ret.d` **`reti.d`**

Delayed slot instructions

All instructions other than those listed below can be used as a delayed slot instruction.

Instructions that cannot be used as a delayed slot instruction

`brk call calla ext halt int intl jpa jpr jr* ret retd reti slp`

The `ext` instruction cannot be used to expand the operand of delayed slot instructions.
A delayed slot instruction is always executed regardless of whether the delayed branch instruction used is conditional or unconditional and whether it branches.
In "non-delayed" branch instructions (those not followed by the extension "`.d`"), the instruction at the address next to the branch instruction is not executed if the program branches; however, if it is a conditional jump and the program does not branch, the instruction at the next address is executed as the one that follows the branch instruction.

The return address saved to the stack by the `call.d` or `calla.d` instruction becomes the address for the next instruction following the delayed slot instruction, so that the delayed slot instruction is not executed when the program returns from the subroutine.

No interrupts occur in between a delayed branch instruction and a delayed slot instruction, as they are masked out by hardware.

Application for leaf subroutines

The following shows an example application of delayed branch instructions for achieving a fast leaf subroutine call.

Example:

```
        jpr.d   SUB       ; Jumps to a subroutine by a delayed branch instruction
        ld.a    %r7,%pc   ; Loads the return address into a general-purpose register by
                          ; a delayed slot instruction
        add.a   %r1,%r2   ; Return address
          :       :
SUB:
          :       :
        jpr     %r7       ; Return
```

Notes:
- The `ld.a` *`%rd`*`,%pc` instruction must be executed as a delayed slot instruction. If it does not follow a delayed branch instruction, the PC value that is loaded into the *rd* register may not be the next instruction address to the `ld.a` instruction.
- The delayed branch instruction listed below can only be used with the `ld.a` *`%rd`*`,%pc` delayed slot instruction.
  - `jpr.d` *`%rb/sign10`*
  - `jr*.d` *`sign7`*
  - `jpa.d` *`%rb/imm7`*

5.9 System Control Instructions

The following five instructions are used to control the system.

`nop`	Only increments the PC, with no other operations performed
`halt`	Places the processor in HALT mode
`slp`	Places the processor in SLEEP mode
`ei`	Enables interrupts
`di`	Disables interrupts

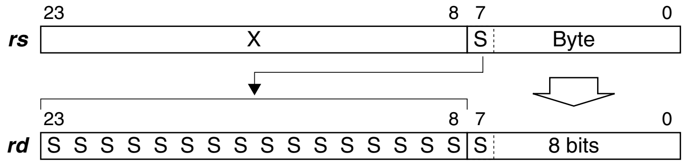
For details on HALT and SLEEP modes, refer to Section 6.4, "Power-Down Mode," and the Technical Manual for each S1C17 model.
For details on the interrupt control, refer to Section 6.3, "Interrupts."

5.10 Conversion Instructions

The 8/16/24/32 data conversion instructions listed below are provided for supporting C compiler.

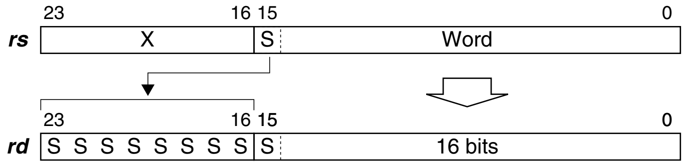
`cv.ab %rd,%rs`

Converts Byte data (8 bits) into 24-bit data with sign extended.



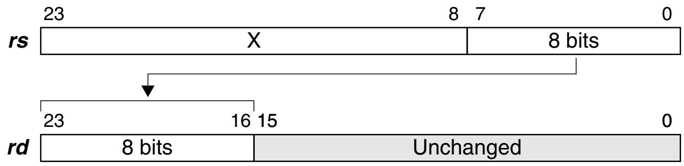
`cv.as %rd,%rs`

Converts 16-bit data into 24-bit data with sign extended.



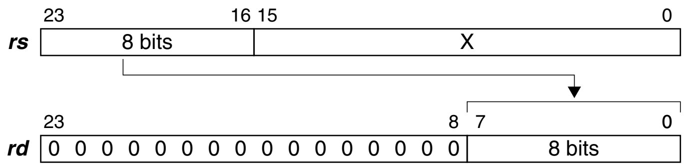
`cv.al %rd,%rs`

Extracts the high-order 8 bits to convert 32-bit data into 24-bit data.



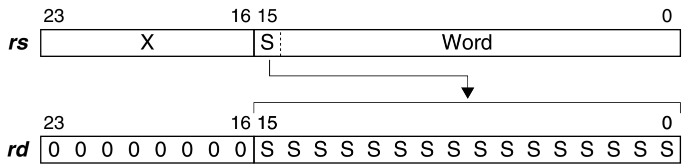
`cv.la %rd,%rs`

Extracts the high-order 8 bits to convert 24-bit data into 32-bit data.



`cv.ls %rd,%rs`

Extends the sign to convert 16-bit data into 32-bit data.



5.11 Coprocessor Instructions

The S1C17 Core incorporates a coprocessor interface and provides the dedicated coprocessor instructions listed below.

`ld.cw`	Transfer data to the coprocessor
`ld.ca`	Transfer data and input the results and flag status to/from the coprocessor
`ld.cf`	Input flag status from the coprocessor

The `ld.cw` and `ld.ca` instructions send two 24-bit data set in the *rd* (data 0) and *rs* (data 1) registers to the coprocessor. Data 1 can also be specified in an immediate *imm7*. In this case, the 7-bit immediate can be extended into *imm20* or *imm24* using the `ext` instruction.
The `ld.ca` instruction inputs the results from the coprocessor to the *rd* register.
The `ld.ca` and `ld.cf` instructions input the flag status from the coprocessor and set it to the PSR (C, V, Z, and N flags).
The concrete commands and status of the coprocessor vary with each coprocessor connected to the chip. Refer to the user's manual for the coprocessor used.

6 Functions

This chapter describes the processing status of the S1C17 Core and outlines the operation.

6.1 Transition of the Processor Status

The diagram below shows the transition of the operating status in the S1C17 Core.

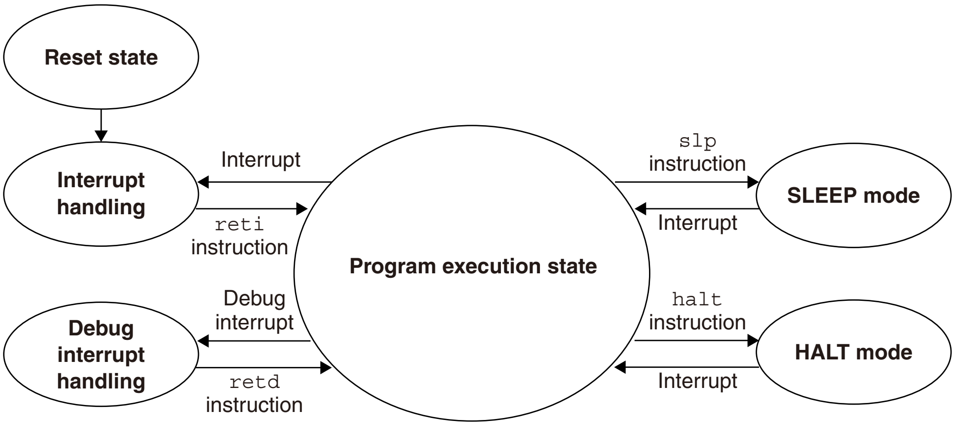


Figure 6.1.1 Processor Status Transition Diagram

6.1.1 Reset State

The processor is initialized when the reset signal is asserted, and then starts processing from the reset vector when the reset signal is deasserted.

6.1.2 Program Execution State

This is a state in which the processor executes the user program sequentially. The processor state transits to another when an interrupt occurs or the `slp` or `halt` instruction is executed.

6.1.3 Interrupt Handling

When a software or other interrupt occurs, the processor enters an interrupt handling state. The following are the possible causes of the need for interrupt handling:

(1) External interrupt
(2) Software interrupt
(3) Address misaligned interrupt
(4) NMI

6.1.4 Debug Interrupt

The S1C17 Core incorporates a debugging assistance facility to increase the efficiency of software development. To use this facility, a dedicated mode known as “debug mode” is provided. The processor can be switched from user mode to this mode by the `brk` instruction or a debug interrupt. The processor does not normally enter this mode.

6.1.5 HALT and SLEEP Modes

The processor is placed in HALT or SLEEP mode to reduce power consumption by executing the `halt` or `slp` instruction in the software (see Section 6.4). Normally the processor can be taken out of HALT or SLEEP mode by NMI or an external interrupt as well as initial reset.

6.2 Program Execution

Following initial reset, the processor loads the reset vector (address of the reset handler routine) into the PC and starts executing instructions beginning with the address. As the instructions in the S1C17 Core are fixed to 16 bits in length, the PC is incremented by 2 each time an instruction is fetched from the address indicated by the PC. In this way, instructions are executed successively.

When a branch instruction is executed, the processor checks the PSR flags and whether the branch conditions have been satisfied, and loads the jump address into the PC.

When an interrupt occurs, the processor loads the address for the interrupt handler routine from the vector table into the PC.
The vector table contains interrupt vectors beginning with the reset vector and is located from the address set in the TTBR register (0xffff80). The start address can be set to the TTBR in the configuration.

6.2.1 Instruction Fetch and Execution

Internally in the S1C17 Core, instructions are processed in three pipelined stages, so that the basic instructions except for the branch instructions and data transfer instructions with the memory address increment/decrement function can be executed in one clock cycle.
Pipelining speeds up instruction processing by executing one instruction while fetching another. In the 3-stage pipeline, each instruction is processed in three stages, with processing of instructions occurring in parallel, for faster instruction execution.

Basic instruction stages

Instruction fetch	Instruction decode	Instruction execution / Memory access / Register write

Hereinafter, each stage is represented by the following symbols:
F (for Fetch): Instruction fetch
D (for Decode): Instruction decode
E (for Execute): Instruction execution, memory access, register write

Pipelined operation

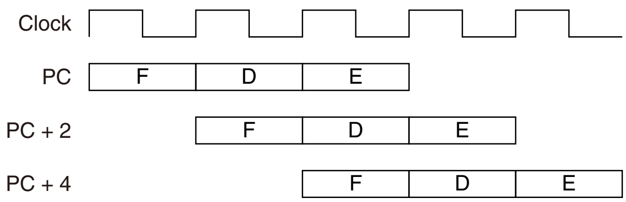


Figure 6.2.1.1 Pipelined Operation

Note: The pipelined operation shown above uses the internal memory. If external memory or low-speed external devices are used, one or more wait cycles may be inserted depending on the devices used, with the E stage kept waiting.

6.2.2 Execution Cycles and Flags

The following shows the number of cycles required for executing each instruction in a 1-cycle accessible memory connected to the Harvard bus and the flag change status.
Depending on the model, clock cycles spent by the external bus arbiter and wait cycles inherent in the external devices may be added.

Table 6.2.2.1 Number of Instruction Execution Cycles and Flag Status

Classification	Mnemonic		Cycle	Flag						Remark
				IL	IE	C	V	Z	N	
Data transfer	ld.b	%rd,%rs	1	–	–	–	–	–	–	*1: 1 cycle when ext is not used 2 cycles when ext is used
		%rd,[%rb]	1–2*1	–	–	–	–	–	–	
		%rd,[%rb]+	2	–	–	–	–	–	–	
		%rd,[%rb]-	2	–	–	–	–	–	–	
		%rd,-[%rb]	2	–	–	–	–	–	–	
		%rd,[%sp+imm7]	2	–	–	–	–	–	–	
		%rd,[imm7]	1	–	–	–	–	–	–	
		[%rb],%rs	1–2*1	–	–	–	–	–	–	
		[%rb]+,%rs	2	–	–	–	–	–	–	
		[%rb]-,%rs	2	–	–	–	–	–	–	
		-[%rb],%rs	2	–	–	–	–	–	–	
		[%sp+imm7],%rs	2	–	–	–	–	–	–	
		[imm7],%rs	1	–	–	–	–	–	–	
	ld.ub	%rd,%rs	1	–	–	–	–	–	–	
		%rd,[%rb]	1–2*1	–	–	–	–	–	–	
		%rd,[%rb]+	2	–	–	–	–	–	–	
		%rd,[%rb]-	2	–	–	–	–	–	–	
		%rd,-[%rb]	2	–	–	–	–	–	–	
		%rd,[%sp+imm7]	2	–	–	–	–	–	–	
		%rd,[imm7]	1	–	–	–	–	–	–	
	ld	%rd,%rs	1	–	–	–	–	–	–	
		%rd,sign7	1	–	–	–	–	–	–	
		%rd,[%rb]	1–2*1	–	–	–	–	–	–	
		%rd,[%rb]+	2	–	–	–	–	–	–	
		%rd,[%rb]-	2	–	–	–	–	–	–	
		%rd,-[%rb]	2	–	–	–	–	–	–	
		%rd,[%sp+imm7]	2	–	–	–	–	–	–	
		%rd,[imm7]	1	–	–	–	–	–	–	
		[%rb],%rs	1–2*1	–	–	–	–	–	–	
		[%rb]+,%rs	2	–	–	–	–	–	–	
		[%rb]-,%rs	2	–	–	–	–	–	–	
		-[%rb],%rs	2	–	–	–	–	–	–	
		[%sp+imm7],%rs	2	–	–	–	–	–	–	
		[imm7],%rs	1	–	–	–	–	–	–	
	ld.a	%rd,%rs	1	–	–	–	–	–	–	
		%rd,imm7	1	–	–	–	–	–	–	
		%rd,[%rb]	1–2*1	–	–	–	–	–	–	
		%rd,[%rb]+	2	–	–	–	–	–	–	
		%rd,[%rb]-	2	–	–	–	–	–	–	
		%rd,-[%rb]	2	–	–	–	–	–	–	
		%rd,[%sp+imm7]	2	–	–	–	–	–	–	
		%rd,[imm7]	1	–	–	–	–	–	–	
		[%rb],%rs	1–2*1	–	–	–	–	–	–	
		[%rb]+,%rs	2	–	–	–	–	–	–	
		[%rb]-,%rs	2	–	–	–	–	–	–	
		-[%rb],%rs	2	–	–	–	–	–	–	
		[%sp+imm7],%rs	2	–	–	–	–	–	–	
		[imm7],%rs	1	–	–	–	–	–	–	
		%rd,%sp	1	–	–	–	–	–	–	
		%rd,%pc	1	–	–	–	–	–	–	
		%rd,[%sp]	1–2*1	–	–	–	–	–	–	
		%rd,[%sp]+	2	–	–	–	–	–	–	
		%rd,[%sp]-	2	–	–	–	–	–	–	
		%rd,-[%sp]	2	–	–	–	–	–	–	

Classification	Mnemonic		Cycle	Flag						Remark
				IL	IE	C	V	Z	N	
Data transfer	ld.a	[%sp],%rs	1–2*1	–	–	–	–	–	–	*1: 1 cycle when ext is not used 2 cycles when ext is used
		[%sp]+,%rs	2	–	–	–	–	–	–	
		[%sp]-,%rs	2	–	–	–	–	–	–	
		-[%sp],%rs	2	–	–	–	–	–	–	
		%sp,%rs	1	–	–	–	–	–	–	
		%sp,*imm7*	1	–	–	–	–	–	–	
Integer arithmetic operation	add	%rd,%rs	1	–	–	↔	↔	↔	↔	
	add/c	%rd,%rs	1	–	–	–	↔	↔	↔	
	add/nc	%rd,%rs	1	–	–	–	↔	↔	↔	
	add	%rd,*imm7*	1	–	–	↔	↔	↔	↔	
	add.a	%rd,%rs	1	–	–	–	–	–	–	
	add.a/c	%rd,%rs	1	–	–	–	–	–	–	
	add.a/nc	%rd,%rs	1	–	–	–	–	–	–	
	add.a	%sp,%rs	1	–	–	–	–	–	–	
		%rd,*imm7*	1	–	–	–	–	–	–	
		%sp,*imm7*	1	–	–	–	–	–	–	
	adc	%rd,%rs	1	–	–	↔	↔	↔	↔	
	adc/c	%rd,%rs	1	–	–	–	↔	↔	↔	
	adc/nc	%rd,%rs	1	–	–	–	↔	↔	↔	
	adc	%rd,*imm7*	1	–	–	↔	↔	↔	↔	
	sub	%rd,%rs	1	–	–	↔	↔	↔	↔	
	sub/c	%rd,%rs	1	–	–	–	↔	↔	↔	
	sub/nc	%rd,%rs	1	–	–	–	↔	↔	↔	
	sub	%rd,*imm7*	1	–	–	↔	↔	↔	↔	
	sub.a	%rd,%rs	1	–	–	–	–	–	–	
	sub.a/c	%rd,%rs	1	–	–	–	–	–	–	
	sub.a/nc	%rd,%rs	1	–	–	–	–	–	–	
	sub.a	%sp,%rs	1	–	–	–	–	–	–	
		%rd,*imm7*	1	–	–	–	–	–	–	
		%sp,*imm7*	1	–	–	–	–	–	–	
	sbc	%rd,%rs	1	–	–	↔	↔	↔	↔	
	sbc/c	%rd,%rs	1	–	–	–	↔	↔	↔	
	sbc/nc	%rd,%rs	1	–	–	–	↔	↔	↔	
	sbc	%rd,*imm7*	1	–	–	↔	↔	↔	↔	
	cmp	%rd,%rs	1	–	–	↔	↔	↔	↔	
	cmp/c	%rd,%rs	1	–	–	–	↔	↔	↔	
	cmp/nc	%rd,%rs	1	–	–	–	↔	↔	↔	
	cmp	%rd,*sign7*	1	–	–	↔	↔	↔	↔	
	cmp.a	%rd,%rs	1	–	–	↔	–	↔	–	
	cmp.a/c	%rd,%rs	1	–	–	–	–	↔	–	
	cmp.a/nc	%rd,%rs	1	–	–	–	–	↔	–	
	cmp.a	%rd,*imm7*	1	–	–	↔	–	↔	–	
	cmc	%rd,%rs	1	–	–	↔	↔	↔	↔	
	cmc/c	%rd,%rs	1	–	–	–	↔	↔	↔	
	cmc/nc	%rd,%rs	1	–	–	–	↔	↔	↔	
	cmc	%rd,*sign7*	1	–	–	↔	↔	↔	↔	
Logical operation	and	%rd,%rs	1	–	–	–	0	↔	↔	
	and/c	%rd,%rs	1	–	–	–	0	↔	↔	
	and/nc	%rd,%rs	1	–	–	–	0	↔	↔	
	and	%rd,*sign7*	1	–	–	–	0	↔	↔	
	or	%rd,%rs	1	–	–	–	0	↔	↔	
	or/c	%rd,%rs	1	–	–	–	0	↔	↔	
	or/nc	%rd,%rs	1	–	–	–	0	↔	↔	
	or	%rd,*sign7*	1	–	–	–	0	↔	↔	
	xor	%rd,%rs	1	–	–	–	0	↔	↔	
	xor/c	%rd,%rs	1	–	–	–	0	↔	↔	
	xor/nc	%rd,%rs	1	–	–	–	0	↔	↔	
	xor	%rd,*sign7*	1	–	–	–	0	↔	↔	
	not	%rd,%rs	1	–	–	–	0	↔	↔	
	not/c	%rd,%rs	1	–	–	–	0	↔	↔	
	not/nc	%rd,%rs	1	–	–	–	0	↔	↔	
	not	%rd,*sign7*	1	–	–	–	0	↔	↔	

Classification	Mnemonic		Cycle	Flag						Remark
				IL	IE	C	V	Z	N	
Shift and swap	sr	%rd,%rs	1	–	–	↔	–	↔	↔	
		%rd,imm7	1	–	–	↔	–	↔	↔	
	sa	%rd,%rs	1	–	–	↔	–	↔	↔	
		%rd,imm7	1	–	–	↔	–	↔	↔	
	sl	%rd,%rs	1	–	–	↔	–	↔	↔	
		%rd,imm7	1	–	–	↔	–	↔	↔	
	swap	%rd,%rs	1	–	–	–	–	–	–	
Immediate extension	ext	imm13	1	–	–	–	–	–	–	
Conversion	cv.ab	%rd,%rs	1	–	–	–	–	–	–	
	cv.as	%rd,%rs	1	–	–	–	–	–	–	
	cv.al	%rd,%rs	1	–	–	–	–	–	–	
	cv.la	%rd,%rs	1	–	–	–	–	–	–	
	cv.ls	%rd,%rs	1	–	–	–	–	–	–	
Branch	jpr jpr.d	sign10 %rb	3 2(.d)*3	–	–	–	–	–	–	*2: 2 cycles when not jumped 3 cycles when jumped *3: When a 1-cycle delayed slot instruction follows Same values as one without (.d) when a 2-cycle delayed slot instruction follows
	jpa ipa.d	imm7 %rb	3 2(.d)*3	–	–	–	–	–	–	
	jrgt jrgt.d	sign7	2–3*2 2(.d)*3	–	–	–	–	–	–	
	jrge jrge.d	sign7	2–3*2 2(.d)*3	–	–	–	–	–	–	
	jrlt jrlt.d	sign7	2–3*2 2(.d)*3	–	–	–	–	–	–	
	jrle jrle.d	sign7	2–3*2 2(.d)*3	–	–	–	–	–	–	
	jrugt jrugt.d	sign7	2–3*2 2(.d)*3	–	–	–	–	–	–	
	jruge jruge.d	sign7	2–3*2 2(.d)*3	–	–	–	–	–	–	
	jrult jrult.d	sign7	2–3*2 2(.d)*3	–	–	–	–	–	–	
	jrule jrule.d	sign7	2–3*2 2(.d)*3	–	–	–	–	–	–	
	jreq jreq.d	sign7	2–3*2 2(.d)*3	–	–	–	–	–	–	
	jrne jrne.d	sign7	2–3*2 2(.d)*3	–	–	–	–	–	–	
	call call.d	sign10 %rb	4 3(.d)*3	–	–	–	–	–	–	
	calla calla.d	imm7 %rb	4 3(.d)*3	–	–	–	–	–	–	
	ret ret.d		3 2(.d)*3	–	–	–	–	–	–	
	int	imm5	3	–	0	–	–	–	–	
	intl	imm5,imm3	3	↔	0	–	–	–	–	
	reti reti.d		3 2(.d)*3	↔	↔	↔	↔	↔	↔	
	brk		4	–	0	–	–	–	–	
	retd		4	↔	↔	↔	↔	↔	↔	
System control	nop		1	–	–	–	–	–	–	
	halt		6	–	–	–	–	–	–	
	slp		6	–	–	–	–	–	–	
	ei		1	–	1	–	–	–	–	
	di		1	–	0	–	–	–	–	
Coprocessor control	ld.cw	%rd,%rs %rd,imm7	1	–	–	–	–	–	–	
	ld.ca	%rd,%rs %rd,imm7	1	–	–	↔	↔	↔	↔	
	ld.cf	%rd,%rs %rd,imm7	1	–	–	↔	↔	↔	↔	

6.3 Interrupts

When an interrupt occurs during program execution, the processor enters an interrupt handling state. The interrupt handling state is a process by which the processor branches to the corresponding user's service routine for the interrupt that occurred. The processor returns after branching and starts executing the program from where it left off.

6.3.1 Priority of Interrupts

The interrupts supported by the S1C17 Core, their vector addresses and the priority of these interrupts are listed in the table below.

Table 6.3.1.1 Vector Address and Priority of Interrupts

Interrupt	Vector address (Hex)	Priority
Reset	TTBR + 0x00	High
Address misaligned interrupt	TTBR + 0x04	↑
Debug interrupt	(0xfffc00)	
NMI	TTBR + 0x08	
Software interrupt	TTBR + 0x00 to TTBR + 0x7c	↓
Maskable external interrupt	TTBR + 0x00 to TTBR + 0x7c	Low

When two or more interrupts occur simultaneously, they are processed in order of priority beginning with the one that has the highest priority.

When an interrupt occurs, the processor disables interrupts that would occur thereafter and performs interrupt handling. To support multiple interrupts (or another interrupt from within an interrupt), set the IE flag in the PSR to 1 in the interrupt handler routine to enable interrupts during interrupt handling. Basically, even when multiple interrupts are enabled, interrupts whose priorities are below the one set by the IL[2:0] bits in the PSR are not accepted.

The debug interrupt does not use the vector table and the stack. The PC and PSR are saved in a specific area along with R0.
The table below shows the addresses that are referenced when a debug interrupt occurs.

Table 6.3.1.2 Debug Interrupt Handler Start Address and Register Save Area

Address	Content
0xfffc00	Debug interrupt handler start address
DBRAM set value + 0x00	PC and PSR save area
DBRAM set value + 0x04	R0 save area

(DBRAM: See Section 4.2.3)

During debug interrupt handling, neither other interrupts nor multiple debug interrupts are accepted. They are kept pending until the debug interrupt handling currently underway finishes.

6.3.2 Vector Table

Vector table in the S1C17 Core

The table below lists the interrupts for which the vector table is referenced during interrupt handling.

Table 6.3.2.1 Vector List

Vector No. Software interrupt No.	Interrupt	Vector address
0 (0x00)	Reset	TTBR + 0x00
1 (0x01)	Address misaligned interrupt	TTBR + 0x04
2 (0x02)	NMI	TTBR + 0x08
3 (0x03)	Maskable external interrupt 3	TTBR + 0x0c
:	:	:
31 (0x1f)	Maskable external interrupt 31	TTBR + 0x7c

The vector address is one that contains a vector (or the jump address) for the user's interrupt handler routine that is provided for each interrupt and is executed when the relevant interrupt occurs. Because an address value is stored, each vector address is located at a 16-bit boundary. The memory area in which these vectors are stored is referred to as the "vector table." The "TTBR" in the Vector Address column represents the base (start) address of the vector table. For the TTBR value, refer to the Technical Manual of each model. The set value can be read from TTBR (vector table base register) located at address 0xffff80.

6.3.3 Interrupt Handling

When an interrupt occurs, the processor starts interrupt handling. (This interrupt handling does not apply for reset and debug interrupts.)

The interrupt handling performed by the processor is outlined below.

(1) Suspends the instructions currently being executed.
An interrupt is generated synchronously with the rising edge of the system clock at the end of the cycle of the currently executed instruction.

(2) Saves the contents of the PC and PSR to the stack (SP), in that order.

(3) Clears the IE (interrupt enable) bit in the PSR to disable maskable interrupts that would occur thereafter. If the generated interrupt is a maskable interrupt, the IL (interrupt level) in the PSR is rewritten to that of the generated interrupt.

(4) Reads the vector for the generated interrupt from the vector table, and sets it in the PC. The processor thereby branches to the user's interrupt handler routine.

After branching to the user's interrupt handler routine, when the `reti` instruction is executed at the end of interrupt handling, the saved data is restored from the stack in order of the PC and PSR, and the processing returns to the suspended instructions.

6.3.4 Reset

The processor is reset by applying a low-level pulse to its rst_n pin. All the registers are thereby cleared to 0.
The processor starts operating at the rising edge of the reset pulse to perform a reset sequence. In this reset sequence, the reset vector is read out from the top of the vector table and set in the PC. The processor thereby branches to the user's initialization routine, in which it starts executing the program. The reset sequence has priority over all other processing.

6.3.5 Address Misaligned Interrupt

The load instructions that access memory or I/O areas are characteristic in that the data size to be transferred is predetermined for each instruction used, and that the accessed addresses must be aligned with the respective data-size boundaries.

Instruction	Transfer data size	Address
`ld.b/ld.ub`	Byte (8 bits)	Byte boundary (applies to all addresses)
`ld`	16 bits	16-bit boundary (least significant address bit = 0)
`ld.a`	32 bits	32-bit boundary (two least significant address bits = 00)

If the specified address in a load instruction does not satisfy this condition, the processor assumes an address misaligned interrupt and performs interrupt handling. Even in this case the load instruction is executed as the least significant bit or the two low-order bits of the address set to 0. The PC value saved to the stack in interrupt handling is the address of the load instruction that caused the interrupt.

This interrupt does not occur in the program branch instructions as the least significant bit of the PC is always fixed to 0. The same applies to the vector for interrupt handling.

6.3.6 NMI

An NMI is generated when the nmi_n input on the processor is asserted low. When an NMI occurs, the processor performs interrupt handling after it has finished executing the instruction currently underway.

6.3.7 Maskable External Interrupts

The S1C17 Core can accept up to 32 types of maskable external interrupts (however, the first three interrupt causes use the save vector address as the reset interrupt, address misaligned interrupt, and NMI). It is only when the IE (interrupt enable) flag in the PSR is set that the processor accepts a maskable external interrupt. Furthermore, their acceptable interrupt levels are limited by the IL (interrupt level) field in the PSR. The interrupt levels (0–7) in the IL field dictate the interrupt levels that can be accepted by the processor, and only interrupts with priority levels higher than that are accepted. Interrupts with the same interrupt level as IL cannot be accepted.
The IE flag can be set in the software. When an interrupt occurs, the IE flag is cleared to 0 (interrupts disabled) after the PSR is saved to the stack, and the maskable interrupts remain disabled until the IE flag is set in the handler routine or the handler routine is terminated by the `reti` instruction that restores the PSR from the stack. The IL field is set to the priority level of the interrupt that occurred.
Multiple interrupts or the ability to accept another interrupt during interrupt handling if its priority is higher than that of the currently serviced interrupt can easily be realized by setting the IE flag in the interrupt handler routine.
When the processor is reset, the PSR is initialized to 0 and the maskable interrupts are therefore disabled, and the interrupt level is set to 0 (interrupts with priority levels 1–7 enabled).

The following describes how the maskable interrupts are accepted and processed by the processor.

(1) Suspends the instructions currently being executed.
The interrupt is accepted synchronously with the rising edge of the system clock at the end of the cycle of the currently executed instruction.

(2) Saves the contents of the PC (current value) and PSR to the stack (SP), in that order.

(3) Clears the IE flag in the PSR and copy the priority level of the accepted interrupt to the IL field.

(4) Reads the vector for the interrupt from the vector address in the vector table, and sets it in the PC. The processor then branches to the interrupt handler routine.

In the interrupt handler routine, the `reti` instruction should be executed at the end of processing. In the `reti` instruction, the saved data is restored from the stack in order of the PC and PSR, and the processing returns to the suspended instructions.

6.3.8 Software Interrupts

The S1C17 Core provides the `int` *imm5* and `intl` *imm5,imm3* instructions allowing the software to generate any interrupts. The operand *imm5* specifies a vector number (0–31) in the vector table. In addition to this, the `intl` instruction has the operand *imm3* to specify an interrupt level (0–7) to be set to the IL field in the PSR.
The processor performs the same interrupt handling as that of a hardware interrupt.

6.3.9 Interrupt Masked Period

Address misaligned interrupts, NMIs, debug interrupts, and external maskable interrupts are masked between the specific instructions listed below and cannot be generated during that period (pending state). When the processor exits the masked period, the pending interrupt can be accepted.

(1) Between the `ext` instruction and the next instruction
(2) Between a delayed branch (`.d`) instruction and the delayed slot instruction that follows
(3) Between the `retd` instruction and the next instruction (located at the return address)
(4) Between the `reti` or `reti.d` *1 instruction and the next instruction (located at the return address) *2
(5) Between the `int`, `ei`, `di`, `slp`, or `halt` instruction and the next instruction *2
(6) Between a conditional jump (`jr*`) instruction and the next instruction when the condition has not been met *2

*1 An interrupt that occurs when the `reti.d` instruction is being executed will be accepted after the delayed slot instruction that follows and the next instruction (located at the return address) are executed.

|	
`reti.d`	
Delayed slot instruction	Interrupt masked state
|	
Instruction at return address	Interrupt masked state still continues, so the next instruction will be executed before interrupts can be generated.
Next instruction	Interrupt mask is released.

*2 The debug interrupt may occur even in the conditions (4) to (6).

6.4 Power-Down Mode

The S1C17 Core supports two power-down modes: HALT and SLEEP modes.

HALT mode

Program execution is halted at the same time that the S1C17 Core executes the `halt` instruction, and the processor enters HALT mode.
HALT mode commonly turns off only the S1C17 Core operation, note, however that modules to be turned off depend on the implementation of the clock control circuit outside the core. Refer to the technical manual of each model for details.

SLEEP mode

Program execution is halted at the same time the S1C17 Core executes the `slp` instruction, and the processor enters SLEEP mode.
SLEEP mode commonly turns off the S1C17 Core and on-chip peripheral circuit operations, thereby it significantly reduces the current consumption in comparison to HALT mode. However, modules to be turned off depend on the implementation of the clock control circuit outside the core. Refer to the technical manual of each model for details.

Canceling HALT or SLEEP mode

Initial reset is one cause that can bring the processor out of HALT or SLEEP mode. Other causes depend on the implementation of the clock control circuit outside the S1C17 Core.
Initial reset, maskable external interrupts, NMI, and debug interrupts are commonly used for canceling HALT and SLEEP modes.

The interrupt enable/disable status set in the processor does not affect the cancellation of HALT or SLEEP mode even if an interrupt signal is used as the cancellation. In other words, interrupt signals are able to cancel HALT and SLEEP modes even if the IE flag in PSR or the interrupt enable bits in the interrupt controller (depending on the implementation) are set to disable interrupts.
When the processor is taken out of HALT or SLEEP mode using an interrupt that has been enabled (by the interrupt controller and IE flag), the corresponding interrupt handler routine is executed after executing the instruction next to the `halt` or `slp` instruction.
When the interrupt has been disabled, the processor restarts the program from the instruction next to `halt` or `slp` after the processor is taken out of HALT or SLEEP mode.

6.5 Debug Circuit

The S1C17 Core has a debug circuit to assist in software development by the user.

6.5.1 Debugging Functions

The debug circuit provides the following functions:

• Instruction break
  A debug interrupt is generated before the set instruction address is executed. An instruction break can be set at two addresses.

• Single step
  A debug interrupt is generated every instruction executed.

• Forcible break
  A debug interrupt is generated by an external input signal.

• Software break
  A debug interrupt is generated when the `brk` instruction is executed.

When a debug interrupt occurs, the processor performs the following processing:

(1) Suspends the instructions currently being executed.

(2) Saves the contents of the PC and PSR, and R0, in that order, to the addresses specified below.
PC/PSR → DBRAM + 0x0
R0 → DBRAM + 0x4 (DBRAM: Start address of the work area for debugging in the user RAM)

(3) Loads address 0xfffc00 to PC and branches to the debug interrupt handler routine.

In the interrupt handler routine, the `retd` instruction should be executed at the end of processing to return to the suspended instructions. When returning from the interrupt by the `retd` instruction, the processor restores the saved data in order of the R0 and the PC and PSR.
Neither hardware interrupts nor NMI interrupts are accepted during a debug interrupt.

6.5.2 Resource Requirements and Debugging Tools

The on-chip debug function requires a 64-byte work area. For the work area for debugging, refer to the Technical Manual of each model.

Debugging is performed by connecting a serial ICE to the debug pins of the S1C17 Core and entering debug commands from the debugger being run on a personal computer. The tools listed below are required for debugging.
• S1C17 Family Serial ICE (S5U1C17001H)
• S1C17 Family C Compiler Package

6.5.3 Registers for Debugging

The reserved core I/O area contains the debug registers described below.

0xFFFF90: Debug RAM Base Register (DBRAM)

Register name	Address	Bit	Name	Function	Setting	Init.	R/W	Remarks
Debug RAM base register	**FFFF90 (L)**	D31–24	–	Unused (fixed at 0)	0x0	0x0	R	
		D23 | D0	**DBRAM23 | DBRAM0**	Debug RAM base address DBRAM[5:0] is fixed at 0x0.	0x0–0xFFFDC0 (64 byte units)	*	R	Initial value is set in the C17 RTL-define DBRAM_BASE.

D[23:0] DBRAM[23:0]: Debug RAM Base Address Bits

This is a read-only register that contains the start address of a work area (64 bytes) for debugging.

0xFFFFA0: Debug Control Register (DCR)

Register name	Address	Bit	Name	Function	Setting				Init.	R/W	Remarks
Debug control register	**FFFFA0 (B)**	D7–5	–	Reserved	–				–	–	0 when being read.
		D4	**DR**	Debug request flag	1	Occurred	0	Not occurred	0	R/W	Reset by writing 1.
		D3	**IBE1**	Instruction break #1 enable	1	Enable	0	Disable	0	R/W	
		D2	**IBE0**	Instruction break #0 enable	1	Enable	0	Disable	0	R/W	
		D1	**SE**	Single step enable	1	Enable	0	Disable	0	R/W	
		D0	**DM**	Debug mode	1	Debug mode	0	User mode	0	R	

D[7:5] Reserved

D4 DR: Debug Request Flag

Indicates whether an external debug request has occurred or not.

1 (R): Occurred
0 (R): Not occurred (default)
1 (W): Flag is reset
0 (W): Has no effect

This flag is cleared (reset to 0) by writing 1. The flag must be cleared before the debug handler routine has been terminated by executing the `retd` instruction.

D3 IBE1: Instruction Break #1 Enable Bit

Enables/disables instruction break #1.

1 (R/W): Enable
0 (R/W): Disable (default)

When this bit is set to 1, instruction fetch addresses will be compared with the value set in the Instruction Break Address Register 1 (0xffffb4), and an instruction break will occur if they are matched. Setting this bit to 0 disables the comparison.

D2 IBE0: Instruction Break #0 Enable Bit

Enables/disables instruction break #0.

1 (R/W): Enable
0 (R/W): Disable (default)

When this bit is set to 1, instruction fetch addresses will be compared with the value set in the Instruction Break Address Register 0 (0xffffb0), and an instruction break will occur if they are matched. Setting this bit to 0 disables the comparison.

D1 SE: Single Step Enable Bit

Enables/disables single-step execution.

1 (R/W): Enable
0 (R/W): Disable (default)

D0 DM: Debug Mode Bit

Indicates the current operation mode of the processor (debug mode or user mode).

1 (R): Debug mode
0 (R): User mode (default)

0xFFFFB0: Instruction Break Address Register 0 (IBAR0)

Register name	Address	Bit	Name	Function	Setting	Init.	R/W	Remarks
Instruction break address register 0	**FFFFB0 (L)**	D31–24	–	Unused (fixed at 0)	0x0	0x0	R	
		D23 | D0	**IBAR023** | **IBAR00**	Instruction break address #0 IBAR00 is fixed at 0.	0x0–0xFFFDE	0x0	R/W	

D[23:0] IBAR0[23:0]: Instruction Break Address #0

This register is used to set instruction break address #0. (Default: 0x000000)

0xFFFFB4: Instruction Break Address Register 1 (IBAR1)

Register name	Address	Bit	Name	Function	Setting	Init.	R/W	Remarks
Instruction break address register 1	**FFFFB4 (L)**	D31–24	–	Unused (fixed at 0)	0x0	0x0	R	
		D23 | D0	**IBAR123** | **IBAR10**	Instruction break address #1 IBAR10 is fixed at 0.	0x0–0xFFFDE	0x0	R/W	

D[23:0] IBAR1[23:0]: Instruction Break Address #1

This register is used to set instruction break address #1. (Default: 0x000000)

0xFFFFC0: Serial Status Register for Debugging (SSR)

Register name	Address	Bit	Name	Function	Setting				Init.	R/W	Remarks
Serial status register for debugging	**FFFFC0 (B)**	D7–3	–	Reserved	–				–	–	0 when being read.
		D2	**RXDEN**	Receive disable	1	Disable	0	Enable	1	R/W	
		D1	**TDBE**	Transmit data buffer empty flag	1	Empty	0	Not empty	1	R	
		D0	**RDBF**	Receive data buffer full flag	1	Full	0	Not full	0	R	

D[7:3] Reserved

D2 RXDEN: Receive Disable Bit

Enables/disables receive operation in the serial interface for the on-chip debug monitor.

1 (R/W): Disable (default)
0 (R/W): Enable

D1 TDBE: Transmit Data Buffer Empty Flag

Indicates transmit buffer status in the serial interface for the on-chip debug monitor.

1 (R): Empty (default)
0 (R): Not empty

D0 RDBF: Receive Data Buffer Full Flag

Indicates receive buffer status in the serial interface for the on-chip debug monitor.

1 (R): Full
0 (R): Not full (default)

0xFFFFC2: Serial Transmit/Receive Data Register for Debugging (SDR)

Register name	Address	Bit	Name	Function	Setting	Init.	R/W	Remarks
Serial transmit/ receive data register for debugging	**FFFFC2 (B)**	D7 | | D0	**TXRXD7** | | **TXRXD0**	Transmit/receive data	0x0–0xFF	0x0	R/W	

D[7:0] TXRXD[7:0]: Transmit/Receive Data

This is the transmit/receive data register of the serial interface for the on-chip debug monitor used to set transmit data and to store received data. (Default: 0x00)

7 Details of Instructions

This section explains all the instructions in alphabetical order.

Symbols in the instruction reference

↔	Indicates that the bit is set (= 1) or reset (= 0) by instruction execution
1	Indicates that the bit is set (= 1) by instruction execution
0	Indicates that the bit is reset (= 0) by instruction execution

Registers/Register Data

`%rd`, *rd*:	A general-purpose register (R0–R7) used as the destination register or its contents
`%rs`, *rs*:	A general-purpose register (R0–R7) used as the source register or its contents
`%rb`, *rb*:	A general-purpose register (R0–R7) that has stored a base address to be accessed in the register indirect addressing mode or its contents
`%sp`, sp:	Stack pointer (SP) or its contents
`%pc`, pc:	Program counter (PC) or its contents

The register field (*rd*, *rs*) in the code contains a general-purpose register number.
R0 = 0b000, R1 = 0b001 . . . R7 = 0b111

Memory/Addresses/Memory Data

`[%rb]`, `[%sp]`:	Specification for register indirect addressing
`[%rb]+`, `[%sp]+`:	Specification for register indirect addressing with post-increment
`[%rb]-`, `[%sp]-`:	Specification for register indirect addressing with post-decrement
`-[%rb]`, `-[%sp]`:	Specification for register indirect addressing with pre-decrement
`[%sp+immX]`:	Specification for register indirect addressing with a displacement
`[imm7]`:	Specification for a memory address with an immediate data
`B[XXX]`:	An address specified with *XXX*, or the byte data stored in the address
`W[XXX]`:	A 16-bit address specified with *XXX*, or the word data stored in the address
`A[XXX]`:	A 32-bit address specified with *XXX*, or the 24-bit or 32-bit data stored in the address

Immediate

`immX`:	A *X*-bit unsigned immediate data. The *X* contains a number representing the bit length of the immediate.
`signX`:	A *X*-bit signed immediate data. The *X* contains a number representing the bit length of the immediate. Furthermore, the most significant bit is handled as the sign bit.

Bit Field

(*X*):	Bit *X* of data
(*X*:*Y*):	A bit field from bit *X* to bit *Y*
{*X*,*Y*...}:	Indicates a bit (data) configuration.

Code

rd, *rs*, *rb*:	Register number (R0 = 0 ... R7 = 7)
d:	Delayed bit (0: Standard branch instruction, 1: Delayed branch instruction)

Functions

←:	Indicates that the right item is loaded or set to the left item.
+:	Addition
-:	Subtraction
&:	AND
|:	OR
^:	XOR
!:	NOT

Flags

IL:	Interrupt level
IE:	Interrupt enable flag
C:	Carry flag
V:	Overflow flag
Z:	Zero flag
N:	Negative flag
–:	Not changed
↔:	Set (1) or reset (0)
1:	Set (1)
0:	Reset (0)

adc %rd, %rs
adc/c %rd, %rs
adc/nc %rd, %rs

Function

16-bit addition with carry

Standard) *rd*(15:0) ← *rd*(15:0) + *rs*(15:0) + C, *rd*(23:16) ← 0

Extension 1) *rd*(15:0) ← *rs*(15:0) + *imm13*(zero extended) + C, *rd*(23:16) ← 0

Extension 2) *rd*(15:0) ← *rs*(15:0) + *imm16* + C, *rd*(23:16) ← 0

Code

15	14	13	12	11	10	9–8 (7)	6	5	4	3	2–0	
0	0	1	1	1	0	rd	1	0	0	1	rs	adc
0	0	1	1	1	0	rd	0	0	0	1	rs	adc/c
0	0	1	1	1	0	rd	0	1	0	1	rs	adc/nc

Flag

IL	IE	C	V	Z	N	
–	–	↔	↔	↔	↔	adc
–	–	–	↔	↔	↔	adc/c, adc/nc

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
adc   %rd,%rs          ; rd ← rd + rs + C
```

The content of the *rs* register and C (carry) flag are added to the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(2) Extension 1

```
ext   imm13
adc   %rd,%rs          ; rd ← rs + imm13 + C
```

The 13-bit immediate *imm13* and C (carry) flag are added to the content of the *rs* register after being zero-extended, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0. The content of the *rs* register is not altered.

(3) Extension 2

```
ext   imm3             ; imm3(2:0)  = imm16(15:13)
ext   imm13            ; = imm16(12:0)
adc   %rd,%rs          ; rd ← rs + imm16 + C
```

The 16-bit immediate *imm16* and C (carry) flag are added to the content of the *rs* register, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0. The content of the *rs* register is not altered.

(4) Conditional execution

The `/c` or `/nc` suffix on the opcode specifies conditional execution.

`adc/c` Executed as `adc` when the C flag is 1 or executed as `nop` when the flag is 0

`adc/nc` Executed as `adc` when the C flag is 0 or executed as `nop` when the flag is 1

In this case, the `ext` instruction can be used to extend the operand.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
adc   %r0,%r1          ; r0 = r0 + r1 + C
```

(2) Addition of 32-bit data, data 1 = {r2, r1}, data 2 = {r4, r3}, result = {r2, r1}

```
add   %r1,%r3          ; Addition of the low-order word
adc   %r2,%r4          ; Addition of the high-order word
```

adc %rd, imm7

Function

16-bit addition with carry

Standard) *rd*(15:0) ← *rd*(15:0) + *imm7*(zero extended) + C, *rd*(23:16) ← 0

Extension 1) *rd*(15:0) ← *rd*(15:0) + *imm16* + C, *rd*(23:16) ← 0

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	0	1		*rd*					*imm7*			

Flag

IL	IE	C	V	Z	N
–	–	↔	↔	↔	↔

Mode

Src: Immediate data (unsigned)

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
adc   %rd,imm7          ; rd ← rd + imm7 + C
```

The 7-bit immediate *imm7* and C (carry) flag are added to the *rd* register after being zero-extended. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(2) Extension 1

```
ext   imm9              ; imm9(8:0) = imm16(15:7)
adc   %rd,imm7          ; rd ← rd + imm16 + C, imm7 = imm16(6:0)
```

The 16-bit immediate *imm16* and C (carry) flag are added to the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(3) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
adc   %r0,0x7f          ; r0 = r0 + 0x7f + C
```

(2)
```
ext   0x1ff
adc   %r1,0x7f          ; r1 = r1 + 0xffff + C
```

add %rd, %rs
add/c %rd, %rs
add/nc %rd, %rs

Function

16-bit addition

Standard) *rd*(15:0) ← *rd*(15:0) + *rs*(15:0), *rd*(23:16) ← 0

Extension 1) *rd*(15:0) ← *rs*(15:0) + *imm13*(zero extended), *rd*(23:16) ← 0

Extension 2) *rd*(15:0) ← *rs*(15:0) + *imm16*, *rd*(23:16) ← 0

Code

15	14	13	12	11	10	9–7	6	5	4	3	2–0	
0	0	1	1	1	0	rd	1	0	0	0	rs	add
0	0	1	1	1	0	rd	0	0	0	0	rs	add/c
0	0	1	1	1	0	rd	0	1	0	0	rs	add/nc

Flag

IL	IE	C	V	Z	N	
–	–	↔	↔	↔	↔	add
–	–	–	↔	↔	↔	add/c, add/nc

Mode

Src: Register direct %rs = %r0 to %r7

Dst: Register direct %rd = %r0 to %r7

CLK

One cycle

Description

(1) Standard

```
add   %rd,%rs         ; rd ← rd + rs
```

The content of the *rs* register is added to the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(2) Extension 1

```
ext   imm13
add   %rd,%rs         ; rd ← rs + imm13
```

The 13-bit immediate *imm13* is added to the content of the *rs* register after being zero-extended, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0. The content of the *rs* register is not altered.

(3) Extension 2

```
ext   imm3            ; imm3(2:0) = imm16(15:13)
ext   imm13           ; = imm16(12:0)
add   %rd,%rs         ; rd ← rs + imm16
```

The 16-bit immediate *imm16* is added to the content of the *rs* register, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0. The content of the *rs* register is not altered.

(4) Conditional execution

The /c or /nc suffix on the opcode specifies conditional execution.

add/c Executed as add when the C flag is 1 or executed as nop when the flag is 0

add/nc Executed as add when the C flag is 0 or executed as nop when the flag is 1

In this case, the ext instruction can be used to extend the operand.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the ext instruction cannot be performed.

Example

(1)
```
add   %r0,%r0         ; r0 = r0 + r0
```

(2)
```
ext   0x1
ext   0x1fff
add   %r1,%r2         ; r1 = r2 + 0x3fff
```

add %rd, imm7

Function

16-bit addition

Standard) *rd*(15:0) ← *rd*(15:0) + *imm7*(zero extended), *rd*(23:16) ← 0

Extension 1) *rd*(15:0) ← *rd*(15:0) + *imm16*, *rd*(23:16) ← 0

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	0	0		*r d*					*imm7*			

Flag

IL	IE	C	V	Z	N
–	–	↔	↔	↔	↔

Mode

Src: Immediate data (unsigned)

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
add   %rd,imm7        ; rd ← rd + imm7
```

The 7-bit immediate *imm7* is added to the *rd* register after being zero-extended. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(2) Extension 1

```
ext   imm9            ; imm9(8:0) = imm16(15:7)
add   %rd,imm7        ; rd ← rd + imm16, imm7 = imm16(6:0)
```

The 16-bit immediate *imm16* is added to the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(3) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
add   %r0,0x3f        ; r0 = r0 + 0x3f
```

(2)
```
ext   0x1ff
add   %r1,0x7f        ; r1 = r1 + 0xffff
```

add.a %rd, %rs
add.a/c %rd, %rs
add.a/nc %rd, %rs

Function

24-bit addition

Standard) *rd*(23:0) ← *rd*(23:0) + *rs*(23:0)
Extension 1) *rd*(23:0) ← *rs*(23:0) + *imm13*(zero extended)
Extension 2) *rd*(23:0) ← *rs*(23:0) + *imm24*

Code

15	14	13	12	11	10	9–7	6	5	4	3	2–0	
0	0	1	1	0	0	rd	1	0	0	0	rs	`add.a`
0	0	1	1	0	0	rd	0	0	0	0	rs	`add.a/c`
0	0	1	1	0	0	rd	0	1	0	0	rs	`add.a/nc`

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
add.a   %rd,%rs        ; rd ← rd + rs
```

The content of the *rs* register is added to the *rd* register.

(2) Extension 1

```
ext     imm13
add.a   %rd,%rs        ; rd ← rs + imm13
```

The 13-bit immediate *imm13* is added to the content of the *rs* register after being zero-extended, and the result is loaded into the *rd* register. The content of the *rs* register is not altered.

(3) Extension 2

```
ext     imm11          ; imm11(10:0) = imm24(23:13)
ext     imm13          ; = imm24(12:0)
add.a   %rd,%rs        ; rd ← rs + imm24
```

The 24-bit immediate *imm24* is added to the content of the *rs* register, and the result is loaded into the *rd* register. The content of the *rs* register is not altered.

(4) Conditional execution

The `/c` or `/nc` suffix on the opcode specifies conditional execution.

`add.a/c` Executed as `add.a` when the C flag is 1 or executed as `nop` when the flag is 0
`add.a/nc` Executed as `add.a` when the C flag is 0 or executed as `nop` when the flag is 1

In this case, the `ext` instruction can be used to extend the operand.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
add.a   %r0,%r0        ; r0 = r0 + r0
```

(2)
```
ext     0x7ff
ext     0x1fff
add.a   %r1,%r2        ; r1 = r2 + 0xffffff
```

add.a %rd, imm7

Function

24-bit addition

Standard) *rd*(23:0) ← *rd*(23:0) + *imm7*(zero extended)
Extension 1) *rd*(23:0) ← *rd*(23:0) + *imm20*(zero extended)
Extension 2) *rd*(23:0) ← *rd*(23:0) + *imm24*

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	0	*rd*			*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Immediate data (unsigned)
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
add.a   %rd,imm7       ; rd ← rd + imm7
```

The 7-bit immediate *imm7* is added to the *rd* register after being zero-extended.

(2) Extension 1

```
ext     imm13          ; = imm20(19:7)
add.a   %rd,imm7       ; rd ← rd + imm20, imm7 = imm20(6:0)
```

The 20-bit immediate *imm20* is added to the *rd* register after being zero-extended.

(3) Extension 2

```
ext     imm4           ; imm4(3:0) = imm24(23:20)
ext     imm13          ; = imm24(19:7)
add.a   %rd,imm7       ; rd ← rd + imm24, imm7 = imm24(6:0)
```

The 24-bit immediate *imm24* is added to the *rd* register.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
add.a   %r0,0x7f       ; r0 = r0 + 0x7f
```

(2)
```
ext     0xf
ext     0x1fff
add.a   %r1,0x7f       ; r1 = r1 + 0xffffff
```

add.a %sp, %rs

Function

24-bit addition

Standard) sp(23:0) ← sp(23:0) + *rs*(23:0)

Extension 1) sp(23:0) ← *rs*(23:0) + *imm13*(zero extended)

Extension 2) sp(23:0) ← *rs*(23:0) + *imm24*

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	1	0	0	0	0	0	0	0	0	1	*rs*		

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register direct `%sp`

CLK

One cycle

Description

(1) Standard

```
add.a   %sp,%rs        ; sp ← sp + rs
```

The content of the *rs* register is added to the stack pointer SP.

(2) Extension 1

```
ext     imm13
add.a   %sp,%rs        ; sp ← rs + imm13
```

The 13-bit immediate *imm13* is added to the content of the *rs* register after being zero-extended, and the result is loaded into the stack pointer SP. The content of the *rs* register is not altered.

(3) Extension 2

```
ext     imm11          ; imm11(10:0)  =  imm24(23:13)
ext     imm13          ; = imm24(12:0)
add.a   %sp,%rs        ; sp ← rs + imm24
```

The 24-bit immediate *imm24* is added to the content of the *rs* register, and the result is loaded into the stack pointer SP. The content of the *rs* register is not altered.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
add.a   %sp,%r0        ; sp = sp + r0
```

(2)
```
ext     0x1
ext     0x1ffc
add.a   %sp,%r2        ; sp = r2 + 0x3ffc
```

Caution

The 2 low-order bits of the addition results are always loaded to the SP as 0.

add.a %sp, *imm7*

Function

24-bit addition

Standard) sp(23:0) ← sp(23:0) + *imm7*(zero extended)
Extension 1) sp(23:0) ← sp(23:0) + *imm20*(zero extended)
Extension 2) sp(23:0) ← sp(23:0) + *imm24*

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	0	1	0	0	0	*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Immediate data (unsigned)
Dst: Register direct `%sp`

CLK

One cycle

Description

(1) Standard

```
add.a   %sp,imm7      ; sp ← sp + imm7
```

The 7-bit immediate *imm7* is added to the stack pointer SP after being zero-extended.

(2) Extension 1

```
ext     imm13         ; = imm20(19:7)
add.a   %sp,imm7      ; sp ← sp + imm20, imm7 = imm20(6:0)
```

The 20-bit immediate *imm20* is added to the stack pointer SP after being zero-extended.

(3) Extension 2

```
ext     imm4          ; imm4(3:0) = imm24(23:20)
ext     imm13         ; = imm24(19:7)
add.a   %sp,imm7      ; sp ← sp + imm24, imm7 = imm24(6:0)
```

The 24-bit immediate *imm24* is added to the stack pointer SP.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
add.a   %sp,0x7c      ; sp = sp + 0x7c
```

(2)
```
ext     0x1fff
add.a   %sp,0x7c      ; sp = sp + 0xffffc
```

Caution

The 2 low-order bits of the addition results are always loaded to the SP as 0.

and %rd, %rs
and/c %rd, %rs
and/nc %rd, %rs

Function

16-bit logical AND

Standard) *rd*(15:0) ← *rd*(15:0) & *rs*(15:0), *rd*(23:16) ← 0

Extension 1) *rd*(15:0) ← *rs*(15:0) & *imm13*(zero extended), *rd*(23:16) ← 0

Extension 2) *rd*(15:0) ← *rs*(15:0) & *imm16*, *rd*(23:16) ← 0

Code

15	14	13	12	11	10	9–7	6	5	4	3	2–0	
0	0	1	0	1	1	rd	1	0	0	0	rs	and
0	0	1	0	1	1	rd	0	0	0	0	rs	and/c
0	0	1	0	1	1	rd	0	1	0	0	rs	and/nc

Flag

IL	IE	C	V	Z	N
–	–	–	0	↔	↔

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
and  %rd,%rs          ; rd ← rd & rs
```

The content of the *rs* register and that of the *rd* register are logically AND'ed, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(2) Extension 1

```
ext  imm13
and  %rd,%rs          ; rd ← rs & imm13
```

The content of the *rs* register and the zero-extended 13-bit immediate *imm13* are logically AND'ed, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0. The content of the *rs* register is not altered.

(3) Extension 2

```
ext  imm3             ; imm3(2:0) = imm16(15:13)
ext  imm13            ; = imm16(12:0)
and  %rd,%rs          ; rd ← rs & imm16
```

The content of the *rs* register and the 16-bit immediate *imm16* are logically AND'ed, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0. The content of the *rs* register is not altered.

(4) Conditional execution

The `/c` or `/nc` suffix on the opcode specifies conditional execution.

`and/c` Executed as `and` when the C flag is 1 or executed as `nop` when the flag is 0

`and/nc` Executed as `and` when the C flag is 0 or executed as `nop` when the flag is 1

In this case, the `ext` instruction can be used to extend the operand.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
and  %r0,%r0          ; r0 = r0 & r0
```

(2)
```
ext  0x1
ext  0x1fff
and  %r1,%r2          ; r1 = r2 & 0x3fff
```

and %rd, sign7

Function

16-bit logical AND

Standard) *rd*(15:0) ← *rd*(15:0) & *sign7*(sign extended), *rd*(23:16) ← 0

Extension 1) *rd*(15:0) ← *rd*(15:0) & *sign16*, *rd*(23:16) ← 0

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	0	0	*rd*			*sign7*						

Flag

IL	IE	C	V	Z	N
–	–	–	0	↔	↔

Mode

Src: Immediate data (signed)

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
and   %rd,sign7        ; rd ← rd & sign7
```

The content of the *rd* register and the sign-extended 7-bit immediate *sign7* are logically AND'ed, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(2) Extension 1

```
ext   imm9             ; imm9(8:0) = sign16(15:7)
and   %rd,sign7        ; rd ← rd & sign16, sign7 = sign16(6:0)
```

The content of the *rd* register and the 16-bit immediate *sign16* are logically AND'ed, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(3) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
and   %r0,0x7e         ; r0 = r0 & 0xfffe
```

(2)
```
ext   0x3f
and   %r1,0x7f         ; r1 = r1 & 0x1fff
```

brk

Function

Debugging interrupt

Standard) A[DBRAM] ← {psr, pc + 2}, A[DBRAM + 0x4] ← r0, pc ← 0xfffc00

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	0	0	1	0	1	1	0	0	0	0	0

Flag

IL	IE	C	V	Z	N
–	0	–	–	–	–

Mode

–

CLK

Four cycles

Description

Calls a debugging handler routine.

The `brk` instruction stores the address (PC + 2) that follows this instruction, the contents of the PSR, and the contents of the R0 register into the work area for debugging (DBRAM), then sets the mini-monitor start address (0xfffc00) to the PC. Thus the program branches to the debug-handler routine. Furthermore the processor enters the debug mode.

The `retd` instruction must be used for return from the debug-handler routine.

This instruction is provided for debug firmware. Do not use it in the user program.

Example

```
brk                 ; Executes the debug-handler routine
```

call %*rb*
call.d %*rb*

Function

PC relative subroutine call

Standard) `call`: sp ← sp - 4, A[sp] ← pc + 2, pc ← pc + 2 + *rb*

`call.d`: sp ← sp - 4, A[sp] ← pc + 4, pc ← pc + 2 + *rb*

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	0	0	0	0	0	1	0	0	0	0	0	*rb*			`call`
0	0	0	0	0	0	0	1	1	0	0	0	0	*rb*			`call.d`

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Register direct `%rb` = `%r0` to `%r7`

CLK

`call` Four cycles

`call.d` Three cycles (when a one-cycle delayed slot instruction follows), Four cycles (other)

Description

(1) Standard

`call %rb`

Stores the address of the following instruction into the stack, then adds the contents of the *rb* register to the PC (PC + 2) for calling the subroutine that starts from the address set to the PC. The LSB of the *rb* register is invalid and is always handled as 0. When the `ret` instruction is executed in the subroutine, the program flow returns to the instruction following the `call` instruction.

(2) Delayed branch (d bit (bit 7) = 1)

`call.d %rb`

When `call.d %rb` is specified, the d bit (bit 7) in the instruction code is set and the following instruction becomes a delayed slot instruction.

The delayed slot instruction is executed before branching to the subroutine. Therefore the address (PC + 4) of the instruction that follows the delayed slot instruction is stored into the stack as the return address.

When the `call.d` instruction is executed, interrupts cannot occur because they are masked between the `call.d` and delayed slot instructions.

Example

```
call  %r0   ; Calls the subroutine that starts from pc + 2 + r0.
```

Caution

When the `call.d` instruction (delayed branch) is used, be careful to ensure that the next instruction is limited to those that can be used as a delayed slot instruction. If any other instruction is executed, the program may operate indeterminately. For the usable instructions, refer to the instruction list in the Appendix.

call *sign10*
call.d *sign10*

Function

PC relative subroutine call

Standard) `call`: sp ← sp - 4, A[sp] ← pc + 2, pc ← pc + 2 + *sign10* × 2
`call.d`: sp ← sp - 4, A[sp] ← pc + 4, pc ← pc + 2 + *sign10* × 2
Extension 1) `call`: sp ← sp - 4, A[sp] ← pc + 2, pc ← pc + 2 + sign24
`call.d`: sp ← sp - 4, A[sp] ← pc + 4, pc ← pc + 2 + sign24
Extension 2) Unusable

Code

15	14	13	12	11	10	9–0	
0	0	0	1	1	0	*sign10*	`call`
0	0	0	1	1	1	*sign10*	`call.d`

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Signed PC relative

CLK

`call` Four cycles
`call.d` Three cycles (when a one-cycle delayed slot instruction follows), Four cycles (other)

Description

(1) Standard

```
call   sign10   ; = "call sign11"
                ; sign10 = sign11(10:1), sign11(0) = 0
```

Stores the address of the following instruction into the stack, then doubles the signed 10-bit immediate *sign10* and adds it to the PC (PC + 2) for calling the subroutine that starts from the address. The *sign10* specifies a word address in 16-bit units. When the `ret` instruction is executed in the subroutine, the program flow returns to the instruction following the `call` instruction.
The *sign10* (×2) allows branches within the range of PC - 1,022 to PC + 1,024.

(2) Extension 1

```
ext    imm13    ; = sign24(23:11)
call   sign10   ; = "call sign24"
                ; sign10 = sign24(10:1), sign24(0) = 0
```

The `ext` instruction extends the displacement into 24 bits using its 13-bit immediate *imm13*. The 24-bit displacement is added to the PC.
The *sign24* allows branches within the range of PC - 8,388,606 to PC + 8,388,608.

(3) Delayed branch (d bit (bit 10) = 1)

```
call.d  sign10
```

When `call.d sign10` is specified, the d bit (bit 10) in the instruction code is set and the following instruction becomes a delayed slot instruction. The delayed slot instruction is executed before branching to the subroutine. Therefore the address (PC + 4) of the instruction that follows the delayed slot instruction is stored into the stack as the return address.
When the `call.d` instruction is executed, interrupts cannot occur because they are masked between the `call.d` and delayed slot instructions.

Example

```
ext    0x1fff
call   0x0         ; Calls the subroutine that starts from the
                   ; address specified by pc + 2 - 0x800.
```

Caution

When the `call.d` instruction (delayed branch) is used, be careful to ensure that the next instruction is limited to those that can be used as a delayed slot instruction. If any other instruction is executed, the program may operate indeterminately. For the usable instructions, refer to the instruction list in the Appendix.

calla %rb
calla.d %rb

Function

PC absolute subroutine call

Standard) `calla`: sp ← sp - 4, A[sp] ← pc + 2, pc ← *rb*
`calla.d`: sp ← sp - 4, A[sp] ← pc + 4, pc ← *rb*

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	0	0	0	0	0	1	0	0	0	0	1	*rb*			calla
0	0	0	0	0	0	0	1	1	0	0	0	1	*rb*			calla.d

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

PC absolute

CLK

`calla` Four cycles
`calla.d` Three cycles (when a one-cycle delayed slot instruction follows), Four cycles (other)

Description

(1) Standard

`calla %rb`

Stores the address of the following instruction into the stack, then sets the contents of the *rb* register to the PC for calling the subroutine that starts from the address set to the PC. The LSB of the *rb* register is invalid and is always handled as 0. When the `ret` instruction is executed in the subroutine, the program flow returns to the instruction following the `calla` instruction.

(2) Delayed branch (d bit (bit 7) = 1)

`calla.d %rb`

When `calla.d` is specified, the d bit (bit 7) in the instruction code is set and the following instruction becomes a delayed slot instruction.
The delayed slot instruction is executed before branching to the subroutine. Therefore the address (PC + 4) of the instruction that follows the delayed slot instruction is stored into the stack as the return address.
When the `calla.d` instruction is executed, interrupts cannot occur because they are masked between the `calla.d` and delayed slot instructions.

Example

```
calla  %r0       ; Calls the subroutine that starts from the
                 ; address stored in the r0 register.
```

Caution

When the `calla.d` instruction (delayed branch) is used, be careful to ensure that the next instruction is limited to those that can be used as a delayed slot instruction. If any other instruction is executed, the program may operate indeterminately. For the usable instructions, refer to the instruction list in the Appendix.

calla *imm7*
calla.d *imm7*

Function

PC absolute subroutine call

Standard)	`calla:`	sp ← sp - 4, A[sp] ← pc + 2, pc ← *imm7*
	`calla.d:`	sp ← sp - 4, A[sp] ← pc + 4, pc ← *imm7*
Extension 1)	`calla:`	sp ← sp - 4, A[sp] ← pc + 2, pc ← *imm20*
	`calla.d:`	sp ← sp - 4, A[sp] ← pc + 4, pc ← *imm20*
Extension 2)	`calla:`	sp ← sp - 4, A[sp] ← pc + 2, pc ← *imm24*
	`calla.d:`	sp ← sp - 4, A[sp] ← pc + 4, pc ← *imm24*

Code

15	14	13	12	11	10	9	8	7	6–0	
0	0	0	0	0	1	0	1	0	*imm7*	`calla`
0	0	0	0	0	1	0	1	1	*imm7*	`calla.d`

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

PC absolute

CLK

`calla` Four cycles
`calla.d` Three cycles (when a one-cycle delayed slot instruction follows), Four cycles (other)

Description

(1) Standard

```
calla  imm7
```

Stores the address of the following instruction into the stack, then sets the 7-bit immediate *imm7* to the PC for calling the subroutine that starts from the address set to the PC. The LSB of the *imm7* is invalid and is always handled as 0. When the `ret` instruction is executed in the subroutine, the program flow returns to the instruction following the `calla` instruction.

(2) Extension 1

```
ext    imm13   ; = imm20(19:7)
call   imm7    ; = "call imm20", imm7 = imm20(6:0)
```

The `ext` instruction extends the destination address into 20 bits using its 13-bit immediate *imm13*. The 20-bit destination address is set to the PC.

(3) Extension 2

```
ext    imm4    ; imm4(3:0) = imm24(23:20)
ext    imm13   ; = imm24(19:7)
call   imm7    ; = "call imm24", imm7 = imm24(6:0)
```

The 24-bit destination address is set to the PC.

(4) Delayed branch (d bit (bit 7) = 1)

```
calla.d  imm7
```

When `calla.d` is specified, the d bit (bit 7) in the instruction code is set and the following instruction becomes a delayed slot instruction. The delayed slot instruction is executed before branching to the subroutine. Therefore the address (PC + 4) of the instruction that follows the delayed slot instruction is stored into the stack as the return address.
When the `calla.d` instruction is executed, interrupts cannot occur because they are masked between the `calla.d` and delayed slot instructions.

Example

```
ext     0x1fff
calla   0x0       ; Calls the subroutine that starts from
                  ; address 0xfff80.
```

Caution

When the `calla.d` instruction (delayed branch) is used, be careful to ensure that the next instruction is limited to those that can be used as a delayed slot instruction. If any other instruction is executed, the program may operate indeterminately. For the usable instructions, refer to the instruction list in the Appendix.

cmc %rd, %rs
cmc/c %rd, %rs
cmc/nc %rd, %rs

Function

16-bit comparison with carry

Standard) *rd*(15:0) - *rs*(15:0) - C
Extension 1) *rs*(15:0) - *imm13*(zero extended) - C
Extension 2) *rs*(15:0) - *imm16* - C

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	1	1	1	1	*rd*			1	0	0	1	*rs*			cmc
0	0	1	1	1	1	*rd*			0	0	0	1	*rs*			cmc/c
0	0	1	1	1	1	*rd*			0	1	0	1	*rs*			cmc/nc

Flag

IL	IE	C	V	Z	N	
–	–	↔	↔	↔	↔	cmc
–	–	–	↔	↔	↔	cmc/c, cmc/nc

Mode

Src: Register direct `%rs` = `%r0` to `%r7`
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
cmc   %rd,%rs          ; rd - rs - C
```

Subtracts the contents of the *rs* register and C (carry) flag from the contents of the *rd* register, and sets or resets the flags (C, V, Z and N) according to the results. The operation is performed in 16-bit size. It does not change the contents of the *rd* register.

(2) Extension 1

```
ext   imm13
cmc   %rd,%rs          ; rs - imm13 - C
```

Subtracts the contents of the 13-bit immediate *imm13* and C (carry) flag from the contents of the *rs* register, and sets or resets the flags (C, V, Z and N) according to the results. The *imm13* is zero-extended into 16 bits prior to the operation. The operation is performed in 16-bit size. It does not change the contents of the *rd* and *rs* registers.

* This combination does not use the *rd* register value for comparison.

(3) Extension 2

```
ext   imm3             ; imm3(2:0) = imm16(15:13)
ext   imm13            ; = imm16(12:0)
cmc   %rd,%rs          ; rs - imm16 - C
```

Subtracts the contents of the 16-bit immediate *imm16* and C (carry) flag from the contents of the *rs* register, and sets or resets the flags (C, V, Z and N) according to the results. The operation is performed in 16-bit size. It does not change the contents of the *rd* and *rs* registers.

* This combination does not use the *rd* register value for comparison.

(4) Conditional execution

The `/c` or `/nc` suffix on the opcode specifies conditional execution.
`cmc/c` Executed as `cmc` when the C flag is 1 or executed as `nop` when the flag is 0
`cmc/nc` Executed as `cmc` when the C flag is 0 or executed as `nop` when the flag is 1

In this case, the `ext` instruction can be used to extend the operand.
The conditional execution instruction above sets/resets the flags (V, Z and N) according to the results if it is executed.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
cmc  %r0,%r1  ; Changes the flags according to the results of
              ; r0 - r1 - C.
```

(2)
```
ext  0x1fff
cmc  %r1,%r2  ; Changes the flags according to the results of
              ; r2 - 0x1fff - C.
```

cmc %rd, sign7

Function

16-bit comparison with carry

Standard) *rd*(15:0) - *sign7*(sign extended) - C

Extension 1) *rd*(15:0) - *sign16* - C

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	1	0	1	*rd*			*sign7*						

Flag

IL	IE	C	V	Z	N
–	–	↔	↔	↔	↔

Mode

Src: Immediate data (signed)

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
cmc   %rd,sign7       ; rd - sign7 - C
```

Subtracts the contents of the signed 7-bit immediate *sign7* and C (carry) flag from the contents of the *rd* register, and sets or resets the flags (C, V, Z and N) according to the results. The *sign7* is sign-extended into 16 bits prior to the operation. The operation is performed in 16-bit size. It does not change the contents of the *rd* register.

(2) Extension 1

```
ext   imm9            ; imm9(8:0) = sign16(15:7)
cmc   %rd,sign7       ; rd - sign16 - C, sign7 = sign16(6:0)
```

Subtracts the contents of the signed 16-bit immediate *sign16* and C (carry) flag from the contents of the *rd* register, and sets or resets the flags (C, V, Z and N) according to the results. The operation is performed in 16-bit size. It does not change the contents of the *rd* register.

(3) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
cmc   %r0,0x7f ; Changes the flags according to the results of
               ; r0 - 0x7f - C.
```

(2)
```
ext   0x1ff
cmc   %r1,0x7f ; Changes the flags according to the results of
               ; r1 - 0xffff - C.
```

cmp *%rd, %rs*
cmp/c *%rd, %rs*
cmp/nc *%rd, %rs*

Function

16-bit comparison

Standard) *rd*(15:0) - *rs*(15:0)
Extension 1) *rs*(15:0) - *imm13*(zero extended)
Extension 2) *rs*(15:0) - *imm16*

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	1	1	1	1		*rd*		1	0	0	0		*rs*		`cmp`
0	0	1	1	1	1		*rd*		0	0	0	0		*rs*		`cmp/c`
0	0	1	1	1	1		*rd*		0	1	0	0		*rs*		`cmp/nc`

Flag

IL	IE	C	V	Z	N	
–	–	↔	↔	↔	↔	`cmp`
–	–	–	↔	↔	↔	`cmp/c`, `cmp/nc`

Mode

Src: Register direct `%rs` = `%r0` to `%r7`
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
cmp   %rd,%rs          ; rd - rs
```

Subtracts the contents of the *rs* register from the contents of the *rd* register, and sets or resets the flags (C, V, Z and N) according to the results. The operation is performed in 16-bit size. It does not change the contents of the *rd* register.

(2) Extension 1

```
ext   imm13
cmp   %rd,%rs          ; rs - imm13
```

Subtracts the 13-bit immediate *imm13* from the contents of the *rs* register, and sets or resets the flags (C, V, Z and N) according to the results. The *imm13* is zero-extended into 16 bits prior to the operation. The operation is performed in 16-bit size. It does not change the contents of the *rd* and *rs* registers.

∗ This combination does not use the *rd* register value for comparison.

(3) Extension 2

```
ext   imm3             ; imm3(2:0)  = imm16(15:13)
ext   imm13            ; = imm16(12:0)
cmp   %rd,%rs          ; rs - imm16
```

Subtracts the 16-bit immediate *imm16* from the contents of the *rs* register, and sets or resets the flags (C, V, Z and N) according to the results. The operation is performed in 16-bit size. It does not change the contents of the *rd* and *rs* registers.

∗ This combination does not use the *rd* register value for comparison.

(4) Conditional execution

The `/c` or `/nc` suffix on the opcode specifies conditional execution.
`cmp/c` Executed as `cmp` when the C flag is 1 or executed as `nop` when the flag is 0
`cmp/nc` Executed as `cmp` when the C flag is 0 or executed as `nop` when the flag is 1

In this case, the `ext` instruction can be used to extend the operand.
The conditional execution instruction above sets/resets the flags (V, Z and N) according to the results if it is executed.

(5) Delayed slot instruction
This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
cmp  %r0,%r1  ; Changes the flags according to the results of
              ; r0 - r1.
```

(2)
```
ext  0x1
ext  0x1fff   ; Changes the flags according to the results of
cmp  %r1,%r2  ; r2 - 0x3fff.
```

cmp %rd, sign7

Function

16-bit comparison

Standard) *rd*(15:0) - *sign7*(sign extended)

Extension 1) *rd*(15:0) - *sign16*

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	1	0	0		*r d*				*sign7*				

Flag

IL	IE	C	V	Z	N
–	–	↔	↔	↔	↔

Mode

Src: Immediate data (signed)

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
cmp   %rd,sign7        ; rd - sign7
```

Subtracts the signed 7-bit immediate *sign7* from the contents of the *rd* register, and sets or resets the flags (C, V, Z and N) according to the results. The *sign7* is sign-extended into 16 bits prior to the operation. The operation is performed in 16-bit size. It does not change the contents of the *rd* register.

(2) Extension 1

```
ext   imm9             ; imm9(8:0) = sign16(15:7)
cmp   %rd,sign7        ; rd - sign16, sign7 = sign16(6:0)
```

Subtracts the signed 16-bit immediate *sign16* from the contents of the *rd* register, and sets or resets the flags (C, V, Z and N) according to the results. The operation is performed in 16-bit size. It does not change the contents of the *rd* register.

(3) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
cmp   %r0,0x3f ; Changes the flags according to the results of
               ; r0 - 0x3f.
```

(2)
```
ext   0x1ff
cmp   %r1,0x7f ; Changes the flags according to the results of
               ; r1 - 0xffff.
```

cmp.a %rd, %rs
cmp.a/c %rd, %rs
cmp.a/nc %rd, %rs

Function

24-bit comparison

Standard) *rd*(23:0) - *rs*(23:0)
Extension 1) *rs*(23:0) - *imm13*(zero extended)
Extension 2) *rs*(23:0) - *imm24*

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	1	1	0	1	rd			1	0	0	0	rs			`cmp.a`
0	0	1	1	0	1	rd			0	0	0	0	rs			`cmp.a/c`
0	0	1	1	0	1	rd			0	1	0	0	rs			`cmp.a/nc`

Flag

IL	IE	C	V	Z	N	
–	–	↔	–	↔	–	`cmp.a`
–	–	–	–	↔	–	`cmp.a/c`, `cmp.a/nc`

Mode

Src: Register direct `%rs` = `%r0` to `%r7`
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
cmp.a   %rd,%rs       ; rd - rs
```

Subtracts the contents of the *rs* register from the contents of the *rd* register, and sets or resets the flags (C and Z) according to the results. It does not change the contents of the *rd* register.

(2) Extension 1

```
ext     imm13
cmp.a   %rd,%rs       ; rs - imm13
```

Subtracts the 13-bit immediate *imm13* from the contents of the *rs* register, and sets or resets the flags (C and Z) according to the results. The *imm13* is zero-extended into 24 bits prior to the operation. It does not change the contents of the *rd* and *rs* registers.

∗ This combination does not use the *rd* register value for comparison.

(3) Extension 2

```
ext     imm11         ; imm11(10:0) = imm24(23:13)
ext     imm13         ; = imm24(12:0)
cmp.a   %rd,%rs       ; rs - imm24
```

Subtracts the 24-bit immediate *imm24* from the contents of the *rs* register, and sets or resets the flags (C and Z) according to the results. It does not change the contents of the *rd* and *rs* registers.

∗ This combination does not use the *rd* register value for comparison.

(4) Conditional execution

The `/c` or `/nc` suffix on the opcode specifies conditional execution.
`cmp.a/c` Executed as `cmp.a` when the C flag is 1 or executed as `nop` when the flag is 0
`cmp.a/nc` Executed as `cmp.a` when the C flag is 0 or executed as `nop` when the flag is 1

In this case, the `ext` instruction can be used to extend the operand.
The conditional execution instruction above sets/resets the flags (V and Z) according to the results if it is executed.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
cmp.a  %r0,%r1   ; Changes the flags according to the results of
                 ; r0 - r1.
```

(2)
```
ext    0x1
ext    0x1fff
cmp.a  %r1,%r2   ; Changes the flags according to the results of
                 ; r2 - 0x3fff.
```

cmp.a *%rd, imm7*

Function

24-bit comparison

Standard) *rd*(23:0) - *imm7*(zero extended)
Extension 1) *rd*(23:0) - *imm20*(zero extended)
Extension 2) *rd*(23:0) - *imm24*

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	0	0		*rd*					*imm7*			

Flag

IL	IE	C	V	Z	N
–	–	↔	–	↔	–

Mode

Src: Immediate data (unsigned)
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
cmp.a   %rd,imm7      ; rd - imm7
```

Subtracts the 7-bit immediate *imm7* from the contents of the *rd* register, and sets or resets the flags (C and Z) according to the results. The *imm7* is zero-extended into 24 bits prior to the operation. It does not change the contents of the *rd* register.

(2) Extension 1

```
ext     imm13         ; = imm20(19:7)
cmp.a   %rd,imm7      ; rd - imm20, imm7 = imm20(6:0)
```

Subtracts the 20-bit immediate *imm20* from the contents of the *rd* register, and sets or resets the flags (C and Z) according to the results. The *imm20* is zero-extended into 24 bits prior to the operation. It does not change the contents of the *rd* register.

(3) Extension 2

```
ext     imm4          ; imm4(3:0) = imm24(23:20)
ext     imm13         ; = imm24(19:7)
cmp.a   %rd,imm7      ; rd - imm24, imm7 = imm24(6:0)
```

Subtracts the 24-bit immediate *imm24* from the contents of the *rd* register, and sets or resets the flags (C and Z) according to the results. It does not change the contents of the *rd* register.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
cmp.a   %r0,0x7f  ; Changes the flags according to the results of
                  ; r0 - 0x7f.
```

(2)
```
ext     0xf
ext     0x1fff
cmp.a   %r1,0x7f  ; Changes the flags according to the results of
                  ; r1 - 0xffffff.
```

cv.ab *%rd*, *%rs*

Function

Data conversion from byte to 24 bits

Standard) *rd*(23:8) ← *rs*(7), *rd*(7:0) ← *rs*(7:0)

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	1	0		*rd*		0	1	1	1		*rs*	

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register direct `%rd` = `%r0` to `%r7`

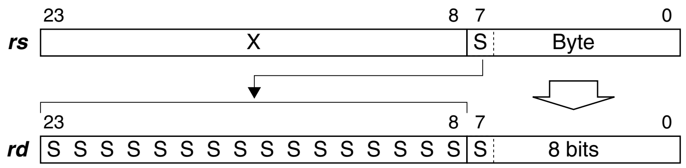
CLK

One cycle

Description

(1) Standard

The eight low-order bits of the *rs* register are transferred to the *rd* register after being sign-extended to 24 bits.



(2) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit.

Example

When the R1 register contains 0x80

```
cv.ab  %r0,%r1   ; r0 = 0xffff80
```

cv.al *%rd, %rs*

Function

Data conversion from 32 bits to 24 bits

Standard) *rd*(23:16) ← *rs*(7:0), *rd*(15:0) ← *rd*(15:0)
Extension 1) Unusable
Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	1	0		*rd*		1	1	1	1		*rs*	

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`
Dst: Register direct `%rd` = `%r0` to `%r7`

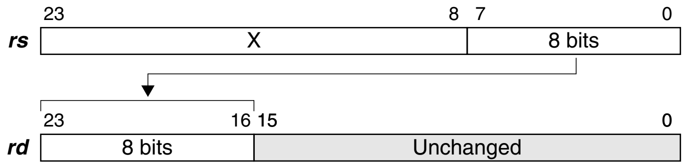
CLK

One cycle

Description

(1) Standard

The eight low-order bits of the *rs* register are transferred to the eight high-order bits of the *rd* register.



(2) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit.

Example

When the R1 register contains 0xff and the R0 register contains 0x0

```
cv.al  %r0,%r1   ; r0 = 0xff0000
```

cv.as %rd, %rs

Function

Data conversion from 16 bits to 24 bits

Standard) *rd*(23:16) ← *rs*(15), *rd*(15:0) ← *rs*(15:0)

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	1	0		*rd*		1	0	1	1		*rs*	

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register direct `%rd` = `%r0` to `%r7`

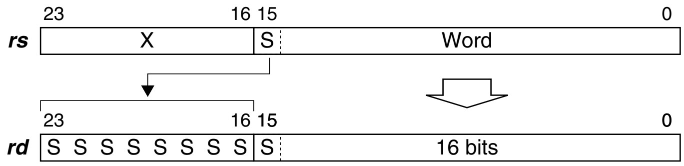
CLK

One cycle

Description

(1) Standard

The 16 low-order bits of the *rs* register are transferred to the *rd* register after being sign-extended to 24 bits.



(2) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit.

Example

When the R1 register contains 0x8000

```
cv.as   %r0,%r1   ; r0 = 0xff8000
```

cv.la %rd, %rs

Function

Data conversion from 24 bits to 32 bits

Standard) *rd*(23:8) ← 0, *rd*(7:0) ← *rs*(23:16)

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	1	0		*rd*		0	1	1	0		*rs*	

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register direct `%rd` = `%r0` to `%r7`

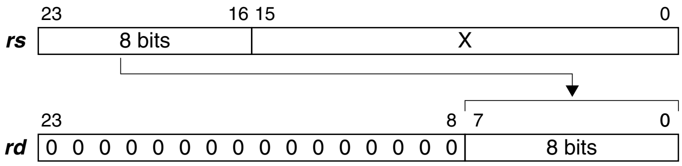
CLK

One cycle

Description

(1) Standard

The eight high-order bits of the *rs* register are transferred to the eight low-order bits of the *rd* register. The 16 high-order bits of the *rd* register are set to 0.



(2) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit.

Example

When the R1 register contains 0x800000

```
cv.la  %r0,%r1   ; r0 = 0x000080
```

cv.ls %rd, %rs

Function

Data conversion from 16 bits to 32 bits

Standard) *rd*(23:16) ← 0, *rd*(15:0) ← *rs*(15)

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	1	0		*rd*		1	0	1	0		*rs*	

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

Bit 15 (sign bit of 16-bit data) of the *rs* register is transferred to the 16 low-order bits of the *rd* register. The eight high-order bits of the *rd* register are set to 0.

(2) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit.

Example

When the R1 register contains 0x008000

```
cv.ls   %r0,%r1   ; r0 = 0x00ffff
```

di

Function

Disable interrupts
Standard) psr(IE) ← 0
Extension 1) Unusable
Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	0	0	0	0	1	0	0	0	0	0	0

Flag

IL	IE	C	V	Z	N
–	0	–	–	–	–

Mode

–

CLK

One cycle

Description

(1) Standard

Resets the IE bit in the PSR to disable external maskable interrupts.

The reset interrupt, address misaligned interrupt, and NMI will be accepted even if the IE bit is set to 0.

(2) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit.

Example

```
di        ; Disables external maskable interrupts.
```

Caution

Maskable interrupts are disabled from the third cycle after the `di` instruction has been executed.

`di`
Instruction 1 ← 1-cycle instruction
Instruction 2 ← 1-cycle instruction
Instruction 3 ← Interrupts are disabled from this instruction.

Example: Interrupt disabled periods using the `di` and `ei` instructions

```
ld %r2,%r3      ← Interrupt enabled
di              ← Interrupt enabled
ld.a %r0,%r1    ← Interrupt enabled
ld.b %r2,%r3    ← Interrupt enabled
ld %r4,%r5      ← Interrupt disabled
ei              ← Interrupt disabled
add %r4,%r5     ← Interrupt disabled
sub %r6,%r7     ← Interrupt disabled
cmp %r0,%r1     ← Interrupt enabled
```

ei

Function

Enable interrupts

Standard) psr(IE) ← 1

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	0	0	0	0	0	1	0	0	0	0	0

Flag

IL	IE	C	V	Z	N
–	1	–	–	–	–

Mode

–

CLK

One cycle

Description

(1) Standard

Sets the IE bit in the PSR to enable external maskable interrupts.

(2) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit.

Example

```
ei        ; Enables external maskable interrupts.
```

Caution

Maskable interrupts are enabled from the third cycle after the `ei` instruction has been executed.

`ei`	
Instruction 1	← 1-cycle instruction
Instruction 2	← 1-cycle instruction
Instruction 3	← Interrupts are enabld from this instruction.

Example: Interrupt disabled periods using the `di` and `ei` instructions

`ld %r2,%r3`	← Interrupt enabled
`di`	← Interrupt enabled
`ld.a %r0,%r1`	← Interrupt enabled
`ld.b %r2,%r3`	← Interrupt enabled
`ld %r4,%r5`	← Interrupt disabled
`ei`	← Interrupt disabled
`add %r4,%r5`	← Interrupt disabled
`sub %r6,%r7`	← Interrupt disabled
`cmp %r0,%r1`	← Interrupt enabled

ext *imm13*

Function

Immediate extension

Standard) Extends the immediate data/operand of the following instruction
Extension 1) Unusable
Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	0	*imm13*												

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Immediate data (unsigned)

CLK

One cycle

Description

Extends the immediate data or operand of the following instruction.
When extending an immediate data, the immediate data in the `ext` instruction will be placed on the high-order side and the immediate data in the target instruction to be extended is placed on the low-order side.

Up to two `ext` *`imm3`* instructions can be used sequentially. In this case, the immediate data in the first `ext` instruction is placed on the most upper part. When three or more `ext` instructions have been described sequentially, the last two are effective and others are ignored.
See descriptions of each instruction for the extension contents and the usage.

Interrupts for the `ext` instruction (not including reset and debug break) are masked in the hardware, and interrupt handling is determined when the target instruction to be extended is executed. In this case, the return address from interrupt handling is the beginning of the `ext` instruction.

Example

```
ext     0x7ff
ext     0x1fff
add.a   %r1,%r2   ; r1 = r2 + 0xffffff
```

Caution

When a load instruction that transfers data between memory and a register follows the `ext` instruction, an address misaligned interrupt may occur before executing the load instruction (if the address that is specified with the immediate data in the `ext` instruction as the displacement is not a boundary address according to the transfer data size). When an address misaligned interrupt occurs, the interrupt handling saves the address of the load instruction into the stack as the return address. If the interrupt handler routine is returned by simply executing the `reti` instruction, the previous `ext` instruction is invalidated. Therefore, it is necessary to modify the return address in that case.

halt

Function

HALT

Standard) Sets the processor to HALT mode

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	0

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

–

CLK

Six cycles

Description

Sets the processor to HALT mode for power saving.
Program execution is halted at the same time that the S1C17 Core executes the `halt` instruction, and the processor enters HALT mode.
HALT mode commonly turns off only the S1C17 Core operation, note, however that modules to be turned off depend on the implementation of the clock control circuit outside the core.

Initial reset is one cause that can bring the processor out of HALT mode. Other causes depend on the implementation of the clock control circuit outside the S1C17 Core.
Initial reset, maskable external interrupts, NMI, and debug interrupts are commonly used for canceling HALT mode.
The interrupt enable/disable status set in the processor does not affect the cancellation of HALT mode even if an interrupt signal is used as the cancellation. In other words, interrupt signals are able to cancel HALT mode even if the IE flag in PSR or the interrupt enable bits in the interrupt controller (depending on the implementation) are set to disable interrupts.
When the processor is taken out of HALT mode using an interrupt that has been enabled (by the interrupt controller and IE flag), the corresponding interrupt handler routine is executed. Therefore, when the interrupt handler routine is terminated by the `reti` instruction, the processor returns to the instruction next to `halt`.
When the interrupt has been disabled, the processor restarts the program from the instruction next to `halt` after the processor is taken out of HALT mode.

Refer to the technical manual of each model for details of HALT mode.

Example

```
halt              ; Sets the processor in HALT mode.
```

int *imm5*

Function

Software interrupt

Standard) sp ← sp - 4, A[sp] ← {psr, pc + 2}, pc ← TTBR + (vector No. = *imm5*) × 4
Extension 1) Unusable
Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	0	1	0	0	0	*imm5*					0	1

Flag

IL	IE	C	V	Z	N
–	0	–	–	–	–

Mode

Immediate data (unsigned)

CLK

Three cycles

Description

Generates the interrupt of the vector number specified with the *imm5*.
The `int` instruction saves the address of the next instruction and the contents of the PSR into the stack, then reads the specified interrupt vector from the vector table and sets it to the PC. By this processing, the program flow branches to the specified interrupt handler routine.

imm5	Vector No.	Vector address	Cause of interrupt
0x00	0	TTBR + 0x00	Reset interrupt
0x01	1	TTBR + 0x04	Address misaligned interrupt
0x02	2	TTBR + 0x08	NMI
0x03	3	TTBR + 0x0c	External maskable interrupt 0x03
:	:	:	:
0x1f	31	TTBR + 0x7c	External maskable interrupt 0x1f

The TTBR is the vector table base address.
The `reti` instruction should be used for return from the handler routine.

Example

```
int  2    ; Generates an NMI.
```

intl *imm5*, *imm3*

Function

Software interrupt with interrupt level setting

Standard) sp ← sp - 4, A[sp] ← {psr, pc + 2}, pc ← TTBR + (vector No. = *imm5*) × 4, psr(IL) ← *imm3*

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	0	1	*imm3*			*imm5*					1	1

Flag

IL	IE	C	V	Z	N
↔	0	–	–	–	–

Mode

Immediate data (unsigned)

CLK

Three cycles

Description

Generates the interrupt of the vector number specified with the *imm5*.

The `intl` instruction saves the address of the next instruction and the contents of the PSR into the stack, then reads the specified interrupt vector from the vector table and sets it to the PC. By this processing, the program flow branches to the specified interrupt handler routine. In addition to this, the *imm3* value is set to the IL bits in the PSR (interrupt level) to disable interrupts of which the interrupt level is lower than the *imm3* while the interrupt handler routine is executed.

The altered IL bits are restored to the value before the `intl` instruction is executed when the interrupt handler routine is terminated by the `reti` instruction.

Example

```
intl  0x3,0x2    ; Generates an external maskable interrupt 0x3
                 ; and set the IL bits to 0x2.
```

jpa %rb
jpa.d %rb

Function Unconditional PC absolute jump

Standard) pc ← *rb*

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	0	0	0	0	0	1	0	1	0	0	1	*rb*			`jpa`
0	0	0	0	0	0	0	1	1	1	0	0	1	*rb*			`jpa.d`

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode PC absolute

CLK

`jpa` Three cycles

`jpa.d` Two cycles (when a one-cycle delayed slot instruction follows), Three cycles (other)

Description (1) Standard

```
jpa   %rb
```

The content of the *rb* register is loaded to the PC, and the program branches to that address. The LSB of the *rb* register is ignored and is always handled as 0.

(2) Delayed branch (d bit (bit 7) = 1)

```
jpa.d   %rb
```

For the `jpa.d` instruction, the next instruction becomes a delayed slot instruction. A delayed slot instruction is executed before the program branches. Interrupts are masked in intervals between the `jpa.d` instruction and the next instruction, so no interrupts occur.

Example

```
jpa  %r0  ; Jumps to the address specified by the r0 register.
```

Caution When the `jpa.d` instruction (delayed branch) is used, be careful to ensure that the next instruction is limited to those that can be used as a delayed slot instruction. If any other instruction is executed, the program may operate indeterminately. For the usable instructions, refer to the instruction list in the Appendix.

jpa imm7
jpa.d imm7

Function

Unconditional PC absolute jump

Standard) pc ← *imm7*
Extension 1) pc ← *imm20*
Extension 2) pc ← *imm24*

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	0	0	0	0	1	1	0	*imm7*							`jpa`
0	0	0	0	0	0	1	1	1	*imm7*							`jpa.d`

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

PC absolute

CLK

`jpa` Three cycles
`jpa.d` Two cycles (when a one-cycle delayed slot instruction follows), Three cycles (other)

Description

(1) Standard

```
jpa  imm7
```

The 7-bit immediate *imm7* is loaded to the PC, and the program branches to that address. The LSB of the *imm7* is ignored and is always handled as 0.

(2) Extension 1

```
ext  imm13    ; = imm20(19:7)
jpa  imm7     ; = "jpa imm20", imm7 = imm20(6:0)
```

The `ext` instruction extends the destination address into 20 bits using its 13-bit immediate *imm13*. The 20-bit destination address is set to the PC.

(3) Extension 2

```
ext  imm4     ; imm4(3:0) = imm24(23:20)
ext  imm13    ; = imm24(19:7)
jpa  imm7     ; = "jpa imm24", imm7 = imm24(6:0)
```

The 24-bit destination address is set to the PC.

(4) Delayed branch (d bit (bit 7) = 1)

```
jpa.d  imm7
```

For the `jpa.d` instruction, the next instruction becomes a delayed slot instruction. A delayed slot instruction is executed before the program branches. Interrupts are masked in intervals between the `jpa.d` instruction and the next instruction, so no interrupts occur.

Example

```
ext  0x300
jpa  0x00 ; Jumps to the address 0x18000.
```

Caution

When the `jpa.d` instruction (delayed branch) is used, be careful to ensure that the next instruction is limited to those that can be used as a delayed slot instruction. If any other instruction is executed, the program may operate indeterminately. For the usable instructions, refer to the instruction list in the Appendix.

jpr %rb
jpr.d %rb

Function

Unconditional PC relative jump

Standard) pc ← pc + 2 + *rb*

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	0	0	0	0	0	1	0	1	0	0	0	*rb*			jpr
0	0	0	0	0	0	0	1	1	1	0	0	0	*rb*			jpr.d

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Signed PC relative

CLK

`jpr` Three cycles

`jpr.d` Two cycles (when a one-cycle delayed slot instruction follows), Three cycles (other)

Description

(1) Standard

`jpr %rb`

The content of the *rb* register is added to the PC (PC + 2), and the program branches to that address. The LSB of the *rb* register is ignored and is always handled as 0.

(2) Delayed branch (d bit (bit 7) = 1)

`jpr.d %rb`

For the `jpr.d` instruction, the next instruction becomes a delayed slot instruction. A delayed slot instruction is executed before the program branches. Interrupts are masked in intervals between the `jpr.d` instruction and the next instruction, so no interrupts occur.

Example

```
jpr  %r0          ; pc ← pc + 2 + r0
```

Caution

When the `jpr.d` instruction (delayed branch) is used, be careful to ensure that the next instruction is limited to those that can be used as a delayed slot instruction. If any other instruction is executed, the program may operate indeterminately. For the usable instructions, refer to the instruction list in the Appendix.

jpr *sign10*
jpr.d *sign10*

Function

Unconditional PC relative jump

Standard) pc ← pc + 2 + *sign10* × 2
Extension 1) pc ← pc + 2 + *sign24*
Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	0	1	0	0	*sign10*										jpr
0	0	0	1	0	1	*sign10*										jpr.d

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Signed PC relative

CLK

`jpr` Three cycles
`jpr.d` Two cycles (when a one-cycle delayed slot instruction follows), Three cycles (other)

Description

(1) Standard

```
jpr  sign10   ; = "jp sign11", sign10 = sign11(10:1), sign11(0)=0
```

Doubles the signed 10-bit immediate *sign10* and adds it to the PC (PC + 2). The program flow branches to the address. The *sign10* specifies a word address in 16-bit units.
The *sign10* (×2) allows branches within the range of PC - 1,022 to PC + 1,024.

(2) Extension 1

```
ext  imm13    ; = sign24(23:11)
jpr  sign10   ; = "jpr sign24", sign10 = sign24(10:1), sign24(0)=0
```

The `ext` instruction extends the displacement to be added to the PC (PC + 2) into 24 bits using its 13-bit immediate *imm13*.
The *sign24* allows branches within the range of PC - 8,388,606 to PC + 8,388608.

(3) Delayed branch (d bit (bit 10) = 1)

```
jpr.d  sign10
```

For the `jpr.d` instruction, the next instruction becomes a delayed slot instruction. A delayed slot instruction is executed before the program branches. Interrupts are masked in intervals between the `jpr.d` instruction and the next instruction, so no interrupts occur.

Example

```
ext  0x20
jpr  0x00 ; Jumps to the address specified by pc + 2 + 0x10000.
```

Caution

When the `jpr.d` instruction (delayed branch) is used, be careful to ensure that the next instruction is limited to those that can be used as a delayed slot instruction. If any other instruction is executed, the program may operate indeterminately. For the usable instructions, refer to the instruction list in the Appendix.

jreq *sign7*
jreq.d *sign7*

Function

Conditional PC relative jump

Standard) pc ← pc + 2 + *sign7* × 2 if Z is true
Extension 1) pc ← pc + 2 + *sign21* if Z is true
Extension 2) pc ← pc + 2 + *sign24* if Z is true

Code

15	14	13	12	11	10	9	8	7	6–0	
0	0	0	0	1	1	1	0	0	*sign7*	jreq
0	0	0	0	1	1	1	0	1	*sign7*	jreq.d

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Signed PC relative

CLK

`jreq` Two cycles (when not branched), Three cycles (when branched)
`jreq.d` Two cycles (when a one-cycle delayed slot instruction follows), Three cycles (other)

Description

(1) Standard

```
jreq   sign7   ; = "jreq sign8", sign7 = sign8(7:1), sign8(0)=0
```

If the condition below has been met, this instruction doubles the signed 7-bit immediate *sign7* and adds it to the PC (PC + 2) for branching the program flow to the address. It does not branch if the condition has not been met.

• Z flag = 1 (e.g. "A = B" has resulted by `cmp A,B`)

The *sign7* specifies a word address in 16-bit units.

The *sign7* (×2) allows branches within the range of PC - 126 to PC + 128.

(2) Extension 1

```
ext    imm13   ; = sign21(20:8)
jreq   sign7   ; = "jreq sign21", sign7 = sign21(7:1), sign21(0)=0
```

The `ext` instruction extends the displacement to be added to the PC (PC + 2) into signed 21 bits using its 13-bit immediate data *imm13*. The *sign21* allows branches within the range of PC - 1,048,574 to PC + 1,048,576.

(3) Extension 2

```
ext    imm3    ; imm3(2:0)= sign24(23:21)
ext    imm13   ; = sign24(20:8)
jreq   sign7   ; = "jreq sign24", sign7 = sign24(7:1), sign24(0)=0
```

The `ext` instructions extend the displacement to be added to the PC (PC + 2) into signed 24 bits using their immediates (*imm3* and *imm13*). The *sign24* allows branches within the range of PC - 8,388,606 to PC + 8,388,608.

(4) Delayed branch (d bit (bit 7) = 1)

```
jreq.d  sign7
```

For the `jreq.d` instruction, the next instruction becomes a delayed slot instruction. A delayed slot instruction is executed before the program branches. Interrupts are masked in intervals between the `jreq.d` instruction and the next instruction, so no interrupts occur.

Example

```
cmp   %r0,%r1
jreq  0x1        ; Skips the next instruction if r1 = r0.
```

Caution

When the `jreq.d` instruction (delayed branch) is used, be careful to ensure that the next instruction is limited to those that can be used as a delayed slot instruction. If any other instruction is executed, the program may operate indeterminately. For the usable instructions, refer to the instruction list in the Appendix.

jrge *sign7*
jrge.d *sign7*

Function

Conditional PC relative jump (for judgment of signed operation results)

Standard) pc ← pc + 2 + *sign7* × 2 if !(N^V) is true
Extension 1) pc ← pc + 2 + *sign21* if !(N^V) is true
Extension 2) pc ← pc + 2 + *sign24* if !(N^V) is true

Code

15	14	13	12	11	10	9	8	7	6–0	
0	0	0	0	0	1	1	1	0	*sign7*	jrge
0	0	0	0	0	1	1	1	1	*sign7*	jrge.d

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Signed PC relative

CLK

`jrge` Two cycles (when not branched), Three cycles (when branched)
`jrge.d` Two cycles (when a one-cycle delayed slot instruction follows), Three cycles (other)

Description

(1) Standard

```
jrge  sign7   ; = "jrge sign8", sign7 = sign8(7:1), sign8(0)=0
```

If the condition below has been met, this instruction doubles the signed 7-bit immediate *sign7* and adds it to the PC (PC + 2) for branching the program flow to the address. It does not branch if the condition has not been met.

- N flag = V flag (e.g. "A ≥ B" has resulted by `cmp A,B`)

The *sign7* specifies a word address in 16-bit units.
The *sign7* (×2) allows branches within the range of PC - 126 to PC + 128.

(2) Extension 1

```
ext   imm13   ; = sign21(20:8)
jrge  sign7   ; = "jrge sign21", sign7 = sign21(7:1), sign21(0)=0
```

The `ext` instruction extends the displacement to be added to the PC (PC + 2) into signed 21 bits using its 13-bit immediate data *imm13*. The *sign21* allows branches within the range of PC - 1,048,574 to PC + 1,048,576.

(3) Extension 2

```
ext   imm3    ; imm3(2:0)= sign24(23:21)
ext   imm13   ; = sign24(20:8)
jrge  sign7   ; = "jrge sign24", sign7 = sign24(7:1), sign24(0)=0
```

The `ext` instructions extend the displacement to be added to the PC (PC + 2) into signed 24 bits using their immediates (*imm3* and *imm13*). The *sign24* allows branches within the range of PC - 8,388,606 to PC + 8,388,608.

(4) Delayed branch (d bit (bit 7) = 1)

```
jrge.d  sign7
```

For the `jrge.d` instruction, the next instruction becomes a delayed slot instruction. A delayed slot instruction is executed before the program branches. Interrupts are masked in intervals between the `jrge.d` instruction and the next instruction, so no interrupts occur.

Example

```
cmp   %r0,%r1    ; r0 and r1 contain signed data.
jrge  0x1        ; Skips the next instruction if r0 ≥ r1.
```

Caution

When the `jrge.d` instruction (delayed branch) is used, be careful to ensure that the next instruction is limited to those that can be used as a delayed slot instruction. If any other instruction is executed, the program may operate indeterminately. For the usable instructions, refer to the instruction list in the Appendix.

jrgt sign7
jrgt.d sign7

Function

Conditional PC relative jump (for judgment of signed operation results)

Standard) pc ← pc + 2 + *sign7* × 2 if !Z&!(N^V) is true
Extension 1) pc ← pc + 2 + *sign21* if !Z&!(N^V) is true
Extension 2) pc ← pc + 2 + *sign24* if !Z&!(N^V) is true

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	0	0	0	1	1	0	0	*sign7*							jrgt
0	0	0	0	0	1	1	0	1	*sign7*							jrgt.d

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Signed PC relative

CLK

`jrgt` Two cycles (when not branched), Three cycles (when branched)
`jrgt.d` Two cycles (when a one-cycle delayed slot instruction follows), Three cycles (other)

Description

(1) Standard

```
jrgt  sign7   ; = "jrgt sign8", sign7 = sign8(7:1), sign8(0)=0
```

If the condition below has been met, this instruction doubles the signed 7-bit immediate *sign7* and adds it to the PC (PC + 2) for branching the program flow to the address. It does not branch if the condition has not been met.

• Z flag = 0 and N flag = V flag (e.g. "A > B" has resulted by `cmp A,B`)

The *sign7* specifies a word address in 16-bit units.

The *sign7* (×2) allows branches within the range of PC - 126 to PC + 128.

(2) Extension 1

```
ext   imm13   ; = sign21(20:8)
jrgt  sign7   ; = "jrgt sign21", sign7 = sign21(7:1), sign21(0)=0
```

The `ext` instruction extends the displacement to be added to the PC (PC + 2) into signed 21 bits using its 13-bit immediate data *imm13*. The *sign21* allows branches within the range of PC - 1,048,574 to PC + 1,048,576.

(3) Extension 2

```
ext   imm3    ; imm3(2:0)= sign24(23:21)
ext   imm13   ; = sign24(20:8)
jrgt  sign7   ; = "jrgt sign24", sign7 = sign24(7:1), sign24(0)=0
```

The `ext` instructions extend the displacement to be added to the PC (PC + 2) into signed 24 bits using their immediates (*imm3* and *imm13*). The *sign24* allows branches within the range of PC - 8,388,606 to PC + 8,388,608.

(4) Delayed branch (d bit (bit 7) = 1)

```
jrgt.d  sign7
```

For the `jrgt.d` instruction, the next instruction becomes a delayed slot instruction. A delayed slot instruction is executed before the program branches. Interrupts are masked in intervals between the `jrgt.d` instruction and the next instruction, so no interrupts occur.

Example

```
cmp   %r0,%r1    ; r0 and r1 contain signed data.
jrgt  0x1        ; Skips the next instruction if r0 > r1.
```

Caution

When the `jrgt.d` instruction (delayed branch) is used, be careful to ensure that the next instruction is limited to those that can be used as a delayed slot instruction. If any other instruction is executed, the program may operate indeterminately. For the usable instructions, refer to the instruction list in the Appendix.

jrle *sign7*
jrle.d *sign7*

Function

Conditional PC relative jump (for judgment of signed operation results)

Standard) pc ← pc + 2 + *sign7* × 2 if Z | (N^V) is true
Extension 1) pc ← pc + 2 + *sign21* if Z | (N^V) is true
Extension 2) pc ← pc + 2 + *sign24* if Z | (N^V) is true

Code

15	14	13	12	11	10	9	8	7	6–0	
0	0	0	0	1	0	0	1	0	*sign7*	jrle
0	0	0	0	1	0	0	1	1	*sign7*	jrle.d

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Signed PC relative

CLK

`jrle` Two cycles (when not branched), Three cycles (when branched)
`jrle.d` Two cycles (when a one-cycle delayed slot instruction follows), Three cycles (other)

Description

(1) Standard

```
jrle  sign7   ; = "jrle sign8", sign7 = sign8(7:1), sign8(0)=0
```

If the condition below has been met, this instruction doubles the signed 7-bit immediate *sign7* and adds it to the PC (PC + 2) for branching the program flow to the address. It does not branch if the condition has not been met.

• Z flag = 1 or N flag ≠ V flag (e.g. "A ≤ B" has resulted by `cmp A,B`)

The *sign7* specifies a word address in 16-bit units.

The *sign7* (×2) allows branches within the range of PC - 126 to PC + 128.

(2) Extension 1

```
ext   imm13   ; = sign21(20:8)
jrle  sign7   ; = "jrle sign21", sign7 = sign21(7:1), sign21(0)=0
```

The `ext` instruction extends the displacement to be added to the PC (PC + 2) into signed 21 bits using its 13-bit immediate data *imm13*. The *sign21* allows branches within the range of PC - 1,048,574 to PC + 1,048,576.

(3) Extension 2

```
ext   imm3    ; imm3(2:0)= sign24(23:21)
ext   imm13   ; = sign24(20:8)
jrle  sign7   ; = "jrle sign24", sign7 = sign24(7:1), sign24(0)=0
```

The `ext` instructions extend the displacement to be added to the PC (PC + 2) into signed 24 bits using their immediates (*imm3* and *imm13*). The *sign24* allows branches within the range of PC - 8,388,606 to PC + 8,388,608.

(4) Delayed branch (d bit (bit 7) = 1)

```
jrle.d  sign7
```

For the `jrle.d` instruction, the next instruction becomes a delayed slot instruction. A delayed slot instruction is executed before the program branches. Interrupts are masked in intervals between the `jrle.d` instruction and the next instruction, so no interrupts occur.

Example

```
cmp   %r0,%r1   ; r0 and r1 contain signed data.
jrle  0x1       ; Skips the next instruction if r0 ≤ r1.
```

Caution

When the `jrle.d` instruction (delayed branch) is used, be careful to ensure that the next instruction is limited to those that can be used as a delayed slot instruction. If any other instruction is executed, the program may operate indeterminately. For the usable instructions, refer to the instruction list in the Appendix.

jrlt *sign7*
jrlt.d *sign7*

Function

Conditional PC relative jump (for judgment of signed operation results)

Standard) pc ← pc + 2 + *sign7* × 2 if N^V is true

Extension 1) pc ← pc + 2 + *sign21* if N^V is true

Extension 2) pc ← pc + 2 + *sign24* if N^V is true

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	0	0	1	0	0	0	0	*sign7*							`jrlt`
0	0	0	0	1	0	0	0	1	*sign7*							`jrlt.d`

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Signed PC relative

CLK

`jrlt` Two cycles (when not branched), Three cycles (when branched)

`jrlt.d` Two cycles (when a one-cycle delayed slot instruction follows), Three cycles (other)

Description

(1) Standard

```
jrlt  sign7   ; = "jrlt sign8", sign7 = sign8(7:1), sign8(0)=0
```

If the condition below has been met, this instruction doubles the signed 7-bit immediate *sign7* and adds it to the PC (PC + 2) for branching the program flow to the address. It does not branch if the condition has not been met.

• N flag ≠ V flag (e.g. “A < B” has resulted by `cmp A,B`)

The *sign7* specifies a word address in 16-bit units.

The *sign7* (×2) allows branches within the range of PC - 126 to PC + 128.

(2) Extension 1

```
ext   imm13   ; = sign21(20:8)
jrlt  sign7   ; = "jrlt sign21", sign7 = sign21(7:1), sign21(0)=0
```

The `ext` instruction extends the displacement to be added to the PC (PC + 2) into signed 21 bits using its 13-bit immediate data *imm13*. The *sign21* allows branches within the range of PC - 1,048,574 to PC + 1,048,576.

(3) Extension 2

```
ext   imm3    ; imm3(2:0)= sign24(23:21)
ext   imm13   ; = sign24(20:8)
jrlt  sign7   ; = "jrlt sign24", sign7 = sign24(7:1), sign24(0)=0
```

The `ext` instructions extend the displacement to be added to the PC (PC + 2) into signed 24 bits using their immediates (*imm3* and *imm13*). The *sign24* allows branches within the range of PC - 8,388,606 to PC + 8,388,608.

(4) Delayed branch (d bit (bit 7) = 1)

```
jrlt.d  sign7
```

For the `jrlt.d` instruction, the next instruction becomes a delayed slot instruction. A delayed slot instruction is executed before the program branches. Interrupts are masked in intervals between the `jrlt.d` instruction and the next instruction, so no interrupts occur.

Example

```
cmp   %r0,%r1    ; r0 and r1 contain signed data.
jrlt  0x1        ; Skips the next instruction if r0 < r1.
```

Caution

When the `jrlt.d` instruction (delayed branch) is used, be careful to ensure that the next instruction is limited to those that can be used as a delayed slot instruction. If any other instruction is executed, the program may operate indeterminately. For the usable instructions, refer to the instruction list in the Appendix.

jrne *sign7*
jrne.d *sign7*

Function

Conditional PC relative jump

Standard) pc ← pc + 2 + *sign7* × 2 if !Z is true
Extension 1) pc ← pc + 2 + *sign21* if !Z is true
Extension 2) pc ← pc + 2 + *sign24* if !Z is true

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	0	0	1	1	1	1	0	*sign7*							jrne
0	0	0	0	1	1	1	1	1	*sign7*							jrne.d

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Signed PC relative

CLK

`jrne` Two cycles (when not branched), Three cycles (when branched)
`jrne.d` Two cycles (when a one-cycle delayed slot instruction follows), Three cycles (other)

Description

(1) Standard

```
jrne  sign7   ; = "jrne sign8", sign7 = sign8(7:1), sign8(0)=0
```

If the condition below has been met, this instruction doubles the signed 7-bit immediate *sign7* and adds it to the PC (PC + 2) for branching the program flow to the address. It does not branch if the condition has not been met.

• Z flag = 0 (e.g. "A ≠ B" has resulted by `cmp A,B`)

The *sign7* specifies a word address in 16-bit units.
The *sign7* (×2) allows branches within the range of PC - 126 to PC + 128.

(2) Extension 1

```
ext   imm13   ; = sign21(20:8)
jrne  sign7   ; = "jrne sign21", sign7 = sign21(7:1), sign21(0)=0
```

The `ext` instruction extends the displacement to be added to the PC (PC + 2) into signed 21 bits using its 13-bit immediate data *imm13*. The *sign21* allows branches within the range of PC - 1,048,574 to PC + 1,048,576.

(3) Extension 2

```
ext   imm3    ; imm3(2:0)= sign24(23:21)
ext   imm13   ; = sign24(20:8)
jrne  sign7   ; = "jrne sign24", sign7 = sign24(7:1), sign24(0)=0
```

The `ext` instructions extend the displacement to be added to the PC (PC + 2) into signed 24 bits using their immediates (*imm3* and *imm13*). The *sign24* allows branches within the range of PC - 8,388,606 to PC + 8,388,608.

(4) Delayed branch (d bit (bit 7) = 1)

```
jrne.d  sign7
```

For the `jrne.d` instruction, the next instruction becomes a delayed slot instruction. A delayed slot instruction is executed before the program branches. Interrupts are masked in intervals between the `jrne.d` instruction and the next instruction, so no interrupts occur.

Example

```
cmp   %r0,%r1
jrne  0x1        ; Skips the next instruction if r0 ≠ r1.
```

Caution

When the `jrne.d` instruction (delayed branch) is used, be careful to ensure that the next instruction is limited to those that can be used as a delayed slot instruction. If any other instruction is executed, the program may operate indeterminately. For the usable instructions, refer to the instruction list in the Appendix.

jruge *sign7*
jruge.d *sign7*

Function

Conditional PC relative jump (for judgment of unsigned operation results)

Standard) pc ← pc + 2 + *sign7* × 2 if !C is true
Extension 1) pc ← pc + 2 + *sign21* if !C is true
Extension 2) pc ← pc + 2 + *sign24* if !C is true

Code

15	14	13	12	11	10	9	8	7	6–0	
0	0	0	0	1	0	1	1	0	*sign7*	`jruge`
0	0	0	0	1	0	1	1	1	*sign7*	`jruge.d`

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Signed PC relative

CLK

`jruge` Two cycles (when not branched), Three cycles (when branched)
`jruge.d` Two cycles (when a one-cycle delayed slot instruction follows), Three cycles (other)

Description

(1) Standard

```
jruge  sign7  ; = "jruge sign8", sign7 = sign8(7:1), sign8(0)=0
```

If the condition below has been met, this instruction doubles the signed 7-bit immediate *sign7* and adds it to the PC (PC + 2) for branching the program flow to the address. It does not branch if the condition has not been met.

• C flag = 0 (e.g. "A ≥ B" has resulted by `cmp A,B`)

The *sign7* specifies a word address in 16-bit units.
The *sign7* (×2) allows branches within the range of PC - 126 to PC + 128.

(2) Extension 1

```
ext    imm13  ; = sign21(20:8)
jruge  sign7  ; = "jruge sign21", sign7 = sign21(7:1), sign21(0)=0
```

The `ext` instruction extends the displacement to be added to the PC (PC + 2) into signed 21 bits using its 13-bit immediate data *imm13*. The *sign21* allows branches within the range of PC - 1,048,574 to PC + 1,048,576.

(3) Extension 2

```
ext    imm3   ; imm3(2:0)= sign24(23:21)
ext    imm13  ; = sign24(20:8)
jruge  sign7  ; = "jruge sign24", sign7 = sign24(7:1), sign24(0)=0
```

The `ext` instructions extend the displacement to be added to the PC (PC + 2) into signed 24 bits using their immediates (*imm3* and *imm13*). The *sign24* allows branches within the range of PC - 8,388,606 to PC + 8,388,608.

(4) Delayed branch (d bit (bit 7) = 1)

```
jruge.d  sign7
```

For the `jruge.d` instruction, the next instruction becomes a delayed slot instruction. A delayed slot instruction is executed before the program branches. Interrupts are masked in intervals between the `jruge.d` instruction and the next instruction, so no interrupts occur.

Example

```
cmp    %r0,%r1   ; r0 and r1 contain unsigned data.
jruge  0x1       ; Skips the next instruction if r0 ≥ r1.
```

Caution

When the `jruge.d` instruction (delayed branch) is used, be careful to ensure that the next instruction is limited to those that can be used as a delayed slot instruction. If any other instruction is executed, the program may operate indeterminately. For the usable instructions, refer to the instruction list in the Appendix.

jrugt *sign7*
jrugt.d *sign7*

Function

Conditional PC relative jump (for judgment of unsigned operation results)

Standard) pc ← pc + 2 + *sign7* × 2 if !Z&!C is true
Extension 1) pc ← pc + 2 + *sign21* if !Z&!C is true
Extension 2) pc ← pc + 2 + *sign24* if !Z&!C is true

Code

15	14	13	12	11	10	9	8	7	6–0	
0	0	0	0	1	0	1	0	0	*sign7*	jrugt
0	0	0	0	1	0	1	0	1	*sign7*	jrugt.d

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Signed PC relative

CLK

`jrugt` Two cycles (when not branched), Three cycles (when branched)
`jrugt.d` Two cycles (when a one-cycle delayed slot instruction follows), Three cycles (other)

Description

(1) Standard

```
jrugt  sign7  ; = "jrugt sign8", sign7 = sign8(7:1), sign8(0)=0
```

If the condition below has been met, this instruction doubles the signed 7-bit immediate *sign7* and adds it to the PC (PC + 2) for branching the program flow to the address. It does not branch if the condition has not been met.

• Z flag = 0 and C flag = 0 (e.g. "A > B" has resulted by `cmp A,B`)

The *sign7* specifies a word address in 16-bit units.

The *sign7* (×2) allows branches within the range of PC - 126 to PC + 128.

(2) Extension 1

```
ext    imm13  ; = sign21(20:8)
jrugt  sign7  ; = "jrugt sign21", sign7 = sign21(7:1), sign21(0)=0
```

The `ext` instruction extends the displacement to be added to the PC (PC + 2) into signed 21 bits using its 13-bit immediate data *imm13*. The *sign21* allows branches within the range of PC - 1,048,574 to PC + 1,048,576.

(3) Extension 2

```
ext    imm3   ; imm3(2:0)= sign24(23:21)
ext    imm13  ; = sign24(20:8)
jrugt  sign7  ; = "jrugt sign24", sign7 = sign24(7:1), sign24(0)=0
```

The `ext` instructions extend the displacement to be added to the PC (PC + 2) into signed 24 bits using their immediates (*imm3* and *imm13*). The *sign24* allows branches within the range of PC - 8,388,606 to PC + 8,388,608.

(4) Delayed branch (d bit (bit 7) = 1)

```
jrugt.d  sign7
```

For the `jrugt.d` instruction, the next instruction becomes a delayed slot instruction. A delayed slot instruction is executed before the program branches. Interrupts are masked in intervals between the `jrugt.d` instruction and the next instruction, so no interrupts occur.

Example

```
cmp    %r0,%r1   ; r0 and r1 contain unsigned data.
jrugt  0x1       ; Skips the next instruction if r0 > r1.
```

Caution

When the `jrugt.d` instruction (delayed branch) is used, be careful to ensure that the next instruction is limited to those that can be used as a delayed slot instruction. If any other instruction is executed, the program may operate indeterminately. For the usable instructions, refer to the instruction list in the Appendix.

jrule *sign7*
jrule.d *sign7*

Function Conditional PC relative jump (for judgment of unsigned operation results)

Standard) pc ← pc + 2 + *sign7* × 2 if Z | C is true
Extension 1) pc ← pc + 2 + *sign21* if Z | C is true
Extension 2) pc ← pc + 2 + *sign24* if Z | C is true

Code

15	14	13	12	11	10	9	8	7	6–0	
0	0	0	0	1	1	0	1	0	*sign7*	`jrule`
0	0	0	0	1	1	0	1	1	*sign7*	`jrule.d`

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode Signed PC relative

CLK
`jrule` Two cycles (when not branched), Three cycles (when branched)
`jrule.d` Two cycles (when a one-cycle delayed slot instruction follows), Three cycles (other)

Description (1) Standard

```
jrule  sign7  ; = "jrule sign8", sign7 = sign8(7:1), sign8(0)=0
```

If the condition below has been met, this instruction doubles the signed 7-bit immediate *sign7* and adds it to the PC (PC + 2) for branching the program flow to the address. It does not branch if the condition has not been met.
• Z flag = 1 or C flag = 1 (e.g. "A ≤ B" has resulted by `cmp A,B`)
The *sign7* specifies a word address in 16-bit units.
The *sign7* (×2) allows branches within the range of PC - 126 to PC + 128.

(2) Extension 1

```
ext    imm13  ; = sign21(20:8)
jrule  sign7  ; = "jrule sign21", sign7 = sign21(7:1), sign21(0)=0
```

The `ext` instruction extends the displacement to be added to the PC (PC + 2) into signed 21 bits using its 13-bit immediate data *imm13*. The *sign21* allows branches within the range of PC - 1,048,574 to PC + 1,048,576.

(3) Extension 2

```
ext    imm3   ; imm3(2:0)= sign24(23:21)
ext    imm13  ; = sign24(20:8)
jrule  sign7  ; = "jrule sign24", sign7 = sign24(7:1), sign24(0)=0
```

The `ext` instructions extend the displacement to be added to the PC (PC + 2) into signed 24 bits using their immediates (*imm3* and *imm13*). The *sign24* allows branches within the range of PC - 8,388,606 to PC + 8,388,608.

(4) Delayed branch (d bit (bit 7) = 1)

```
jrule.d  sign7
```

For the `jrule.d` instruction, the next instruction becomes a delayed slot instruction. A delayed slot instruction is executed before the program branches. Interrupts are masked in intervals between the `jrule.d` instruction and the next instruction, so no interrupts occur.

Example

```
cmp    %r0,%r1   ; r0 and r1 contain unsigned data.
jrule  0x1       ; Skips the next instruction if r0 ≤ r1.
```

Caution When the `jrule.d` instruction (delayed branch) is used, be careful to ensure that the next instruction is limited to those that can be used as a delayed slot instruction. If any other instruction is executed, the program may operate indeterminately. For the usable instructions, refer to the instruction list in the Appendix.

jrult sign7
jrult.d sign7

Function

Conditional PC relative jump (for judgment of unsigned operation results)

Standard) pc ← pc + 2 + *sign7* × 2 if C is true
Extension 1) pc ← pc + 2 + *sign21* if C is true
Extension 2) pc ← pc + 2 + *sign24* if C is true

Code

15	14	13	12	11	10	9	8	7	6–0	
0	0	0	0	1	1	0	0	0	*sign7*	jrult
0	0	0	0	1	1	0	0	1	*sign7*	jrult.d

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Signed PC relative

CLK

`jrult` Two cycles (when not branched), Three cycles (when branched)
`jrult.d` Two cycles (when a one-cycle delayed slot instruction follows), Three cycles (other)

Description

(1) Standard

```
jrult  sign7  ; = "jrult sign8", sign7 = sign8(7:1), sign8(0)=0
```

If the condition below has been met, this instruction doubles the signed 7-bit immediate *sign7* and adds it to the PC (PC + 2) for branching the program flow to the address. It does not branch if the condition has not been met.

• C flag = 1 (e.g. “A < B” has resulted by `cmp A,B`)

The *sign7* specifies a word address in 16-bit units.

The *sign7* (×2) allows branches within the range of PC - 126 to PC + 128.

(2) Extension 1

```
ext    imm13  ; = sign21(20:8)
jrult  sign7  ; = "jrult sign21", sign7 = sign21(7:1), sign21(0)=0
```

The `ext` instruction extends the displacement to be added to the PC (PC + 2) into signed 21 bits using its 13-bit immediate data *imm13*. The *sign21* allows branches within the range of PC - 1,048,574 to PC + 1,048,576.

(3) Extension 2

```
ext    imm3   ; imm3(2:0)= sign24(23:21)
ext    imm13  ; = sign24(20:8)
jrult  sign7  ; = "jrult sign24", sign7 = sign24(7:1), sign24(0)=0
```

The `ext` instructions extend the displacement to be added to the PC (PC + 2) into signed 24 bits using their their immediates (*imm3* and *imm13*). The *sign24* allows branches within the range of PC - 8,388,606 to PC + 8,388,608.

(4) Delayed branch (d bit (bit 7) = 1)

```
jrult.d  sign7
```

For the `jrult.d` instruction, the next instruction becomes a delayed slot instruction. A delayed slot instruction is executed before the program branches. Interrupts are masked in intervals between the `jrult.d` instruction and the next instruction, so no interrupts occur.

Example

```
cmp    %r0,%r1   ; r0 and r1 contain unsigned data.
jrult  0x1       ; Skips the next instruction if r0 < r1.
```

Caution

When the `jrult.d` instruction (delayed branch) is used, be careful to ensure that the next instruction is limited to those that can be used as a delayed slot instruction. If any other instruction is executed, the program may operate indeterminately. For the usable instructions, refer to the instruction list in the Appendix.

ld %rd, %rs

Function

16-bit data transfer

Standard) *rd*(15:0) ← *rs*(15:0), *rd*(23:16) ← 0

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	1	0		*rd*		0	0	1	0		*rs*	

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

The 16 low-order bits of the *rs* register are transferred to the *rd* register. The eight high-order bits of the *rd* register are set to 0.

(2) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit.

Example

```
ld  %r0,%r1     ; r0 ← r1(15:0)
```

ld %rd, [%rb]

Function

16-bit data transfer

Standard) *rd*(15:0) ← W[*rb*], *rd*(23:16) ← 0

Extension 1) *rd*(15:0) ← W[*rb* + *imm13*], *rd*(23:16) ← 0

Extension 2) *rd*(15:0) ← W[*rb* + *imm24*], *rd*(23:16) ← 0

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	0	0		rd		0	0	1	0		rb	

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register indirect `%rb` = `%r0` to `%r7`

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle (two cycles when the `ext` instruction is used)

Description

(1) Standard

```
ld   %rd,[%rb]        ; memory address = rb
```

The 16-bit data in the specified memory location is transferred to the *rd* register. The *rb* register contains the memory address to be accessed. The eight high-order bits of the *rd* register are set to 0.

(2) Extension 1

```
ext   imm13
ld    %rd,[%rb]       ; memory address = rb + imm13
```

The `ext` instruction changes the addressing mode to register indirect addressing with displacement. As a result, the content of the *rb* register with the 13-bit immediate *imm13* added comprises the memory address, the 16-bit data in which is transferred to the *rd* register. The content of the *rb* register is not altered. The eight high-order bits of the *rd* register are set to 0.

(3) Extension 2

```
ext   imm11           ; imm11(10:0) = imm24(23:13)
ext   imm13           ; = imm24(12:0)
ld    %rd,[%rb]       ; memory address = rb + imm24
```

The addressing mode changes to register indirect addressing with displacement, so the content of the *rb* register with the 24-bit immediate *imm24* added comprises the memory address, the 16-bit data in which is transferred to the *rd* register. The content of the *rb* register is not altered. The eight high-order bits of the *rd* register are set to 0.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Caution

The *rb* register and the displacement must specify a 16-bit boundary address (least significant bit = 0). Specifying an odd address causes an address misaligned interrupt. Note, however, that the data transfer is performed by setting the least significant bit of the address to 0.

ld *%rd*, [*%rb*]+
ld *%rd*, [*%rb*]-
ld *%rd*, -[*%rb*]

Function

16-bit data transfer with address increment/decrement option

ld *%rd*, [*%rb*]+ (with post-increment option)

Standard) *rd*(15:0) ← W[*rb*], *rd*(23:16) ← 0, *rb*(23:0) ← *rb*(23:0) + 2
Extension 1) *rd*(15:0) ← W[*rb*], *rd*(23:16) ← 0, *rb*(23:0) ← *rb*(23:0) + *imm13*
Extension 2) *rd*(15:0) ← W[*rb*], *rd*(23:16) ← 0, *rb*(23:0) ← *rb*(23:0) + *imm24*

ld *%rd*, [*%rb*]- (with post-decrement option)

Standard) *rd*(15:0) ← W[*rb*], *rd*(23:16) ← 0, *rb*(23:0) ← *rb*(23:0) - 2
Extension 1) *rd*(15:0) ← W[*rb*], *rd*(23:16) ← 0, *rb*(23:0) ← *rb*(23:0) - *imm13*
Extension 2) *rd*(15:0) ← W[*rb*], *rd*(23:16) ← 0, *rb*(23:0) ← *rb*(23:0) - *imm24*

ld *%rd*, -[*%rb*] (with pre-decrement option)

Standard) *rb*(23:0) ← *rb*(23:0) - 2, *rd*(15:0) ← W[*rb*], *rd*(23:16) ← 0
Extension 1) *rb*(23:0) ← *rb*(23:0) - *imm13, rd*(15:0) ← W[*rb*], *rd*(23:16) ← 0
Extension 2) *rb*(23:0) ← *rb*(23:0) - *imm24, rd*(15:0) ← W[*rb*], *rd*(23:16) ← 0

Code

15	14	13	12	11	10	9–7	6	5	4	3	2–0	
0	0	1	0	0	0	*rd*	0	1	1	0	*rb*	`ld %rd, [%rb]+`
0	0	1	0	0	0	*rd*	1	1	1	0	*rb*	`ld %rd, [%rb]-`
0	0	1	0	0	0	*rd*	1	0	1	0	*rb*	`ld %rd, -[%rb]`

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register indirect `%rb` = `%r0` to `%r7`
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

Two cycles

Description

(1) Address increment/decrement option

Specifying the `[]+`, `[]-`, or `-[]` option will automatically increment/decrement the memory address. This allows the program to simply perform continuous data transfer.

`ld %rd, [%rb]+` Load instruction with post-increment option
The memory address will be incremented after the data transfer has finished.

`ld %rd, [%rb]-` Load instruction with post-decrement option
The memory address will be decremented after the data transfer has finished.

`ld %rd, -[%rb]` Load instruction with pre-decrement option
The memory address will be decremented before starting the data transfer.

The address increment/decrement sizes are listed below.
When no `ext` is used (standard): 2 (16-bit size)
When one `ext` is used (extension 1): *imm13*
When two `ext` are used (extension 2): *imm24*

(2) Standard (example of post-increment option)

```
ld   %rd,[%rb]-   ; source memory address = rb
                  ; post decrement: rb - imm13
```

The 16-bit data in the specified memory location is transferred to the *rd* register. The *rb* register contains the memory address to be accessed. The eight high-order bits of the *rd* register are set to 0. The memory address will be incremented by two bytes after the data transfer has finished.

(3) Extension 1 (example of post-decrement option)

```
ext  imm13
ld   %rd,[%rb]-   ; source memory address = rb
                  ; post decrement: rb - imm13
```

The 16-bit data in the specified memory location is transferred to the *rd* register. The *rb* register contains the memory address to be accessed. The eight high-order bits of the *rd* register are set to 0. The memory address will be decremented by *imm13* bytes after the data transfer has finished.

(4) Extension 2 (example of pre-decrement option)

```
ext  imm11        ; imm11(10:0) = imm24(23:13)
ext  imm13        ; = imm24(12:0)
ld   %rd,-[%rb]   ; source memory address = rb - imm24
```

After the memory address specified by the *rb* register is decremented by *imm24* bytes, the 16-bit data in the decremented address is transferred to the *rd* register. The eight high-order bits of the rd register are set to 0.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Caution The *rb* register and the immediate value must specify a 16-bit boundary address (least significant bit = 0). Specifying an odd address causes an address misaligned interrupt. Note, however, that the data transfer is performed by setting the least significant bit of the address to 0.

ld %rd, [%sp + imm7]

Function

16-bit data transfer

Standard) $rd(15:0) \leftarrow W[sp + imm7], rd(23:16) \leftarrow 0$

Extension 1) $rd(15:0) \leftarrow W[sp + imm20], rd(23:16) \leftarrow 0$

Extension 2) $rd(15:0) \leftarrow W[sp + imm24], rd(23:16) \leftarrow 0$

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	0	1	0	*rd*			*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register indirect with displacement

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

Two cycles

Description

(1) Standard

```
ld   %rd,[%sp + imm7]        ; memory address = sp + imm7
```

The 16-bit data in the specified memory location is transferred to the *rd* register. The content of the current SP with the 7-bit immediate *imm7* added as displacement comprises the memory address to be accessed. The eight high-order bits of the *rd* register are set to 0.

(2) Extension 1

```
ext  imm13                   ; = imm20(19:7)
ld   %rd,[%sp + imm7]        ; memory address = sp + imm20,
                             ; imm7 = imm20(6:0)
```

The `ext` instruction extends the displacement to a 20-bit quantity. As a result, the content of the SP with the 20-bit immediate *imm20* added comprises the memory address, the 16-bit data in which is transferred to the *rd* register. The eight high-order bits of the *rd* register are set to 0.

(3) Extension 2

```
ext  imm4                    ; imm4(3:0) = imm24(23:20)
ext  imm13                   ; = imm24(19:7)
ld   %rd,[%sp + imm7]        ; memory address = sp + imm24,
                             ; imm7 = imm24(6:0)
```

The two `ext` instructions extend the displacement to a 24-bit quantity. As a result, the content of the SP with the 24-bit immediate *imm24* added comprises the memory address, the 16-bit data in which is transferred to the *rd* register. The eight high-order bits of the *rd* register are set to 0.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

```
ext  0x1
ld   %r0,[%sp + 0x2]    ; r0 ← [sp + 0x82]
```

Caution

The SP and the displacement must specify a 16-bit boundary address (least significant bit = 0). Specifying an odd address causes an address misaligned interrupt. Note, however, that the data transfer is performed by setting the least significant bit of the address to 0.

ld %rd, [imm7]

Function

16-bit data transfer

Standard) *rd*(15:0) ← W[*imm7*], *rd*(23:16) ← 0
Extension 1) *rd*(15:0) ← W[*imm20*], *rd*(23:16) ← 0
Extension 2) *rd*(15:0) ← W[*imm24*], *rd*(23:16) ← 0

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	0	0	1	0	*rd*			*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Immediate data (unsigned)
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
ld   %rd,[imm7]        ; memory address = imm7
```

The 16-bit data in the memory address specified with the 7-bit immediate *imm7* is transferred to the *rd* register. The eight high-order bits of the *rd* register are set to 0.

(2) Extension 1

```
ext   imm13            ; = imm20(19:7)
ld    %rd,[imm7]       ; memory address = imm20,
                       ; imm7 = imm20(6:0)
```

The `ext` instruction extends the memory address to a 20-bit quantity. As a result, the 16-bit data in the memory address specified with the 20-bit immediate *imm20* is transferred to the *rd* register. The eight high-order bits of the *rd* register are set to 0.

(3) Extension 2

```
ext   imm4             ; imm4(3:0) = imm24(23:20)
ext   imm13            ; = imm24(19:7)
ld    %rd,[imm7]       ; memory address = sp + imm24,
                       ; imm7 = imm24(6:0)
```

The two `ext` instructions extend the memory address to a 24-bit quantity. As a result, the 16-bit data in the memory address specified with the 24-bit immediate *imm24* is transferred to the *rd* register. The eight high-order bits of the *rd* register are set to 0.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

```
ext   0x1
ld    %r0,[0x2]   ; r0 ← [0x82]
```

Caution

The *imm7* must specify a 16-bit boundary address (least significant bit = 0). Specifying an odd address causes an address misaligned interrupt. Note, however, that the data transfer is performed by setting the least significant bit of the address to 0.

ld %rd, sign7

Function

16-bit data transfer

Standard) *rd*(6:0) ← *sign7*(6:0), *rd*(15:7) ← *sign7*(6), *rd*(23:16) ← 0

Extension 1) *rd*(15:0) ← *sign16*(15:0), *rd*(23:16) ← 0

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	1	1	0	*rd*			*sign7*						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Immediate data (signed)

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
ld   %rd,sign7        ; rd ← sign7 (sign-extended)
```

The 7-bit immediate *sign7* is loaded to the *rd* register after being sign-extended to a 16-bit quantity.

(2) Extension 1

```
ext   imm13           ; = sign16(15:7)
ld    %rd,sign7       ; rd ← sign16, sign7 = sign16(6:0)
```

The immediate data is extended into a 16-bit quantity by the `ext` instruction and it is loaded to the *rd* register.

(3) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

```
ld   %r0,0x7f     ; r0 ← 0xffff (r0 = 0x00ffff)
```

ld [%rb], %rs

Function

16-bit data transfer

Standard) W[*rb*] ← *rs*(15:0)

Extension 1) W[*rb* + *imm13*] ← *rs*(15:0)

Extension 2) W[*rb* + *imm24*] ← *rs*(15:0)

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	0	1		rs		0	0	1	0		rb	

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register indirect `%rb` = `%r0` to `%r7`

CLK

One cycle (two cycles when the `ext` instruction is used)

Description

(1) Standard

```
ld   [%rb],%rs        ; memory address = rb
```

The 16 low-order bits of the *rs* register are transferred to the specified memory location. The *rb* register contains the memory address to be accessed.

(2) Extension 1

```
ext   imm13
ld    [%rb],%rs       ; memory address = rb + imm13
```

The `ext` instruction changes the addressing mode to register indirect addressing with displacement. As a result, the 16 low-order bits of the *rs* register are transferred to the address indicated by the content of the *rb* register with the 13-bit immediate *imm13* added. The content of the *rb* register is not altered.

(3) Extension 2

```
ext   imm11           ; imm11(10:0) = imm24(23:13)
ext   imm13           ; = imm24(12:0)
ld    [%rb],%rs       ; memory address = rb + imm24
```

The addressing mode changes to register indirect addressing with displacement, so the 16 low-order bits of the *rs* register are transferred to the address indicated by the content of the *rb* register with the 24-bit immediate *imm24* added. The content of the *rb* register is not altered.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Caution

The *rb* register and the displacement must specify a 16-bit boundary address (least significant bit = 0). Specifying an odd address causes an address misaligned interrupt. Note, however, that the data transfer is performed by setting the least significant bit of the address to 0.

ld [%*rb*]+, %*rs*
ld [%*rb*]-, %*rs*
ld -[%*rb*], %*rs*

Function 16-bit data transfer with address increment/decrement option

ld [%*rb*]+, %*rs* (with post-increment option)

Standard) W[*rb*] ← *rs*(15:0), *rb*(23:0) ← *rb*(23:0) + 2
Extension 1) W[*rb*] ← *rs*(15:0), *rb*(23:0) ← *rb*(23:0) + *imm13*
Extension 2) W[*rb*] ← *rs*(15:0), *rb*(23:0) ← *rb*(23:0) + *imm24*

ld [%*rb*]-, %*rs* (with post-decrement option)

Standard) W[*rb*] ← *rs*(15:0), *rb*(23:0) ← *rb*(23:0) - 2
Extension 1) W[*rb*] ← *rs*(15:0), *rb*(23:0) ← *rb*(23:0) - *imm13*
Extension 2) W[*rb*] ← *rs*(15:0), *rb*(23:0) ← *rb*(23:0) - *imm24*

ld -[%*rb*], %*rs* (with pre-decrement option)

Standard) *rb*(23:0) ← *rb*(23:0) - 2, W[*rb*] ← *rs*(15:0)
Extension 1) *rb*(23:0) ← *rb*(23:0) - *imm13,* W[*rb*] ← *rs*(15:0)
Extension 2) *rb*(23:0) ← *rb*(23:0) - *imm24,* W[*rb*] ← *rs*(15:0)

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	1	0	0	1		rs		0	1	1	0		rb		`ld [%rb]+,%rs`
0	0	1	0	0	1		rs		1	1	1	0		rb		`ld [%rb]-,%rs`
0	0	1	0	0	1		rs		1	0	1	0		rb		`ld -[%rb],%rs`

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode Src: Register direct `%rs` = `%r0` to `%r7`
Dst: Register indirect `%rb` = `%r0` to `%r7`

CLK Two cycles

Description (1) Address increment/decrement option

Specifying the `[]+`, `[]-`, or `-[]` option will automatically increment/decrement the memory address. This allows the program to simply perform continuous data transfer.

`ld [%rb]+,%rs` Load instruction with post-increment option
The memory address will be incremented after the data transfer has finished.

`ld [%rb]-,%rs` Load instruction with post-decrement option
The memory address will be decremented after the data transfer has finished.

`ld -[%rb],%rs` Load instruction with pre-decrement option
The memory address will be decremented before starting the data transfer.

The address increment/decrement sizes are listed below.
When no `ext` is used (standard): 2 (16-bit size)
When one `ext` is used (extension 1): *imm13*
When two `ext` are used (extension 2): *imm24*

(2) Standard (example of post-increment option)

```
ld   [%rb]+,%rs       ; Destination memory address = rb
                      ; post increment: rb + 2
```

The 16 low-order bits of the *rs* register are transferred to the specified memory location. The *rb* register contains the memory address to be accessed. The memory address will be incremented by two bytes after the data transfer has finished.

(3) Extension 1 (example of post-decrement option)

```
ext  imm13
ld   [%rb]-,%rs       ; Destination memory address = rb
                      ; post decrement: rb - imm13
```

The 16 low-order bits of the *rs* register are transferred to the specified memory location. The *rb* register contains the memory address to be accessed. The memory address will be decremented by *imm13* bytes after the data transfer has finished.

(4) Extension 2 (example of pre-decrement option)

```
ext  imm11            ; imm11(10:0) = imm24(23:13)
ext  imm13            ; = imm24(12:0)
ld  -[%rb],%rs        ; Destination memory address = rb - imm24
```

After the memory address specified by the *rb* register is decremented by *imm24* bytes, the 16 low-order bits of the *rs* register are transferred to the decremented address.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Caution The *rb* register and the immediate value must specify a 16-bit boundary address (least significant bit = 0). Specifying an odd address causes an address misaligned interrupt. Note, however, that the data transfer is performed by setting the least significant bit of the address to 0.

ld [%sp + *imm7*], %*rs*

Function
16-bit data transfer
Standard) W[sp + *imm7*] ← *rs*(15:0)
Extension 1) W[sp + *imm20*] ← *rs*(15:0)
Extension 2) W[sp + *imm24*] ← *rs*(15:0)

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	1	1	0	*rs*			*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode
Src: Register direct `%rs` = `%r0` to `%r7`
Dst: Register indirect with displacement

CLK
Two cycles

Description
(1) Standard

```
ld   [%sp + imm7],%rs       ; memory address = sp + imm7
```

The 16 low-order bits of the *rs* register are transferred to the specified memory location. The content of the current SP with the 7-bit immediate *imm7* added as displacement comprises the memory address to be accessed.

(2) Extension 1

```
ext  imm13                  ; = imm20(19:7)
ld   [%sp + imm7],%rs       ; memory address = sp + imm20,
                            ; imm7 = imm20(6:0)
```

The `ext` instruction extends the displacement to a 20-bit quantity. As a result, the 16 low-order bits of the *rs* register are transferred to the address indicated by the content of the SP with the 20-bit immediate *imm20* added.

(3) Extension 2

```
ext  imm4                   ; imm4(3:0) = imm24(23:20)
ext  imm13                  ; = imm24(19:7)
ld   [%sp + imm7],%rs       ; memory address = sp + imm24,
                            ; imm7 = imm24(6:0)
```

The two `ext` instructions extend the displacement to a 24-bit quantity. As a result, the 16 low-order bits of the *rs* register are transferred to the address indicated by the content of the SP with the 24-bit immediate *imm24* added.

(4) Delayed slot instruction
This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

```
ext  0x1
ld   [%sp + 0x2],%r0   ; W[sp + 0x82] ← 16 low-order bits of r0
```

Caution
The SP and the displacement must specify a 16-bit boundary address (least significant bit = 0). Specifying an odd address causes an address misaligned interrupt. Note, however, that the data transfer is performed by setting the least significant bit of the address to 0.

ld [*imm7*], %*rs*

Function

16-bit data transfer

Standard) W[*imm7*] ← *rs*(15:0)
Extension 1) W[*imm20*] ← *rs*(15:0)
Extension 2) W[*imm24*] ← *rs*(15:0)

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	0	1	1	0	*rs*			*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`
Dst: Immediate data (unsigned)

CLK

One cycle

Description

(1) Standard

```
ld   [imm7],%rs        ; memory address = imm7
```

The 16 low-order bits of the *rs* register are transferred to the memory address specified with the 7-bit immediate *imm7*.

(2) Extension 1

```
ext   imm13            ; = imm20(19:7)
ld    [imm7],%rs       ; memory address = imm20, imm7 = imm20(6:0)
```

The `ext` instruction extends the memory address to a 20-bit quantity. As a result, the 16 low-order bits of the *rs* register are transferred to the memory address specified with the 20-bit immediate *imm20*.

(3) Extension 2

```
ext   imm4             ; imm4(3:0) = imm24(23:20)
ext   imm13            ; = imm24(19:7)
ld    [imm7],%rs       ; memory address = imm24, imm7 = imm24(6:0)
```

The two `ext` instructions extend the memory address to a 24-bit quantity. As a result, the 16 low-order bits of the *rs* register are transferred to the memory address specified with the 24-bit immediate *imm24*.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

```
ext   0x1
ld    [0x2],%r0        ; W[0x82] ← 16 low-order bits of r0
```

Caution

The *imm7* must specify a 16-bit boundary address (least significant bit = 0). Specifying an odd address causes an address misaligned interrupt. Note, however, that the data transfer is performed by setting the least significant bit of the address to 0.

ld.a *%rd*, %pc

Function

24-bit data transfer

Standard) *rd*(23:0) ← pc(23:0) + 2
Extension 1) Unusable
Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	1	1	1		*rd*		0	1	1	0	0	0	0

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%pc`
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

The content of the PC (PC + 2) is transferred to the *rd* register.

Example

```
ld.a   %r0,%pc    ; r0 ← pc + 2
```

Caution

- When this instruction is executed, a value equal to the PC of this instruction plus 2 is loaded into the register. This instruction must be executed as a delayed slot instruction. If it does not follow a delayed branch instruction, the PC value that is loaded into the *rd* register may not be the next instruction address to the `ld.a` instruction.
- This instruction must be used as a delayed slot instruction for `jr*.d`, `jpr.d` or `jpa.d`.

ld.a *%rd*, *%rs*

Function

24-bit data transfer

Standard) *rd*(23:0) ← *rs*(23:0)
Extension 1) Unusable
Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	1	0		*rd*		0	0	1	1		*rs*	

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard
The content of the *rs* register (24-bit data) is transferred to the *rd* register.

(2) Delayed slot instruction
This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit.

Example

```
ld.a  %r0,%r1   ; r0 ← r1
```

ld.a *%rd*, %sp

Function

24-bit data transfer

Standard) *rd*(23:2) ← sp(23:2), *rd*(1:0) ← 0

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	1	1	1	*rd*			0	0	1	0	0	0	0

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%sp`

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

The content of the SP (24-bit data) is transferred to the *rd* register.

Example

```
ld.a   %r0,%sp    ; r0 ← sp
```

ld.a *%rd*, [*%rb*]

Function

32-bit data transfer

Standard) *rd*(23:0) ← A[*rb*](23:0), ignored ← A[*rb*](31:24)
Extension 1) *rd*(23:0) ← A[*rb* + *imm13*](23:0), ignored ← A[*rb* + *imm13*](31:24)
Extension 2) *rd*(23:0) ← A[*rb* + *imm24*](23:0), ignored ← A[*rb* + *imm24*](31:24)

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	0	0		*r d*		0	0	1	1		*r b*	

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register indirect `%rb` = `%r0` to `%r7`
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle (two cycles when the `ext` instruction is used)

Description

(1) Standard

```
ld.a    %rd,[%rb]       ; memory address = rb
```

The 32-bit data (the eight high-order bits are ignored) in the specified memory location is transferred to the *rd* register. The *rb* register contains the memory address to be accessed.

(2) Extension 1

```
ext     imm13
ld.a    %rd,[%rb]       ; memory address = rb + imm13
```

The `ext` instruction changes the addressing mode to register indirect addressing with displacement. As a result, the content of the *rb* register with the 13-bit immediate *imm13* added comprises the memory address, the 32-bit data (the eight high-order bits are ignored) in which is transferred to the *rd* register. The content of the *rb* register is not altered.

(3) Extension 2

```
ext     imm11           ; imm11(10:0) = imm24(23:13)
ext     imm13           ; = imm24(12:0)
ld.a    %rd,[%rb]       ; memory address = rb + imm24
```

The addressing mode changes to register indirect addressing with displacement, so the content of the *rb* register with the 24-bit immediate *imm24* added comprises the memory address, the 32-bit data (the eight high-order bits are ignored) in which is transferred to the *rd* register. The content of the *rb* register is not altered.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Caution

The *rb* register and the displacement must specify a 32-bit boundary address (two least significant bits = 0). Specifying other address causes an address misaligned interrupt. Note, however, that the data transfer is performed by setting the two least significant bits of the address to 0.

ld.a *%rd*, [*%rb*]+
ld.a *%rd*, [*%rb*]-
ld.a *%rd*, -[*%rb*]

Function

32-bit data transfer with address increment/decrement option

ld.a *%rd*, [*%rb*]+ (with post-increment option)

Standard) *rd*(23:0) ← A[*rb*](23:0), ignored ← A[*rb*](31:24), *rb*(23:0) ← *rb*(23:0) + 4
Extension 1) *rd*(23:0) ← A[*rb*](23:0), ignored ← A[*rb*](31:24), *rb*(23:0) ← *rb*(23:0) + *imm13*
Extension 2) *rd*(23:0) ← A[*rb*](23:0), ignored ← A[*rb*](31:24), *rb*(23:0) ← *rb*(23:0) + *imm24*

ld.a *%rd*, [*%rb*]- (with post-decrement option)

Standard) *rd*(23:0) ← A[*rb*](23:0), ignored ← A[*rb*](31:24), *rb*(23:0) ← *rb*(23:0) - 4
Extension 1) *rd*(23:0) ← A[*rb*](23:0), ignored ← A[*rb*](31:24), *rb*(23:0) ← *rb*(23:0) - *imm13*
Extension 2) *rd*(23:0) ← A[*rb*](23:0), ignored ← A[*rb*](31:24), *rb*(23:0) ← *rb*(23:0) - *imm24*

ld.a *%rd*, -[*%rb*] (with pre-decrement option)

Standard) *rb*(23:0) ← *rb*(23:0) - 4, *rd*(23:0) ← A[*rb*](23:0), ignored ← A[*rb*](31:24)
Extension 1) *rb*(23:0) ← *rb*(23:0) - *imm13*, *rd*(23:0) ← A[*rb*](23:0), ignored ← A[*rb*](31:24)
Extension 2) *rb*(23:0) ← *rb*(23:0) - *imm24*, *rd*(23:0) ← A[*rb*](23:0), ignored ← A[*rb*](31:24)

Code

15	14	13	12	11	10	9–7	6	5	4	3	2–0	
0	0	1	0	0	0	*rd*	0	1	1	1	*rb*	`ld.a %rd, [%rb]+`
0	0	1	0	0	0	*rd*	1	1	1	1	*rb*	`ld.a %rd, [%rb]-`
0	0	1	0	0	0	*rd*	1	0	1	1	*rb*	`ld.a %rd, -[%rb]`

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register indirect `%rb` = `%r0` to `%r7`
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

Two cycles

Description

(1) Address increment/decrement option

Specifying the `[]+`, `[]-`, or `-[]` option will automatically increment/decrement the memory address. This allows the program to simply perform continuous data transfer.

`ld.a %rd, [%rb]+` Load instruction with post-increment option
The memory address will be incremented after the data transfer has finished.

`ld.a %rd, [%rb]-` Load instruction with post-decrement option
The memory address will be decremented after the data transfer has finished.

`ld.a %rd, -[%rb]` Load instruction with pre-decrement option
The memory address will be decremented before starting the data transfer.

The address increment/decrement sizes are listed below.
When no `ext` is used (standard): 4 (32-bit size)
When one `ext` is used (extension 1): *imm13*
When two `ext` are used (extension 2): *imm24*

(2) Standard (example of post-increment option)

```
ld.a   %rd,[%rb]+  ; source memory address = rb
                   ; post increment: rb + 4
```

The 32-bit data (the eight high-order bits are ignored) in the specified memory location is transferred to the *rd* register. The *rb* register contains the memory address to be accessed. The memory address will be incremented by four bytes after the data transfer has finished.

(3) Extension 1 (example of post-decrement option)

```
ext    imm13
ld.a   %rd,[%rb]-  ; source memory address = rb
                   ; post decrement: rb - imm13
```

The 32-bit data (the eight high-order bits are ignored) in the specified memory location is transferred to the *rd* register. The *rb* register contains the memory address to be accessed. The memory address will be decremented by *imm13* bytes after the data transfer has finished.

(4) Extension 2 (example of pre-decrement option)

```
ext   imm11        ; imm11(10:0) = imm24(23:13)
ext   imm13        ; = imm24(12:0)
ld.a    %rd,-[%rb] ; source memory address = rb - imm24
```

After the memory address specified by the *rb* register is decremented by *imm24* bytes, the 32-bit data (the eight high-order bits are ignored) in the decremented address is transferred to the *rd* register.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Caution The *rb* register and the immediate value must specify a 32-bit boundary address (two least significant bits = 0). Specifying other address causes an address misaligned interrupt. Note, however, that the data transfer is performed by setting the two least significant bits of the address to 0.

ld.a *%rd*, [%sp]

Function

32-bit data transfer

Standard) *rd*(23:0) ← A[sp](23:0), ignored ← A[sp](31:24)

Extension 1) *rd*(23:0) ← A[sp + *imm13*](23:0), ignored ← A[sp + *imm13*](31:24)

Extension 2) *rd*(23:0) ← A[sp + *imm24*](23:0), ignored ← A[sp + *imm24*](31:24)

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	1	1	1		*rd*		0	0	1	1	0	0	0

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register indirect `%sp`

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle (two cycles when the `ext` instruction or an increment/decrement option is used)

Description

(1) Standard

```
ld.a    %rd,[%sp]       ; memory address = sp
```

The 32-bit data (the eight high-order bits are ignored) in the specified memory location is transferred to the *rd* register. The SP contains the memory address to be accessed.

(2) Extension 1

```
ext     imm13
ld.a    %rd,[%sp]       ; memory address = sp + imm13
```

The `ext` instruction changes the addressing mode to register indirect addressing with displacement. As a result, the content of the SP with the 13-bit immediate *imm13* added comprises the memory address, the 32-bit data (the eight high-order bits are ignored) in which is transferred to the *rd* register. The content of the SP is not altered.

(3) Extension 2

```
ext     imm11           ; imm11(10:0) = imm24(23:13)
ext     imm13           ; = imm24(12:0)
ld.a    %rd,[%sp]       ; memory address = sp + imm24
```

The addressing mode changes to register indirect addressing with displacement, so the content of the SP with the 24-bit immediate *imm24* added comprises the memory address, the 32-bit data (the eight high-order bits are ignored) in which is transferred to the *rd* register. The content of the SP is not altered.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Caution

The displacement must specify a 32-bit boundary address (two least significant bits = 0). Specifying other address causes an address misaligned interrupt. Note, however, that the data transfer is performed by setting the two least significant bits of the address to 0.

ld.a *%rd*, [%sp]+
ld.a *%rd*, [%sp]-
ld.a *%rd*, -[%sp]

Function

32-bit data transfer with address increment/decrement option

ld.a *%rd*, [%sp]+ (with post-increment option)

Standard) *rd*(23:0) ← A[sp](23:0), ignored ← A[sp](31:24), sp(23:0) ← sp(23:0) + 4
Extension 1) *rd*(23:0) ← A[sp](23:0), ignored ← A[sp](31:24), sp(23:0) ← sp(23:0) + *imm13*
Extension 2) *rd*(23:0) ← A[sp](23:0), ignored ← A[sp](31:24), sp(23:0) ← sp(23:0) + *imm24*

ld.a *%rd*, [%sp]- (with post-decrement option)

Standard) *rd*(23:0) ← A[sp](23:0), ignored ← A[sp](31:24), sp(23:0) ← sp(23:0) - 4
Extension 1) *rd*(23:0) ← A[sp](23:0), ignored ← A[sp](31:24), sp(23:0) ← sp(23:0) - *imm13*
Extension 2) *rd*(23:0) ← A[sp](23:0), ignored ← A[sp](31:24), sp(23:0) ← sp(23:0) - *imm24*

ld.a *%rd*, -[%sp] (with pre-decrement option)

Standard) sp(23:0) ← sp(23:0) - 4, *rd*(23:0) ← A[sp](23:0), ignored ← A[sp](31:24)
Extension 1) sp(23:0) ← sp(23:0) - *imm13, rd*(23:0) ← A[sp](23:0), ignored ← A[sp](31:24)
Extension 2) sp(23:0) ← sp(23:0) - *imm24, rd*(23:0) ← A[sp](23:0), ignored ← A[sp](31:24)

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	1	1	1	1		*rd*		0	1	1	1	0	0	0	`ld.a %rd,[%sp]+`
0	0	1	1	1	1		*rd*		1	1	1	1	0	0	0	`ld.a %rd,[%sp]-`
0	0	1	1	1	1		*rd*		1	0	1	1	0	0	0	`ld.a %rd,-[%sp]`

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register indirect `%sp`
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

Two cycles

Description

(1) Address increment/decrement option

Specifying the `[]+`, `[]-`, or `-[]` option will automatically increment/decrement the memory address. This allows the program to simply perform continuous data transfer.

`ld.a %rd,[%sp]+` Load instruction with post-increment option
The memory address will be incremented after the data transfer has finished.

`ld.a %rd,[%sp]-` Load instruction with post-decrement option
The memory address will be decremented after the data transfer has finished.

`ld.a %rd,-[%sp]` Load instruction with pre-decrement option
The memory address will be decremented before starting the data transfer.

The address increment/decrement sizes are listed below.
When no `ext` is used (standard): 4 (32-bit size)
When one `ext` is used (extension 1): *imm13*
When two `ext` are used (extension 2): *imm24*

(2) Standard (example of post-increment option)

```
ld.a  %rd,[%sp]+  ; source memory address = sp
                  ; post increment: sp + 4
```

The 32-bit data (the eight high-order bits are ignored) in the specified memory location is transferred to the *rd* register. The SP contains the memory address to be accessed. The memory address will be incremented by four bytes after the data transfer has finished.

(3) Extension 1 (example of post-decrement option)

```
ext   imm13
ld.a  %rd,[%sp]-  ; source memory address = sp
                  ; post decrement: sp - imm13
```

The 32-bit data (the eight high-order bits are ignored) in the specified memory location is transferred to the *rd* register. The SP contains the memory address to be accessed. The memory address will be decremented by *imm13* bytes after the data transfer has finished.

(4) Extension 2 (example of pre-decrement option)

```
ext  imm11        ; imm11(10:0) = imm24(23:13)
ext  imm13        ; = imm24(12:0)
ld   %rd,-[%sp]   ; source memory address = sp - imm24
```

After the memory address specified by the SP is decremented by *imm24* bytes, the 32-bit data (the eight high-order bits are ignored) in the decremented address is transferred to the *rd* register.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Caution The immediate must specify a 32-bit boundary address (two least significant bits = 0). Specifying other address causes an address misaligned interrupt. Note, however, that the data transfer is performed by setting the two least significant bits of the address to 0.

ld.a *%rd*, [%sp + *imm7*]

Function

32-bit data transfer

Standard) *rd*(23:0) ← A[sp + *imm7*](23:0), ignored ← A[sp + *imm7*](31:24)
Extension 1) *rd*(23:0) ← A[sp + *imm20*](23:0), ignored ← A[sp + *imm20*](31:24)
Extension 2) *rd*(23:0) ← A[sp + *imm24*](23:0), ignored ← A[sp + *imm24*](31:24)

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	0	1	1	*rd*			*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register indirect with displacement
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

Two cycles

Description

(1) Standard

```
ld.a   %rd,[%sp + imm7]      ; memory address = sp + imm7
```

The 32-bit data (the eight high-order bits are ignored) in the specified memory location is transferred to the *rd* register. The content of the current SP with the 7-bit immediate *imm7* added as displacement comprises the memory address to be accessed.

(2) Extension 1

```
ext    imm13                 ; = imm20(19:7)
ld.a   %rd,[%sp + imm7]      ; memory address = sp + imm20,
                             ; imm7 = imm20(6:0)
```

The `ext` instruction extends the displacement to a 20-bit quantity. As a result, the content of the SP with the 20-bit immediate *imm20* added comprises the memory address, the 32-bit data (the eight high-order bits are ignored) in which is transferred to the *rd* register.

(3) Extension 2

```
ext    imm4                  ; imm4(3:0) = imm24(23:20)
ext    imm13                 ; = imm24(19:7)
ld.a   %rd,[%sp + imm7]      ; memory address = sp + imm24,
                             ; imm7 = imm24(6:0)
```

The two `ext` instructions extend the displacement to a 24-bit quantity. As a result, the content of the SP with the 24-bit immediate *imm24* added comprises the memory address, the 32-bit data (the eight high-order bits are ignored) in which is transferred to the *rd* register.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

```
ext    0x1
ld.a   %r0,[%sp + 0x4]   ; r0 ← [sp + 0x84]
```

Caution

The SP and the displacement must specify a 32-bit boundary address (two least significant bits = 0). Specifying other address causes an address misaligned interrupt. Note, however, that the data transfer is performed by setting the two least significant bits of the address to 0.

ld.a *%rd*, [*imm7*]

Function

32-bit data transfer

Standard) *rd*(23:0) ← A[*imm7*](23:0), ignored ← A[*imm7*](31:24)
Extension 1) *rd*(23:0) ← A[*imm20*](23:0), ignored ← A[*imm20*](31:24)
Extension 2) *rd*(23:0) ← A[*imm24*](23:0), ignored ← A[*imm24*](31:24)

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	0	0	1	1	*rd*			*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Immediate data (unsigned)
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
ld.a   %rd,[imm7]     ; memory address = imm7
```

The 32-bit data (the eight high-order bits are ignored) in the memory address specified with the 7-bit immediate *imm7* is transferred to the *rd* register.

(2) Extension 1

```
ext    imm13          ; = imm20(19:7)
ld.a   %rd,[imm7]     ; memory address = imm20, imm7 = imm20(6:0)
```

The `ext` instruction extends the displacement to a 20-bit quantity. As a result, the 32-bit data (the eight high-order bits are ignored) in the memory address specified with the 20-bit immediate *imm20* is transferred to the *rd* register.

(3) Extension 2

```
ext    imm4           ; imm4(3:0) = imm24(23:20)
ext    imm13          ; = imm24(19:7)
ld.a   %rd,[imm7]     ; memory address = imm24, imm7 = imm24(6:0)
```

The two `ext` instructions extend the displacement to a 24-bit quantity. As a result, the 32-bit data (the eight high-order bits are ignored) in the memory address specified with the 24-bit immediate *imm24* is transferred to the *rd* register.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

```
ext    0x1
ld.a   %r0,[0x4]  ; r0 ← [0x84]
```

Caution

The *imm7* must specify a 32-bit boundary address (two least significant bits = 0). Specifying other address causes an address misaligned interrupt. Note, however, that the data transfer is performed by setting the two least significant bits of the address to 0.

ld.a *%rd, imm7*

Function

24-bit data transfer

Standard) *rd*(6:0) ← *imm7*, *rd*(23:7) ← 0

Extension 1) *rd*(19:0) ← *imm20*, *rd*(23:20) ← 0

Extension 2) *rd*(23:0) ← *imm24*

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	1	1	1	*r d*			*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Immediate data (unsigned)

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
ld.a   %rd,imm7        ; rd ← imm7 (zero-extended)
```

The 7-bit immediate *imm7* is loaded to the *rd* register after being zero-extended.

(2) Extension 1

```
ext    imm13           ; = sign20(19:7)
ld.a   %rd,imm7        ; rd ← imm20 (zero-extended),
                       ; imm7 = imm20(6:0)
```

The immediate data is extended into a 20-bit quantity by the `ext` instruction and it is loaded to the *rd* register after being zero-extended.

(3) Extension 2

```
ext    imm4            ; imm4(3:0) = imm24(23:20)
ext    imm13           ; = imm24(19:7)
ld.a   %rd,imm7        ; rd ← imm24, imm7 = imm24(6:0)
```

The immediate data is extended into a 24-bit quantity by the `ext` instruction and it is loaded to the *rd* register.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

```
ld.a  %r0,0x3f  ; r0 ← 0x00003f
```

ld.a %sp, %*rs*

Function
24-bit data transfer
Standard) sp(23:2) ← *rs*(23:2), sp(1:0) ← 0
Extension 1) Unusable
Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	1	1	1		*rs*		1	0	1	0	0	0	0

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode
Src: Register direct `%rs` = `%r0` to `%r7`
Dst: Register direct `%sp`

CLK
One cycle

Description
The content of the *rs* register is transferred to the SP.

Example
```
ld.a   %sp,%r0    ; sp ← r0
```

Caution
In data transfer to the SP, the low-order two bits of the source data are always handled as 0.

ld.a %sp, *imm7*

Function

24-bit data transfer

Standard) sp(6:2) ← *imm7*(6:2), sp(23:7) ← 0, sp(1:0) ← 0
Extension 1) sp(19:2) ← *imm20*(19:2), sp(23:20) ← 0, sp(1:0) ← 0
Extension 2) sp(23:2) ← *imm24*(23:2), sp(1:0) ← 0

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	1	0	0	0	*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Immediate data (unsigned)
Dst: Register direct `%sp`

CLK

One cycle

Description

(1) Standard

```
ld.a   %sp,imm7        ; sp ← imm7 (zero-extended)
```

The 7-bit immediate *imm7* is loaded to the SP after being zero-extended.

(2) Extension 1

```
ext    imm13           ; = sign20(19:7)
ld.a   %sp,imm7        ; sp ← imm20 (zero-extended),
                       ; imm7 = imm20(6:0)
```

The immediate data is extended into a 20-bit quantity by the `ext` instruction and it is loaded to the SP after being zero-extended.

(3) Extension 2

```
ext    imm4            ; imm4(3:0) = imm24(23:20)
ext    imm13           ; = imm24(19:7)
ld.a   %sp,imm7        ; sp ← imm24, imm7 = imm24(6:0)
```

The immediate data is extended into a 24-bit quantity by the `ext` instruction and it is loaded to the SP.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

```
ext    0x8
ld.a   %sp,0x0    ; sp ← 0x400
```

Caution

In data transfer to the SP, the low-order two bits of the source data are always handled as 0.

ld.a [%*rb*], %*rs*

Function

32-bit data transfer

Standard) A[*rb*](23:0) ← *rs*(23:0), A[*rb*](31:24) ← 0

Extension 1) A[*rb* + *imm13*](23:0) ← *rs*(23:0), A[*rb* + *imm13*](31:24) ← 0

Extension 2) A[*rb* + *imm24*](23:0) ← *rs*(23:0), A[*rb* + *imm24*](31:24) ← 0

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	0	1		*rs*		0	0	1	1		*rb*	

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register indirect `%rb` = `%r0` to `%r7`

CLK

One cycle (two cycles when the `ext` instruction is used

Description

(1) Standard

```
ld.a   [%rb],%rs      ; memory address = rb
```

The content of the *rs* register (24-bit data) is transferred to the specified memory location. The *rb* register contains the memory address to be accessed. This instruction writes 32-bit data with the eight high-order bits set to 0 in the memory.

(2) Extension 1

```
ext    imm13
ld.a   [%rb],%rs      ; memory address = rb + imm13
```

The `ext` instruction changes the addressing mode to register indirect addressing with displacement. As a result, the content of the *rs* register is transferred to the address indicated by the content of the *rb* register with the 13-bit immediate *imm13* added. The content of the *rb* register is not altered.

(3) Extension 2

```
ext    imm11          ; imm11(10:0) = imm24(23:13)
ext    imm13          ; = imm24(12:0)
ld.a   [%rb],%rs      ; memory address = rb + imm24
```

The addressing mode changes to register indirect addressing with displacement, so the content of the *rs* register is transferred to the address indicated by the content of the *rb* register with the 24-bit immediate *imm24* added. The content of the *rb* register is not altered.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Caution

The *rb* register and the displacement must specify a 32-bit boundary address (two least significant bits = 0). Specifying other address causes an address misaligned interrupt. Note, however, that the data transfer is performed by setting the two least significant bits of the address to 0.

ld.a [%rb]+, %rs
ld.a [%rb]-, %rs
ld.a -[%rb], %rs

Function

32-bit data transfer with address increment/decrement option

ld.a [%rb]+, %rs (with post-increment option)

Standard) A[*rb*](23:0) ← *rs*(23:0), A[*rb*](31:24) ← 0, *rb*(23:0) ← *rb*(23:0) + 4
Extension 1) A[*rb*](23:0) ← *rs*(23:0), A[*rb*](31:24) ← 0, *rb*(23:0) ← *rb*(23:0) + *imm13*
Extension 2) A[*rb*](23:0) ← *rs*(23:0), A[*rb*](31:24) ← 0, *rb*(23:0) ← *rb*(23:0) + *imm24*

ld.a [%rb]-, %rs (with post-decrement option)

Standard) A[*rb*](23:0) ← *rs*(23:0), A[*rb*](31:24) ← 0, *rb*(23:0) ← *rb*(23:0) - 4
Extension 1) A[*rb*](23:0) ← *rs*(23:0), A[*rb*](31:24) ← 0, *rb*(23:0) ← *rb*(23:0) - *imm13*
Extension 2) A[*rb*](23:0) ← *rs*(23:0), A[*rb*](31:24) ← 0, *rb*(23:0) ← *rb*(23:0) - *imm24*

ld.a -[%rb], %rs (with pre-decrement option)

Standard) *rb*(23:0) ← *rb*(23:0) - 4, A[*rb*](23:0) ← *rs*(23:0), A[*rb*](31:24) ← 0
Extension 1) *rb*(23:0) ← *rb*(23:0) - *imm13,* A[*rb*](23:0) ← *rs*(23:0), A[*rb*](31:24) ← 0
Extension 2) *rb*(23:0) ← *rb*(23:0) - *imm24,* A[*rb*](23:0) ← *rs*(23:0), A[*rb*](31:24) ← 0

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	1	0	0	1		rs		0	1	1	1		rb		`ld.a [%rb]+,%rs`
0	0	1	0	0	1		rs		1	1	1	1		rb		`ld.a [%rb]-,%rs`
0	0	1	0	0	1		rs		1	0	1	1		rb		`ld.a -[%rb],%rs`

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`
Dst: Register indirect `%rb` = `%r0` to `%r7`

CLK

Two cycles

Description

(1) Address increment/decrement option

Specifying the `[]+`, `[]-`, or `-[]` option will automatically increment/decrement the memory address. This allows the program to simply perform continuous data transfer.

`ld.a [%rb]+,%rs` Load instruction with post-increment option
The memory address will be incremented after the data transfer has finished.

`ld.a [%rb]-,%rs` Load instruction with post-decrement option
The memory address will be decremented after the data transfer has finished.

`ld.a -[%rb],%rs` Load instruction with pre-decrement option
The memory address will be decremented before starting the data transfer.

The address increment/decrement sizes are listed below.
When no `ext` is used (standard): 4 (32-bit size)
When one `ext` is used (extension 1): *imm13*
When two `ext` are used (extension 2): *imm24*

(2) Standard (example of post-increment option)

```
ld.a   [%rb]+,%rs     ; Destination memory address = rb
                      ; post increment: rb + 4
```

The content of the *rs* register (24-bit data) is transferred to the specified memory location. The *rb* register contains the memory address to be accessed. This instruction writes 32-bit data with the eight high-order bits set to 0 in the memory. The memory address will be incremented by two bytes after the data transfer has finished.

(3) Extension 1 (example of post-decrement option)

```
ext    imm13
ld.a   [%rb]-,%rs     ; Destination memory address = rb
                      ; post decrement: rb - imm13
```

The content of the *rs* register (24-bit data) is transferred to the specified memory location. The *rb* register contains the memory address to be accessed. This instruction writes 32-bit data with the eight high-order bits set to 0 in the memory. The memory address will be decremented by *imm13* bytes after the data transfer has finished.

(4) Extension 2 (example of pre-decrement option)

```
ext    imm11          ; imm11(10:0) = imm24(23:13)
ext    imm13          ; = imm24(12:0)
ld.a   -[%rb],%rs     ; Destination memory address = rb - imm24
```

After the memory address specified by the *rb* register is decremented by *imm24* bytes, the content of the *rs* register (24-bit data) is transferred to the decremented address. This instruction writes 32-bit data with the eight high-order bits set to 0 in the memory.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Caution The *rb* register and the immediate value must specify a 32-bit boundary address (two least significant bits = 0). Specifying other address causes an address misaligned interrupt. Note, however, that the data transfer is performed by setting the two least significant bits of the address to 0.

ld.a [%sp], *%rs*

Function

32-bit data transfer

Standard) A[sp](23:0) ← *rs*(23:0), A[sp](31:24) ← 0

Extension 1) A[sp + *imm13*](23:0) ← *rs*(23:0), A[sp + *imm13*](31:24) ← 0

Extension 2) A[sp + *imm24*](23:0) ← *rs*(23:0), A[sp + *imm24*](31:24) ← 0

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	1	1	1		*rs*		0	0	1	1	1	0	0

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register indirect `%sp`

CLK

One cycle (two cycles when the `ext` instruction is used)

Description

(1) Standard

```
ld.a   [%sp],%rs      ; memory address = sp
```

The content of the *rs* register (24-bit data) is transferred to the specified memory location. The SP contains the memory address to be accessed. This instruction writes 32-bit data with the eight high-order bits set to 0 in the memory.

(2) Extension 1

```
ext    imm13
ld.a   [%sp],%rs      ; memory address = sp + imm13
```

The `ext` instruction changes the addressing mode to register indirect addressing with displacement. As a result, the content of the *rs* register is transferred to the address indicated by the content of the SP with the 13-bit immediate *imm13* added. The content of the SP is not altered.

(3) Extension 2

```
ext    imm11          ; imm11(10:0) = imm24(23:13)
ext    imm13          ; = imm24(12:0)
ld.a   [%sp],%rs      ; memory address = sp + imm24
```

The addressing mode changes to register indirect addressing with displacement, so the content of the *rs* register is transferred to the address indicated by the content of the SP with the 24-bit immediate *imm24* added. The content of the SP is not altered.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Caution

The SP and the displacement must specify a 32-bit boundary address (two least significant bits = 0). Specifying other address causes an address misaligned interrupt. Note, however, that the data transfer is performed by setting the two least significant bits of the address to 0.

ld.a [%sp]+, *%rs*
ld.a [%sp]-, *%rs*
ld.a -[%sp], *%rs*

Function

32-bit data transfer with address increment/decrement option

ld.a [%sp]+, *%rs* (with post-increment option)

Standard) A[sp](23:0) ← *rs*(23:0), A[sp](31:24) ← 0, sp(23:0) ← sp(23:0) + 4
Extension 1) A[sp](23:0) ← *rs*(23:0), A[sp](31:24) ← 0, sp(23:0) ← sp(23:0) + *imm13*
Extension 2) A[sp](23:0) ← *rs*(23:0), A[sp](31:24) ← 0, sp(23:0) ← sp(23:0) + *imm24*

ld.a [%sp]-, *%rs* (with post-decrement option)

Standard) A[sp](23:0) ← *rs*(23:0), A[sp](31:24) ← 0, sp(23:0) ← sp(23:0) - 4
Extension 1) A[sp](23:0) ← *rs*(23:0), A[sp](31:24) ← 0, sp(23:0) ← sp(23:0) - *imm13*
Extension 2) A[sp](23:0) ← *rs*(23:0), A[sp](31:24) ← 0, sp(23:0) ← sp(23:0) - *imm24*

ld.a -[%sp], *%rs* (with pre-decrement option)

Standard) sp(23:0) ← sp(23:0) - 4, A[sp](23:0) ← *rs*(23:0), A[sp](31:24) ← 0
Extension 1) sp(23:0) ← sp(23:0) - *imm13,* A[sp](23:0) ← *rs*(23:0), A[sp](31:24) ← 0
Extension 2) sp(23:0) ← sp(23:0) - *imm24,* A[sp](23:0) ← *rs*(23:0), A[sp](31:24) ← 0

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	1	1	1	1		*rs*		0	1	1	1	1	0	0	`ld.a [%sp]+,%rs`
0	0	1	1	1	1		*rs*		1	1	1	1	1	0	0	`ld.a [%sp]-,%rs`
0	0	1	1	1	1		*rs*		1	0	1	1	1	0	0	`ld.a -[%sp],%rs`

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`
Dst: Register indirect `%sp`

CLK

Two cycles

Description

(1) Address increment/decrement option

Specifying the `[]+`, `[]-`, or `-[]` option will automatically increment/decrement the memory address. This allows the program to simply perform continuous data transfer.

`ld.a [%sp]+,%rs` Load instruction with post-increment option
The memory address will be incremented after the data transfer has finished.

`ld.a [%sp]-,%rs` Load instruction with post-decrement option
The memory address will be decremented after the data transfer has finished.

`ld.a -[%sp],%rs` Load instruction with pre-decrement option
The memory address will be decremented before starting the data transfer.

The address increment/decrement sizes are listed below.
When no `ext` is used (standard): 4 (32-bit size)
When one `ext` is used (extension 1): *imm13*
When two `ext` are used (extension 2): *imm24*

(2) Standard (example of post-increment option)

```
ld.a   [%sp]+,%rs     ; Destination memory address = sp
                      ; post increment: sp + 4
```

The content of the *rs* register (24-bit data) is transferred to the specified memory location. The SP contains the memory address to be accessed. This instruction writes 32-bit data with the eight high-order bits set to 0 in the memory. The memory address will be incremented by two bytes after the data transfer has finished.

(3) Extension 1 (example of post-decrement option)

```
ext    imm13
ld.a   [%sp]-,%rs     ; Destination memory address = sp
                      ; post decrement: sp - imm13
```

The content of the *rs* register (24-bit data) is transferred to the specified memory location. The SP contains the memory address to be accessed. This instruction writes 32-bit data with the eight high-order bits set to 0 in the memory. The memory address will be decremented by *imm13* bytes after the data transfer has finished.

(4) Extension 2 (example of pre-decrement option)

```
ext    imm11          ; imm11(10:0) = imm24(23:13)
ext    imm13          ; = imm24(12:0)
ld.a   -[%sp],%rs     ; Destination memory address = sp - imm24
```

After the memory address specified by the SP is decremented by *imm24* bytes, the content of the *rs* register (24-bit data) is transferred to the decremented address. This instruction writes 32-bit data with the eight high-order bits set to 0 in the memory.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Caution The SP and the immediate value must specify a 32-bit boundary address (two least significant bits = 0). Specifying other address causes an address misaligned interrupt. Note, however, that the data transfer is performed by setting the two least significant bits of the address to 0.

ld.a [%sp + *imm7*], %*rs*

Function

32-bit data transfer

Standard) A[sp + *imm7*](23:0) ← *rs*(23:0), A[sp + *imm7*](31:24) ← 0
Extension 1) A[sp + *imm20*](23:0) ← *rs*(23:0), A[sp + *imm20*](31:24) ← 0
Extension 2) A[sp + *imm24*](23:0) ← *rs*(23:0), A[sp + *imm24*](31:24) ← 0

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	1	1	1	*rs*			*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`
Dst: Register indirect with displacement

CLK

Two cycles

Description

(1) Standard

```
ld.a   [%sp + imm7],%rs      ; memory address = sp + imm7
```

The content of the *rs* register is transferred to the specified memory location. The content of the current SP with the 7-bit immediate *imm7* added as displacement comprises the memory address to be accessed. This instruction writes 32-bit data with the eight high-order bits set to 0 in the memory.

(2) Extension 1

```
ext    imm13                 ; = imm20(19:7)
ld.a   [%sp + imm7],%rs      ; memory address = sp + imm20,
                             ; imm7 = imm20(6:0)
```

The `ext` instruction extends the displacement to a 20-bit quantity. As a result, the content of the *rs* register is transferred to the address indicated by the content of the SP with the 20-bit immediate *imm20* added.

(3) Extension 2

```
ext    imm4                  ; imm4(3:0) = imm24(23:20)
ext    imm13                 ; = imm24(19:7)
ld.a   [%sp + imm7],%rs      ; memory address = sp + imm24,
                             ; imm7 = imm24(6:0)
```

The two `ext` instructions extend the displacement to a 24-bit quantity. As a result, the content of the *rs* register is transferred to the address indicated by the content of the SP with the 24-bit immediate *imm24* added.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

```
ext   0x1
ld.a  [%sp + 0x4],%r0  ; [sp + 0x84] ← r0
```

Caution

The SP and the displacement must specify a 32-bit boundary address (two least significant bits = 0). Specifying other address causes an address misaligned interrupt. Note, however, that the data transfer is performed by setting the two least significant bits of the address to 0.

ld.a [*imm7*], %*rs*

Function

32-bit data transfer

Standard) A[*imm7*](23:0) ← *rs*(23:0), A[*imm7*](31:24) ← 0
Extension 1) A[*imm20*](23:0) ← *rs*(23:0), A[*imm20*](31:24) ← 0
Extension 2) A[*imm24*](23:0) ← *rs*(23:0), A[*imm24*](31:24) ← 0

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	0	1	1	1	*rs*			*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`
Dst: Immediate data (unsigned)

CLK

One cycle

Description

(1) Standard

```
ld.a   [imm7],%rs     ; memory address = imm7
```

The content of the *rs* register is transferred to the memory address specified with the 7-bit immediate *imm7*. This instruction writes 32-bit data with the eight high-order bits set to 0 in the memory.

(2) Extension 1

```
ext    imm13          ; = imm20(19:7)
ld.a   [imm7],%rs     ; memory address = imm20, imm7 = imm20(6:0)
```

The `ext` instruction extends the displacement to a 20-bit quantity. As a result, the content of the *rs* register is transferred to the memory address specified with the 20-bit immediate *imm20*.

(3) Extension 2

```
ext    imm4           ; imm4(3:0) = imm24(23:20)
ext    imm13          ; = imm24(19:7)
ld.a   [imm7],%rs     ; memory address = imm24, imm7 = imm24(6:0)
```

The two `ext` instructions extend the displacement to a 24-bit quantity. As a result, the content of the *rs* register is transferred to the memory address specified with the 24-bit immediate *imm24*.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

```
ext    0x1
ld.a   [0x4],%r0         ; [0x84] ← r0
```

Caution

The *imm7* must specify a 32-bit boundary address (two least significant bits = 0). Specifying other address causes an address misaligned interrupt. Note, however, that the data transfer is performed by setting the two least significant bits of the address to 0.

ld.b *%rd, %rs*

Function

Signed byte data transfer

Standard) *rd*(7:0) ← *rs*(7:0), *rd*(15:8) ← *rs*(7), *rd*(23:16) ← 0

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	1	0		*rd*		0	0	0	0		*rs*	

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

The eight low-order bits of the *rs* register are transferred to the *rd* register after being sign-extended to 16 bits. The eight high-order bits of the *rd* register are set to 0.

(2) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit.

Example

```
ld.b  %r0,%r1   ; r0 ← r1(7:0) sign-extended
```

ld.b *%rd*, [*%rb*]

Function

Signed byte data transfer

Standard) *rd*(7:0) ← B[*rb*], *rd*(15:8) ← B[*rb*](7), *rd*(23:16) ← 0

Extension 1) *rd*(7:0) ← B[*rb* + *imm13*], *rd*(15:8) ← B[*rb* + *imm13*](7), *rd*(24:16) ← 0

Extension 2) *rd*(7:0) ← B[*rb* + *imm24*], *rd*(15:8) ← B[*rb* + *imm24*](7), *rd*(24:16) ← 0

Code

15	14	13	12	11	10	9–7	6	5	4	3	2–0	
0	0	1	0	0	0	*rd*	0	0	0	0	*rb*	`ld.b %rd,[%rb]`

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src:Register indirect `%rb` = `%r0` to `%r7`

Dst:Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle (two cycles when the `ext` instruction is used)

Description

(1) Standard

```
ld.b   %rd,[%rb]      ; memory address = rb
```

The byte data in the specified memory location is transferred to the *rd* register after being sign-extended to 16 bits. The *rb* register contains the memory address to be accessed. The eight high-order bits of the *rd* register are set to 0.

(2) Extension 1

```
ext    imm13
ld.b   %rd,[%rb]      ; memory address = rb + imm13
```

The `ext` instruction changes the addressing mode to register indirect addressing with displacement. As a result, the content of the *rb* register with the 13-bit immediate *imm13* added comprises the memory address, the byte data in which is transferred to the *rd* register after being sign-extended to 16 bits. The eight high-order bits of the *rd* register are set to 0. The content of the *rb* register is not altered.

(3) Extension 2

```
ext    imm11          ; imm11(10:0) = imm24(23:13)
ext    imm13          ; = imm24(12:0)
ld.b   %rd,[%rb]      ; memory address = rb + imm24
```

The addressing mode changes to register indirect addressing with displacement, so the content of the *rb* register with the 24-bit immediate *imm24* added comprises the memory address, the byte data in which is transferred to the *rd* register after being sign-extended to 16 bits. The eight high-order bits of the *rd* register are set to 0. The content of the *rb* register is not altered.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

ld.b *%rd*, [*%rb*]+
ld.b *%rd*, [*%rb*]-
ld.b *%rd*, -[*%rb*]

Function

Signed byte data transfer with address increment/decrement option

ld.b *%rd*, [*%rb*]+ (with post-increment option)

Standard) *rd*(7:0) ← B[*rb*], *rd*(15:8) ← B[*rb*](7), *rd*(23:16) ← 0, *rb*(23:0) ← *rb*(23:0) + 1
Extension 1) *rd*(7:0) ← B[*rb*], *rd*(15:8) ← B[*rb*](7), *rd*(24:16) ← 0, *rb*(23:0) ← *rb*(23:0) + *imm13*
Extension 2) *rd*(7:0) ← B[*rb*], *rd*(15:8) ← B[*rb*](7), *rd*(24:16) ← 0, *rb*(23:0) ← *rb*(23:0) + *imm24*

ld.b *%rd*, [*%rb*]- (with post-decrement option)

Standard) *rd*(7:0) ← B[*rb*], *rd*(15:8) ← B[*rb*](7), *rd*(23:16) ← 0, *rb*(23:0) ← *rb*(23:0) - 1
Extension 1) *rd*(7:0) ← B[*rb*], *rd*(15:8) ← B[*rb*](7), *rd*(24:16) ← 0, *rb*(23:0) ← *rb*(23:0) - *imm13*
Extension 2) *rd*(7:0) ← B[*rb*], *rd*(15:8) ← B[*rb*](7), *rd*(24:16) ← 0, *rb*(23:0) ← *rb*(23:0) - *imm24*

ld.b *%rd*, -[*%rb*] (with pre-decrement option)

Standard) *rb*(23:0) ← *rb*(23:0) - 1, *rd*(7:0) ← B[*rb*], *rd*(15:8) ← B[*rb*](7), *rd*(23:16) ← 0
Extension 1) *rb*(23:0) ← *rb*(23:0) - *imm13, rd*(7:0) ← B[*rb*], *rd*(15:8) ← B[*rb*](7), *rd*(24:16) ← 0
Extension 2) *rb*(23:0) ← *rb*(23:0) - *imm24, rd*(7:0) ← B[*rb*], *rd*(15:8) ← B[*rb*](7), *rd*(24:16) ← 0

Code

15	14	13	12	11	10	9–7	6	5	4	3	2–0	
0	0	1	0	0	0	*rd*	0	1	0	0	*rb*	`ld.b %rd,[%rb]+`
0	0	1	0	0	0	*rd*	1	1	0	0	*rb*	`ld.b %rd,[%rb]-`
0	0	1	0	0	0	*rd*	1	0	0	0	*rb*	`ld.b %rd,-[%rb]`

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register indirect `%rb` = `%r0` to `%r7`
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

Two cycles

Description

(1) Address increment/decrement option

Specifying the `[]+`, `[]-`, or `-[]` option will automatically increment/decrement the memory address. This allows the program to simply perform continuous data transfer.

`ld.b %rd,[%rb]+` Load instruction with post-increment option
The memory address will be incremented after the data transfer has finished.

`ld.b %rd,[%rb]-` Load instruction with post-decrement option
The memory address will be decremented after the data transfer has finished.

`ld.b %rd,-[%rb]` Load instruction with pre-decrement option
The memory address will be decremented before starting the data transfer.

The address increment/decrement sizes are listed below.
When no `ext` is used (standard): 1 (byte size)
When one `ext` is used (extension 1): *imm13*
When two `ext` are used (extension 2): *imm24*

(2) Standard (example of post-increment option)

```
ld.b   %rd,[%rb]+  ; source memory address = rb
                   ; post increment: rb + 1
```

The byte data in the specified memory location is transferred to the *rd* register after being sign-extended to 16 bits. The *rb* register contains the memory address to be accessed. The eight high-order bits of the *rd* register are set to 0. The memory address will be incremented by two bytes after the data transfer has finished.

(3) Extension 1 (example of post-decrement option)

```
ext    imm13
ld.b   %rd,[%rb]-  ; source memory address = rb
                   ; post decrement: rb - imm13
```

The byte data in the specified memory location is transferred to the *rd* register after being sign-extended to 16 bits. The *rb* register contains the memory address to be accessed. The eight high-order bits of the *rd* register are set to 0. The memory address will be decremented by *imm13* bytes after the data transfer has finished.

(4) Extension 2 (example of pre-decrement option)

```
ext    imm11       ; imm11(10:0) = imm24(23:13)
ext    imm13       ; = imm24(12:0)
ld.b   %rd,-[%rb]  ; source memory address = rb - imm24
```

After the memory address specified by the *rb* register is decremented by *imm24* bytes, the byte data in the decremented address is transferred to the *rd* register after being sign-extended to 16 bits. The eight high-order bits of the *rd* register are set to 0.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

ld.b *%rd*, [%sp + *imm7*]

Function

Signed byte data transfer

Standard) *rd*(7:0) ← B[sp + *imm7*], *rd*(15:8) ← B[sp + *imm7*](7), *rd*(23:16) ← 0

Extension 1) *rd*(7:0) ← B[sp + *imm20*], *rd*(15:8) ← B[sp + *imm20*](7), *rd*(23:16) ← 0

Extension 2) *rd*(7:0) ← B[sp + *imm24*], *rd*(15:8) ← B[sp + *imm24*](7), *rd*(23:16) ← 0

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	0	0	0	*rd*			*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register indirect with displacement

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

Two cycles

Description

(1) Standard

```
ld.b   %rd,[%sp + imm7]      ; memory address = sp + imm7
```

The byte data in the specified memory location is transferred to the *rd* register after being sign-extended to 16 bits. The content of the current SP with the 7-bit immediate *imm7* added as displacement comprises the memory address to be accessed. The eight high-order bits of the *rd* register are set to 0.

(2) Extension 1

```
ext    imm13                 ; = imm20(19:7)
ld.b   %rd,[%sp + imm7]      ; memory address = sp + imm20,
                             ; imm7 = imm20(6:0)
```

The `ext` instruction extends the displacement to a 20-bit quantity. As a result, the content of the SP with the 20-bit immediate *imm20* added comprises the memory address, the byte data in which is transferred to the *rd* register after being sign-extended to 16 bits. The eight high-order bits of the *rd* register are set to 0.

(3) Extension 2

```
ext    imm4                  ; imm4(3:0)= imm24(23:20)
ext    imm13                 ; = imm24(19:7)
ld.b   %rd,[%sp + imm7]      ; memory address = sp + imm24,
                             ; imm7 ← imm24(6:0)
```

The two `ext` instructions extend the displacement to a 24-bit quantity. As a result, the content of the SP with the 24-bit immediate *imm24* added comprises the memory address, the byte data in which is transferred to the *rd* register after being sign-extended to 16 bits. The eight high-order bits of the *rd* register are set to 0.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

```
ext    0x1
ld.b   %r0,[%sp + 0x1]   ; r0 ← [sp + 0x81] sign-extended
```

ld.b *%rd*, [*imm7*]

Function

Signed byte data transfer

Standard) *rd*(7:0) ← B[*imm7*], *rd*(15:8) ← B[*imm7*](7), *rd*(23:16) ← 0

Extension 1) *rd*(7:0) ← B[*imm20*], *rd*(15:8) ← B[*imm20*](7), *rd*(23:16) ← 0

Extension 2) *rd*(7:0) ← B[*imm24*], *rd*(15:8) ← B[*imm24*](7), *rd*(23:16) ← 0

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	0	0	0	0	*rd*			*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Immediate data (unsigned)

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
ld.b   %rd,[imm7]      ; memory address = imm7
```

The byte data in the memory address specified with the 7-bit immediate *imm7* is transferred to the *rd* register after being sign-extended to 16 bits. The eight high-order bits of the *rd* register are set to 0.

(2) Extension 1

```
ext    imm13           ; = imm20(19:7)
ld.b   %rd,[imm7]      ; memory address = imm20, imm7 = imm20(6:0)
```

The `ext` instruction extends the displacement to a 20-bit quantity. As a result, the byte data in the memory address specified with the 20-bit immediate *imm20* is transferred to the *rd* register after being sign-extended to 16 bits. The eight high-order bits of the *rd* register are set to 0.

(3) Extension 2

```
ext    imm4            ; imm4(3:0)= imm24(23:20)
ext    imm13           ; = imm24(19:7)
ld.b   %rd,[imm7]      ; memory address = imm24, imm7 ← imm24(6:0)
```

The two `ext` instructions extend the displacement to a 24-bit quantity. As a result, the byte data in the memory address specified with the 24-bit immediate *imm24* is transferred to the *rd* register after being sign-extended to 16 bits. The eight high-order bits of the *rd* register are set to 0.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

```
ext    0x1
ld.b   %r0,[0x1] ; r0 ← [0x81] sign-extended
```

ld.b [%rb], %rs

Function

Signed byte data transfer

Standard) B[*rb*] ← *rs*(7:0)

Extension 1) B[*rb* + *imm13*] ← *rs*(7:0)

Extension 2) B[*rb* + *imm24*] ← *rs*(7:0)

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	0	1		rs		0	0	0	0		rb	

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register indirect `%rb` = `%r0` to `%r7`

CLK

One cycle (two cycles when the `ext` instruction is used)

Description

(1) Standard

```
ld.b   [%rb],%rs      ; memory address = rb
```

The eight low-order bits of the *rs* register are transferred to the specified memory location. The *rb* register contains the memory address to be accessed.

(2) Extension 1

```
ext    imm13
ld.b   [%rb],%rs      ; memory address = rb + imm13
```

The `ext` instruction changes the addressing mode to register indirect addressing with displacement. As a result, the eight low-order bits of the *rs* register are transferred to the address indicated by the content of the *rb* register with the 13-bit immediate *imm13* added. The content of the *rb* register is not altered.

(3) Extension 2

```
ext    imm11          ; imm11(10:0) = imm24(23:13)
ext    imm13          ; = imm24(12:0)
ld.b   [%rb],%rs      ; memory address = rb + imm24
```

The addressing mode changes to register indirect addressing with displacement, so the eight low-order bits of the *rs* register are transferred to the address indicated by the content of the *rb* register with the 24-bit immediate *imm24* added. The content of the *rb* register is not altered.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

ld.b [%*rb*]+, %*rs*
ld.b [%*rb*]-, %*rs*
ld.b -[%*rb*], %*rs*

Function

Signed byte data transfer with address increment/decrement option

ld.b [%*rb*]+, %*rs* (with post-increment option)

Standard) B[*rb*] ← *rs*(7:0), *rb*(23:0) ← *rb*(23:0) + 1
Extension 1) B[*rb*] ← *rs*(7:0), *rb*(23:0) ← *rb*(23:0) + *imm13*
Extension 2) B[*rb*] ← *rs*(7:0), *rb*(23:0) ← *rb*(23:0) + *imm24*

ld.b [%*rb*]-, %*rs* (with post-decrement option)

Standard) B[*rb*] ← *rs*(7:0), *rb*(23:0) ← *rb*(23:0) - 1
Extension 1) B[*rb*] ← *rs*(7:0), *rb*(23:0) ← *rb*(23:0) - *imm13*
Extension 2) B[*rb*] ← *rs*(7:0), *rb*(23:0) ← *rb*(23:0) - *imm24*

ld.b -[%*rb*], %*rs* (with pre-decrement option)

Standard) *rb*(23:0) ← *rb*(23:0) - 1, B[*rb*] ← *rs*(7:0)
Extension 1) *rb*(23:0) ← *rb*(23:0) - *imm13*, B[*rb*] ← *rs*(7:0)
Extension 2) *rb*(23:0) ← *rb*(23:0) - *imm24*, B[*rb*] ← *rs*(7:0)

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	1	0	0	1		*rs*		0	1	0	0		*rb*		`ld.b [%rb]+,%rs`
0	0	1	0	0	1		*rs*		1	1	0	0		*rb*		`ld.b [%rb]-,%rs`
0	0	1	0	0	1		*rs*		1	0	0	0		*rb*		`ld.b -[%rb],%rs`

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`
Dst: Register indirect `%rb` = `%r0` to `%r7`

CLK

Two cycles

Description

(1) Address increment/decrement option

Specifying the `[]+`, `[]-`, or `-[]` option will automatically increment/decrement the memory address. This allows the program to simply perform continuous data transfer.

`ld.b [%rb]+,%rs` Load instruction with post-increment option
The memory address will be incremented after the data transfer has finished.

`ld.b [%rb]-,%rs` Load instruction with post-decrement option
The memory address will be decremented after the data transfer has finished.

`ld.b -[%rb],%rs` Load instruction with pre-decrement option
The memory address will be decremented before starting the data transfer.

The address increment/decrement sizes are listed below.
When no `ext` is used (standard): 1 (byte size)
When one `ext` is used (extension 1): *imm13*
When two `ext` are used (extension 2): *imm24*

(2) Standard (example of post-increment option)

```
ld.b  [%rb]+,%rs    ; Destination memory address = rb
                    ; post increment: rb + 1
```

The eight low-order bits of the *rs* register are transferred to the specified memory location. The *rb* register contains the memory address to be accessed. The memory address will be incremented by two bytes after the data transfer has finished.

(3) Extension 1 (example of post-decrement option)

```
ext   imm13
ld.b  [%rb]-,%rs    ; Destination memory address = rb
                    ; post decrement: rb - imm13
```

The eight low-order bits of the *rs* register are transferred to the specified memory location. The *rb* register contains the memory address to be accessed. The memory address will be decremented by *imm13* bytes after the data transfer has finished.

(4) Extension 2 (example of pre-decrement option)

```
ext   imm11         ; imm11(10:0) = imm24(23:13)
ext   imm13         ; = imm24(12:0)
ld.b  -[%rb],%rs    ; Destination memory address = rb - imm24
```

After the memory address specified by the *rb* register is decremented by *imm24* bytes, the eight low-order bits of the *rs* register are transferred to the decremented address.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

ld.b [%sp + *imm7*], %*rs*

Function

Signed byte data transfer

Standard) B[sp + *imm7*] ← *rs*(7:0)
Extension 1) B[sp + *imm20*] ← *rs*(7:0)
Extension 2) B[sp + *imm24*] ← *rs*(7:0)

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	1	0	0	*r s*			*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct %*rs* = %r0 to %r7
Dst: Register indirect with displacement

CLK

Two cycles

Description

(1) Standard

```
ld.b   [%sp + imm7],%rs     ; memory address = sp + imm7
```

The eight low-order bits of the *rs* register are transferred to the specified memory location. The content of the current SP with the 7-bit immediate *imm7* added as displacement comprises the memory address to be accessed.

(2) Extension 1

```
ext    imm13                ; = imm20(19:7)
ld.b   [%sp + imm7],%rs     ; memory address = sp + imm20,
                            ; imm7 = imm20(6:0)
```

The `ext` instruction extends the displacement to a 20-bit quantity. As a result, the eight low-order bits of the *rs* register are transferred to the address indicated by the content of the SP with the 20-bit immediate *imm20* added.

(3) Extension 2

```
ext    imm4                 ; imm4(3:0) = imm24(23:20)
ext    imm13                ; = imm24(19:7)
ld.b   [%sp + imm7],%rs     ; memory address = sp + imm24,
                            ; imm7 = imm24(6:0)
```

The two `ext` instructions extend the displacement to a 24-bit quantity. As a result, the eight low-order bits of the *rs* register are transferred to the address indicated by the content of the SP with the 24-bit immediate *imm24* added.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

```
ext    0x1
ld.b   [%sp + 0x1],%r0  ; B[sp + 0x81] ← 8 low-order bits of r0
```

ld.b [*imm7*], %*rs*

Function

Signed byte data transfer

Standard) B[*imm7*] ← *rs*(7:0)

Extension 1) B[*imm20*] ← *rs*(7:0)

Extension 2) B[*imm24*] ← *rs*(7:0)

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	0	1	0	0	*rs*			*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct %*rs* = %r0 to %r7

Dst: Immediate data (unsigned)

CLK

One cycle

Description

(1) Standard

```
ld.b   [imm7],%rs     ; memory address = sp + imm7
```

The eight low-order bits of the *rs* register are transferred to the memory address specified with the 7-bit immediate *imm7*.

(2) Extension 1

```
ext    imm13          ; = imm20(19:7)
ld.b   [imm7],%rs     ; memory address = imm20, imm7 = imm20(6:0)
```

The ext instruction extends the displacement to a 20-bit quantity. As a result, the eight low-order bits of the *rs* register are transferred to the memory address specified with the 20-bit immediate *imm20*.

(3) Extension 2

```
ext    imm4           ; imm4(3:0) = imm24(23:20)
ext    imm13          ; = imm24(19:7)
ld.b   [imm7],%rs     ; memory address = imm24, imm7 = imm24(6:0)
```

The two ext instructions extend the displacement to a 24-bit quantity. As a result, the eight low-order bits of the *rs* register are transferred to the memory address specified with the 24-bit immediate *imm24*.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the ext instruction cannot be performed.

Example

```
ext    0x1
ld.b   [0x1],%r0          ; B[0x81] ← 8 low-order bits of r0
```

ld.ca *%rd, %rs*

Function

Transfer data to the coprocessor and get the results

Standard) co_dout0 ← *rd*, co_dout1 ← *rs*, *rd* ← co_din, psr(C, V, Z, N) ← co_cvzn

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	1	0	1		*rd*		0	0	1	1		*rs*	

Flag

IL	IE	C	V	Z	N
–	–	↔	↔	↔	↔

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
ld.ca  %rd,%rs      ; co_dout0 data = rd, co_dout1 data = rs
```

Transfers data set in the *rd* and *rs* registers to the coprocessor and gets the operation results by the coprocessor. The results are loaded to the *rd* register and the C, V, Z, and N flags in the PSR.

(2) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit.

ld.ca *%rd, imm7*

Function

Transfer data to the coprocessor and get the results

Standard) co_dout0 ← *rd*, co_dout1 ← *imm7*, *rd* ← co_din, psr(C, V, Z, N) ← co_cvzn
Extension 1) co_dout0 ← *rd*, co_dout1 ← *imm20*, *rd* ← co_din, psr(C, V, Z, N) ← co_cvzn
Extension 2) co_dout0 ← *rd*, co_dout1 ← *imm24*, *rd* ← co_din, psr(C, V, Z, N) ← co_cvzn

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	1	1		*rd*					*imm7*			

Flag

IL	IE	C	V	Z	N
–	–	↔	↔	↔	↔

Mode

Src: Immediate data (unsigned)
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
ld.ca   %rd,imm7      ; co_dout0 data = rd, co_dout1 data = imm7
```

Transfers data set in the *rd* register and 7-bit immediate *imm7* to the coprocessor and gets the operation results by the coprocessor. The results are loaded to the *rd* register and the C, V, Z, and N flags in the PSR.

(2) Extension 1

```
ext     imm13         ; = imm20(19:7)
ld.ca   %rd,imm7      ; co_dout0 data = rd
                      ; co_dout1 data = imm20, imm7 = imm20(6:0)
```

The `ext` instruction extends the immediate to a 20-bit quantity. As a result, data set in the *rd* register and 20-bit immediate *imm20* are transferred to the coprocessor and the results are loaded to the *rd* register and the C, V, Z, and N flags in the PSR.

(3) Extension 2

```
ext     imm4          ; imm4(3:0) = imm24(23:20)
ext     imm13         ; = imm24(19:7)
ld.ca   %rd,imm7      ; co_dout0 data = rd
                      ; co_dout1 data = imm24, imm7 ← imm24(6:0)
```

The two `ext` instructions extend the displacement to a 24-bit quantity. As a result, data set in the *rd* register and 24-bit immediate *imm24* are transferred to the coprocessor and the results are loaded to the *rd* register and the C, V, Z, and N flags in the PSR.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

ld.cf *%rd, %rs*

Function

Transfer data to the coprocessor and get the flag status

Standard) co_dout0 ← *rd*, co_dout1 ← *rs*, psr(C, V, Z, N) ← co_cvzn
Extension 1) Unusable
Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	1	0	1		*rd*		0	0	0	1		*rs*	

Flag

IL	IE	C	V	Z	N
–	–	↔	↔	↔	↔

Mode

Src: Register direct `%rs` = `%r0` to `%r7`
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
ld.cf  %rd,%rs      ; co_dout0 data = rd, co_dout1 data = rs
```

Transfers data set in the *rd* and *rs* registers to the coprocessor and gets the flag status of the coprocessor to the C, V, Z, and N flags in the PSR.

(2) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit.

ld.cf *%rd, imm7*

Function

Transfer data to the coprocessor and get the flag status

Standard) co_dout0 ← *rd*, co_dout1 ← *imm7*, psr(C, V, Z, N) ← co_cvzn

Extension 1) co_dout0 ← *rd*, co_dout1 ← *imm20*, psr(C, V, Z, N) ← co_cvzn

Extension 2) co_dout0 ← *rd*, co_dout1 ← *imm24*, psr(C, V, Z, N) ← co_cvzn

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	0	1	0	1		*r d*					*imm7*			

Flag

IL	IE	C	V	Z	N
–	–	↔	↔	↔	↔

Mode

Src: Immediate data (unsigned)

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
ld.cf   %rd,imm7       ; co_dout0 data = rd, co_dout1 data = imm7
```

Transfers data set in the *rd* register and 7-bit immediate *imm7* to the coprocessor and gets the flag status of the coprocessor to the C, V, Z, and N flags in the PSR.

(2) Extension 1

```
ext     imm13          ; = imm20(19:7)
ld.cf   %rd,imm7       ; co_dout0 data = rd
                       ; co_dout1 data = imm20, imm7 = imm20(6:0)
```

The `ext` instruction extends the immediate to a 20-bit quantity. As a result, data set in the *rd* register and 20-bit immediate *imm20* are transferred to the coprocessor and the flag status is loaded to the C, V, Z, and N flags in the PSR.

(3) Extension 2

```
ext     imm4           ; imm4(3:0) = imm24(23:20)
ext     imm13          ; = imm24(19:7)
ld.cf   %rd,imm7       ; co_dout0 data = rd
                       ; co_dout1 data = imm24, imm7 ← imm24(6:0)
```

The two `ext` instructions extend the displacement to a 24-bit quantity. As a result, data set in the *rd* register and 24-bit immediate *imm24* are transferred to the coprocessor and the flag status is loaded to the C, V, Z, and N flags in the PSR.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

ld.cw *%rd, %rs*

Function

Transfer data to the coprocessor

Standard) co_dout0 ← *rd*, co_dout1 ← *rs*

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	1	0	1		*rd*		0	0	1	0		*rs*	

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
ld.cw  %rd,%rs       ; co_dout0 data = rd, co_dout1 data = rs
```

Transfers data set in the *rd* and *rs* registers to the coprocessor. The *rd* register and the C, V, Z, and N flags in the PSR are not altered.

(2) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit.

ld.cw %rd, imm7

Function

Transfer data to the coprocessor

Standard) co_dout0 ← *rd*, co_dout1 ← *imm7*

Extension 1) co_dout0 ← *rd*, co_dout1 ← *imm20*

Extension 2) co_dout0 ← *rd*, co_dout1 ← *imm24*

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	1	1	0	rd			imm7						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Immediate data (unsigned)

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
ld.cw   %rd,imm7     ; co_dout0 data = rd, co_dout1 data = imm7
```

Transfers data set in the *rd* register and 7-bit immediate *imm7* to the coprocessor. The *rd* register and the C, V, Z, and N flags in the PSR are not altered.

(2) Extension 1

```
ext     imm13        ; = imm20(19:7)
ld.cw   %rd,imm7     ; co_dout0 data = rd
                     ; co_dout1 data = imm20, imm7 = imm20(6:0)
```

The `ext` instruction extends the immediate to a 20-bit quantity. As a result, data set in the *rd* register and 20-bit immediate *imm20* are transferred to the coprocessor. The *rd* register and the C, V, Z, and N flags in the PSR are not altered.

(3) Extension 2

```
ext     imm4         ; imm4(3:0) = imm24(23:20)
ext     imm13        ; = imm24(19:7)
ld.cw   %rd,imm7     ; co_dout0 data = rd
                     ; co_dout1 data = imm24, imm7 ← imm24(6:0)
```

The two `ext` instructions extend the displacement to a 24-bit quantity. As a result, data set in the *rd* register and 24-bit immediate *imm24* are transferred to the coprocessor. The *rd* register and the C, V, Z, and N flags in the PSR are not altered.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

ld.ub *%rd, %rs*

Function

Unsigned byte data transfer

Standard) *rd*(7:0) ← *rs*(7:0), *rd*(15:8) ← 0, *rd*(23:16) ← 0

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	1	0		*rd*		0	0	0	1		*rs*	

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

The eight low-order bits of the *rs* register are transferred to the *rd* register after being zero-extended to 16 bits. The eight high-order bits of the *rd* register are set to 0.

(2) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit.

Example

```
ld.ub  %r0,%r1   ; r0 ← r1(7:0) zero-extended
```

ld.ub *%rd*, [*%rb*]

Function

Unsigned byte data transfer

Standard) *rd*(7:0) ← B[*rb*], *rd*(15:8) ← 0, *rd*(23:16) ← 0

Extension 1) *rd*(7:0) ← B[*rb* + *imm13*], *rd*(15:8) ← 0, *rd*(24:16) ← 0

Extension 2) *rd*(7:0) ← B[*rb* + *imm24*], *rd*(15:8) ← 0, *rd*(24:16) ← 0

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	0	0	*rd*			0	0	0	1	*rb*		

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register indirect `%rb` = `%r0` to `%r7`

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle (two cycles when the `ext` instruction is used)

Description

(1) Standard

```
ld.ub   %rd,[%rb]     ; memory address = rb
```

The byte data in the specified memory location is transferred to the *rd* register after being zero-extended to 16 bits. The *rb* register contains the memory address to be accessed. The eight high-order bits of the *rd* register are set to 0.

(2) Extension 1

```
ext     imm13
ld.ub   %rd,[%rb]     ; memory address = rb + imm13
```

The `ext` instruction changes the addressing mode to register indirect addressing with displacement. As a result, the content of the *rb* register with the 13-bit immediate *imm13* added comprises the memory address, the byte data in which is transferred to the *rd* register after being zero-extended to 16 bits. The eight high-order bits of the *rd* register are set to 0. The content of the *rb* register is not altered.

(3) Extension 2

```
ext     imm11         ; imm11(10:0) = imm24(23:13)
ext     imm13         ; = imm24(12:0)
ld.ub   %rd,[%rb]     ; memory address = rb + imm24
```

The addressing mode changes to register indirect addressing with displacement, so the content of the *rb* register with the 24-bit immediate *imm24* added comprises the memory address, the byte data in which is transferred to the *rd* register after being zero-extended to 16 bits. The eight high-order bits of the *rd* register are set to 0. The content of the *rb* register is not altered.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

ld.ub *%rd*, [*%rb*]+
ld.ub *%rd*, [*%rb*]-
ld.ub *%rd*, -[*%rb*]

Function

Unsigned byte data transfer with address increment/decrement option

ld.ub *%rd*, [*%rb*]+ (with post-increment option)

Standard) *rd*(7:0) ← B[*rb*], *rd*(15:8) ← 0, *rd*(23:16) ← 0, *rb*(23:0) ← *rb*(23:0) + 1
Extension 1) *rd*(7:0) ← B[*rb*], *rd*(15:8) ← 0, *rd*(24:16) ← 0, *rb*(23:0) ← *rb*(23:0) + *imm13*
Extension 2) *rd*(7:0) ← B[*rb*], *rd*(15:8) ← 0, *rd*(24:16) ← 0, *rb*(23:0) ← *rb*(23:0) + *imm24*

ld.ub *%rd*, [*%rb*]- (with post-decrement option)

Standard) *rd*(7:0) ← B[*rb*], *rd*(15:8) ← 0, *rd*(23:16) ← 0, *rb*(23:0) ← *rb*(23:0) - 1
Extension 1) *rd*(7:0) ← B[*rb*], *rd*(15:8) ← 0, *rd*(24:16) ← 0, *rb*(23:0) ← *rb*(23:0) - *imm13*
Extension 2) *rd*(7:0) ← B[*rb*], *rd*(15:8) ← 0, *rd*(24:16) ← 0, *rb*(23:0) ← *rb*(23:0) - *imm24*

ld.ub *%rd*, -[*%rb*] (with pre-decrement option)

Standard) *rb*(23:0) ← *rb*(23:0) - 1, *rd*(7:0) ← B[*rb*], *rd*(15:8) ← 0, *rd*(23:16) ← 0
Extension 1) *rb*(23:0) ← *rb*(23:0) - *imm13, rd*(7:0) ← B[*rb*], *rd*(15:8) ← 0, *rd*(24:16) ← 0
Extension 2) *rb*(23:0) ← *rb*(23:0) - *imm24, rd*(7:0) ← B[*rb*], *rd*(15:8) ← 0, *rd*(24:16) ← 0

Code

15	14	13	12	11	10	9–7	6	5	4	3	2–0	
0	0	1	0	0	0	rd	0	1	0	1	rb	`ld.ub %rd,[%rb]+`
0	0	1	0	0	0	rd	1	1	0	1	rb	`ld.ub %rd,[%rb]-`
0	0	1	0	0	0	rd	1	0	0	1	rb	`ld.ub %rd,-[%rb]`

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register indirect `%rb` = `%r0` to `%r7`
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

Two cycles

Description

(1) Address increment/decrement option

Specifying the `[]+`, `[]-`, or `-[]` option will automatically increment/decrement the memory address. This allows the program to simply perform continuous data transfer.

`ld.ub %rd,[%rb]+` Load instruction with post-increment option
The memory address will be incremented after the data transfer has finished.

`ld.ub %rd,[%rb]-` Load instruction with post-decrement option
The memory address will be decremented after the data transfer has finished.

`ld.ub %rd,-[%rb]` Load instruction with pre-decrement option
The memory address will be decremented before starting the data transfer.

The address increment/decrement sizes are listed below.
When no `ext` is used (standard): 1 (byte size)
When one `ext` is used (extension 1): *imm13*
When two `ext` are used (extension 2): *imm24*

(2) Standard (example of post-increment option)

```
ld.ub   %rd,[%rb]+ ; source memory address = rb
                   ; post increment: rb + 1
```

The byte data in the specified memory location is transferred to the *rd* register after being zero-extended to 16 bits. The *rb* register contains the memory address to be accessed. The eight high-order bits of the *rd* register are set to 0. The memory address will be incremented by two bytes after the data transfer has finished.

(3) Extension 1 (example of post-decrement option)

```
ext     imm13
ld.ub   %rd,[%rb]- ; source memory address = rb
                   ; post decrement: rb - imm13
```

The byte data in the specified memory location is transferred to the *rd* register after being zero-extended to 16 bits. The *rb* register contains the memory address to be accessed. The eight high-order bits of the *rd* register are set to 0. The memory address will be decremented by *imm13* bytes after the data transfer has finished.

(4) Extension 2 (example of pre-decrement option)

```
ext     imm11      ; imm11(10:0) = imm24(23:13)
ext     imm13      ; = imm24(12:0)
ld.ub   %rd,-[%rb] ; source memory address = rb - imm24
```

After the memory address specified by the *rb* register is decremented by *imm24* bytes, the byte data in the decremented address is transferred to the *rd* register after being zero-extended to 16 bits. The eight high-order bits of the *rd* register are set to 0.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

ld.ub *%rd*, [%sp + *imm7*]

Function

Unsigned byte data transfer

Standard) *rd*(7:0) ← B[sp + *imm7*], *rd*(15:8) ← 0, *rd*(23:16) ← 0
Extension 1) *rd*(7:0) ← B[sp + *imm20*], *rd*(15:8) ← 0, *rd*(23:16) ← 0
Extension 2) *rd*(7:0) ← B[sp + *imm24*], *rd*(15:8) ← 0, *rd*(23:16) ← 0

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	0	0	1	*rd*			*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register indirect with displacement
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

Two cycles

Description

(1) Standard

```
ld.ub   %rd,[%sp + imm7]     ; memory address = sp + imm7
```

The byte data in the specified memory location is transferred to the *rd* register after being zero-extended to 16 bits. The content of the current SP with the 7-bit immediate *imm7* added as displacement comprises the memory address to be accessed. The eight high-order bits of the *rd* register are set to 0.

(2) Extension 1

```
ext     imm13                ; = imm20(19:7)
ld.ub   %rd,[%sp + imm7]     ; memory address = sp + imm20,
                             ; imm7 = imm20(6:0)
```

The `ext` instruction extends the displacement to a 20-bit quantity. As a result, the content of the SP with the 20-bit immediate *imm20* added comprises the memory address, the byte data in which is transferred to the *rd* register after being zero-extended to 16 bits. The eight high-order bits of the *rd* register are set to 0.

(3) Extension 2

```
ext     imm4                 ; imm4(3:0) = imm24(23:20)
ext     imm13                ; = imm24(19:7)
ld.ub   %rd,[%sp + imm7]     ; memory address = sp + imm24,
                             ; imm7 ← imm24(6:0)
```

The two `ext` instructions extend the displacement to a 24-bit quantity. As a result, the content of the SP with the 24-bit immediate *imm24* added comprises the memory address, the byte data in which is transferred to the *rd* register after being zero-extended to 16 bits. The eight high-order bits of the *rd* register are set to 0.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

```
ext     0x1
ld.ub   %r0,[%sp + 0x1] ; r0 ← [sp + 0x81] zero-extended
```

ld.ub *%rd*, [*imm7*]

Function Unsigned byte data transfer

Standard) *rd*(7:0) ← B[*imm7*], *rd*(15:8) ← 0, *rd*(23:16) ← 0
Extension 1) *rd*(7:0) ← B[*imm20*], *rd*(15:8) ← 0, *rd*(23:16) ← 0
Extension 2) *rd*(7:0) ← B[*imm24*], *rd*(15:8) ← 0, *rd*(23:16) ← 0

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	0	0	0	1		*rd*					*imm7*			

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode Src: Immediate data (unsigned)
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK One cycle

Description (1) Standard

```
ld.ub   %rd,[imm7]    ; memory address = imm7
```

The byte data in the memory address specified with the 7-bit immediate *imm7* is transferred to the *rd* register after being zero-extended to 16 bits. The eight high-order bits of the *rd* register are set to 0.

(2) Extension 1

```
ext     imm13         ; = imm20(19:7)
ld.ub   %rd,[imm7]    ; memory address = imm20, imm7 = imm20(6:0)
```

The `ext` instruction extends the displacement to a 20-bit quantity. As a result, the byte data in the memory address specified with the 20-bit immediate *imm20* is transferred to the *rd* register after being zero-extended to 16 bits. The eight high-order bits of the *rd* register are set to 0.

(3) Extension 2

```
ext     imm4          ; imm4(3:0) = imm24(23:20)
ext     imm13         ; = imm24(19:7)
ld.ub   %rd,[imm7]    ; memory address = imm24, imm7 ← imm24(6:0)
```

The two `ext` instructions extend the displacement to a 24-bit quantity. As a result, the byte data in the memory address specified with the 24-bit immediate *imm24* is transferred to the *rd* register after being zero-extended to 16 bits. The eight high-order bits of the *rd* register are set to 0.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

```
ext     0x1
ld.ub   %r0,[0x1] ; r0 ← [0x81] zero-extended
```

nop

Function

No operation

Standard) No operation
Extension 1) Unusable
Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

–

CLK

One cycle

Description

(1) Standard

The `nop` instruction just takes one cycle and no operation results. The PC is incremented (+2).

(2) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit.

Example

```
nop
nop             ; Waits 2 cycles
```

not %rd, %rs
not/c %rd, %rs
not/nc %rd, %rs

Function

16-bit logical negation

Standard) *rd*(15:0) ← !*rs*(15:0), *rd*(23:16) ← 0

Extension 1) *rd*(15:0) ← !*imm13*(zero extended), *rd*(23:16) ← 0

Extension 2) *rd*(15:0) ← !*imm16*, *rd*(23:16) ← 0

Code

15	14	13	12	11	10	9–7	6	5	4	3	2–0	
0	0	1	0	1	1	rd	1	0	1	1	rs	not
0	0	1	0	1	1	rd	0	0	1	1	rs	not/c
0	0	1	0	1	1	rd	0	1	1	1	rs	not/nc

Flag

IL	IE	C	V	Z	N
–	–	–	0	↔	↔

Mode

Src: Register direct %*rs* = %r0 to %r7

Dst: Register direct %*rd* = %r0 to %r7

CLK

One cycle

Description

(1) Standard

```
not   %rd,%rs          ; rd ← !rs
```

The low-order 16 bits of the *rs* register are reversed, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(2) Extension 1

```
ext   imm13
not   %rd,%rs          ; rd ← !imm13
```

All the bits of the zero-extended 13-bit immediate *imm13* are reversed after zero-extended into 16 bits, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(3) Extension 2

```
ext   imm3             ; imm3(2:0) = imm16(15:13)
ext   imm13            ; = imm16(12:0)
not   %rd,%rs          ; rd ← !imm16
```

All the bits of the 16-bit immediate *imm16* are reversed, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(4) Conditional execution

The `/c` or `/nc` suffix on the opcode specifies conditional execution.

`not/c` Executed as `not` when the C flag is 1 or executed as `nop` when the flag is 0

`not/nc` Executed as `not` when the C flag is 0 or executed as `nop` when the flag is 1

In this case, the `ext` instruction can be used to extend the operand.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

When r1 = 0x555555

```
not  %r0,%r1    ; r0 = 0x00aaaa
```

not %rd, sign7

Function

16-bit logical negation

Standard) *rd*(15:0) ← !*sign7*(sign extended), *rd*(23:16) ← 0

Extension 1) *rd*(15:0) ← !*sign16*, *rd*(23:16) ← 0

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	1	1	*rd*			*sign7*						

Flag

IL	IE	C	V	Z	N
–	–	–	0	↔	↔

Mode

Src: Immediate data (signed)

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
not  %rd,sign7      ; rd ← !sign7
```

All the bits of the sign-extended 7-bit immediate *sign7* are reversed after sign-extended into 16 bits, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(2) Extension 1

```
ext  imm9           ; imm9(8:0) = sign16(15:7)
not  %rd,sign7      ; rd ← !sign16, sign7 = sign16(6:0)
```

All the bits of the sign-extended 16-bit immediate *sign16* are reversed, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(3) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
not  %r0,0x3f       ; r0 = 0x00ffc0
```

(2)
```
ext  0x1ff
not  %r1,0x7f       ; r1 = 0x000000
```

or %rd, %rs
or/c %rd, %rs
or/nc %rd, %rs

Function

16-bit logical OR

Standard) *rd*(15:0) ← *rd*(15:0) | *rs*(15:0), *rd*(23:16) ← 0

Extension 1) *rd*(15:0) ← *rs*(15:0) | *imm13*(zero extended), *rd*(23:16) ← 0

Extension 2) *rd*(15:0) ← *rs*(15:0) | *imm16*, *rd*(23:16) ← 0

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	1	0	1	1	rd			1	0	0	1	rs			or
0	0	1	0	1	1	rd			0	0	0	1	rs			or/c
0	0	1	0	1	1	rd			0	1	0	1	rs			or/nc

Flag

IL	IE	C	V	Z	N
–	–	–	0	↔	↔

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
or    %rd,%rs          ; rd ← rd | rs
```

The content of the *rs* register and that of the *rd* register are logically OR'ed, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(2) Extension 1

```
ext   imm13
or    %rd,%rs          ; rd ← rs | imm13
```

The content of the *rs* register and the zero-extended 13-bit immediate *imm13* are logically OR'ed, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0. The content of the *rs* register is not altered.

(3) Extension 2

```
ext   imm3             ; imm3(2:0)  = imm16(15:13)
ext   imm13            ; = imm16(12:0)
or    %rd,%rs          ; rd ← rs | imm16
```

The content of the *rs* register and the zero-extended 16-bit immediate *imm16* are logically OR'ed, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0. The content of the *rs* register is not altered.

(4) Conditional execution

The `/c` or `/nc` suffix on the opcode specifies conditional execution.

`or/c` Executed as `or` when the C flag is 1 or executed as `nop` when the flag is 0

`or/nc` Executed as `or` when the C flag is 0 or executed as `nop` when the flag is 1

In this case, the `ext` instruction can be used to extend the operand.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
or    %r0,%r0          ; r0 = r0 | r0
```

(2)
```
ext   0x1
ext   0x1fff
or    %r1,%r2          ; r1 = r2 | 0x3fff
```

or %rd, sign7

Function

16-bit logical OR

Standard) *rd*(15:0) ← *rd*(15:0) | *sign7*(sign extended), *rd*(23:16) ← 0

Extension 1) *rd*(15:0) ← *rd*(15:0) | *sign16*, *rd*(23:16) ← 0

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	0	1	*rd*			*sign7*						

Flag

IL	IE	C	V	Z	N
–	–	–	0	↔	↔

Mode

Src: Immediate data (signed)

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
or   %rd,sign7        ; rd ← rd | sign7
```

The content of the *rd* register and the sign-extended 7-bit immediate *sign7* are logically OR'ed, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(2) Extension 1

```
ext  imm9             ; imm9(8:0) = sign16(15:7)
or   %rd,sign7        ; rd ← rd | sign16, sign7 = sign16(6:0)
```

The content of the *rd* register and the 16-bit immediate *sign16* are logically OR'ed, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(3) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
or   %r0,0x7e         ; r0 = r0 | 0xfffe
```

(2)
```
ext  0xff
or   %r1,0x7f         ; r1 = r1 | 0x7fff
```

ret
ret.d

Function

Return from subroutine

Standard) pc ← A[sp](23:0), sp ← sp + 4
Extension 1) Unusable
Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	0	0	0	0	0	1	0	0	1	0	0	0	0	0	`ret`
0	0	0	0	0	0	0	1	1	0	1	0	0	0	0	0	`ret.d`

Flag

IE	C	V	Z	N
–	–	–	–	–

Mode

–

CLK

`ret` Three cycles
`ret.d` Two cycles (when a one-cycle delayed slot instruction follows), Three cycles (other)

Description

(1) Standard

`ret`

Restores the PC value (return address) that was saved into the stack when the `call`/`calla` instruction was executed for returning the program flow from the subroutine to the routine that called the subroutine. The SP is incremented by four bytes.
If the SP has been modified in the subroutine, it is necessary to return the SP value before executing the `ret` instruction.

(2) Delayed branch (d bit (bit 7) = 1)

`ret.d`

For the `ret.d` instruction, the next instruction becomes a delayed slot instruction. A delayed slot instruction is executed before the program returns from the subroutine. Interrupts are masked in intervals between the `ret.d` instruction and the next instruction, so no interrupts occur.

Example

```
ret.d
add   %r0,%r1    ; Executed before return from the subroutine
```

Caution

When the `ret.d` instruction (delayed branch) is used, be careful to ensure that the next instruction is limited to those that can be used as a delayed slot instruction. If any other instruction is executed, the program may operate indeterminately. For the usable instructions, refer to the instruction list in the Appendix.

retd

Function

Return from a debug-interrupt handler routine

Standard) r0 ← A[DBRAM + 0x4](23:0), {psr, pc} ← A[DBRAM + 0x0]
Extension 1) Unusable
Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	0	0	1	0	1	1	0	1	0	0	0

Flag

IL	IE	C	V	Z	N
↔	↔	↔	↔	↔	↔

Mode

–

CLK

Four cycles

Description

Restore the contents of the R0, PSR, and PC that were saved to the work area for debugging (DBRAM) when an debug interrupt occurred to the respective registers, and return from the debug interrupt handler routine.
This instruction is provided for debug firmware. Do not use it in the user program.

Example

```
retd            ; Return from a debug interrupt handler routine
```

reti
reti.d

Function

Return from interrupt handler routine

Standard) {psr, pc} ← A[sp], sp ← sp + 4
Extension 1) Unusable
Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	0	0	0	0	0	1	0	0	1	0	1	0	0	0	reti
0	0	0	0	0	0	0	1	1	0	1	0	1	0	0	0	reti.d

Flag

IL	IE	C	V	Z	N
↔	↔	↔	↔	↔	↔

Mode

–

CLK

`reti` Three cycles
`reti.d` Two cycles (when a one-cycle delayed slot instruction follows), Three cycles (other)

Description

(1) Standard

`reti`

Restores the contents of the PC and PSR that were saved to the stack when an interrupt occurred to the respective registers, and return from the interrupt handler routine. The SP is incremented by an amount equivalent to four bytes.

(2) Delayed branch (d bit (bit 7) = 1)

`reti.d`

For the `reti.d` instruction, the next instruction becomes a delayed slot instruction. A delayed slot instruction is executed before the program returns from the interrupt handler routine. Interrupts are masked in intervals between the `reti.d` instruction and the next instruction, so no interrupts occur.

Example

```
reti                ; Return from a interrupt handler routine
```

sa %rd, %rs

Function

Arithmetic shift to the right

Standard) Shift the content of *rd* to right as many bits as specified by *rs* (0–3, 4, or 8 bits), MSB ← MSB (sign bit)

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	1	1		*rd*		1	1	0	1		*rs*	

Flag

IL	IE	C	V	Z	N
–	–	↔	–	↔	↔

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

The *rd* register is shifted as shown in the diagram below.

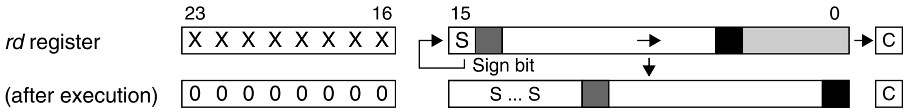
The number of bits to be shifted is specified by the *rs* register value as follows:

rs = 0–3: 0–3 bits

rs = 4–7: 4 bits

rs = 8 or more: 8 bits

The sign bit is copied to bit 15 of the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.



(2) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit included.

sa %rd, imm7

Function

Arithmetic shift to the right

Standard) Shift the content of *rd* to right as many bits as specified by *imm7* (0–3, 4, or 8 bits), MSB ← MSB (sign bit)

Extension 1) *imm7* is extended to *imm20*

Extension 2) *imm7* is extended to *imm24*

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	0	1	*rd*			*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	↔	–	↔	↔

Mode

Src: Immediate (unsigned)

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

The *rd* register is shifted as shown in the diagram below.

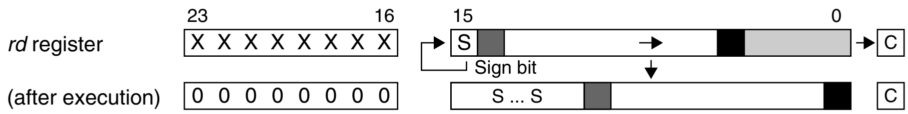
The number of bits to be shifted is specified by the 7-bit immediate *imm7* as follows:

imm7 = 0–3: 0–3 bits

imm7 = 4–7: 4 bits

imm7 = 8 or more: 8 bits

The sign bit is copied to the most significant bit of the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.



(2) Extension

Using the `ext` instruction extends the 7-bit immediate *imm7* to 20-bit immediate *imm20* or 24-bit immediate *imm24*. However, there is no difference in operation from the standard instruction without extension.

(3) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit included. In this case, extension of the immediate by the `ext` instruction cannot be performed.

sbc %rd, %rs
sbc/c %rd, %rs
sbc/nc %rd, %rs

Function

16-bit subtraction with borrow

Standard) *rd*(15:0) ← *rd*(15:0) - *rs*(15:0) - C, *rd*(23:16) ← 0

Extension 1) *rd*(15:0) ← *rs*(15:0) - *imm13*(zero extended) - C, *rd*(23:16) ← 0

Extension 2) *rd*(15:0) ← *rs*(15:0) - *imm16* - C, *rd*(23:16) ← 0

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
0	0	1	1	1	0	rd			1	0	1	1	rs			sbc
0	0	1	1	1	0	rd			0	0	1	1	rs			sbc/c
0	0	1	1	1	0	rd			0	1	1	1	rs			sbc/nc

Flag

IL	IE	C	V	Z	N	
–	–	↔	↔	↔	↔	sbc
–	–	–	↔	↔	↔	sbc/c, sbc/nc

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
sbc  %rd,%rs          ; rd ← rd - rs - C
```

The content of the *rs* register and C (carry) flag are subtracted from the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(2) Extension 1

```
ext  imm13
sbc  %rd,%rs          ; rd ← rs - imm13 - C
```

The 13-bit immediate *imm13* and C (carry) flag are subtracted from the *rs* register after being zero-extended, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0. The content of the *rs* register is not altered.

(3) Extension 2

```
ext  imm3             ; imm3(2:0) = imm16(15:13)
ext  imm13            ; = imm16(12:0)
sbc  %rd,%rs          ; rd ← rs - imm16 - C
```

The 16-bit immediate *imm16* and C (carry) flag are subtracted from the *rs* register, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0. The content of the *rs* register is not altered.

(4) Conditional execution

The `/c` or `/nc` suffix on the opcode specifies conditional execution.

`sbc/c` Executed as `sbc` when the C flag is 1 or executed as `nop` when the flag is 0

`sbc/nc` Executed as `sbc` when the C flag is 0 or executed as `nop` when the flag is 1

In this case, the `ext` instruction can be used to extend the operand.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
sbc  %r0,%r1          ; r0 = r0 - r1 - C
```

(2) Subtraction of 32-bit data

data 1 = {r2, r1}, data2 = {r4, r3}, result = {r2, r1}

```
sub  %r1,%r3          ; Subtraction of the low-order word
sbc  %r2,%r4          ; Subtraction of the high-order word
```

sbc %rd, imm7

Function

16-bit subtraction with borrow

Standard) *rd*(15:0) ← *rd*(15:0) - *imm7*(zero extended) - C, *rd*(23:16) ← 0
Extension 1) *rd*(15:0) ← *rd*(15:0) - *imm16* - C, *rd*(23:16) ← 0
Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	1	1		*rd*					*imm7*			

Flag

IL	IE	C	V	Z	N
–	–	↔	↔	↔	↔

Mode

Src: Immediate data (unsigned)
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
sbc   %rd,imm7          ; rd ← rd - imm7 - C
```

The 7-bit immediate *imm7* and C (carry) flag are subtracted from the *rd* register after being zero-extended. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(2) Extension 1

```
ext   imm9              ; imm9(8:0) = imm16(15:7)
sbc   %rd,imm7          ; rd ← rd - imm16 - C, imm7 = imm16(6:0)
```

The 16-bit immediate *imm16* and C (carry) flag are subtracted from the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(3) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
sbc   %r0,0x7f          ; r0 = r0 - 0x7f - C
```

(2)
```
ext   0x1ff
sbc   %r1,0x7f          ; r1 = r1 - 0xffff - C
```

sl %rd, %rs

Function

Logical shift to the left

Standard) Shift the content of *rd* to left as many bits as specified by *rs* (0–3, 4, or 8 bits), LSB ← 0

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	1	1		*rd*		1	1	1	0		*rs*	

Flag

IL	IE	C	V	Z	N
–	–	↔	–	↔	↔

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

The *rd* register is shifted as shown in the diagram below.

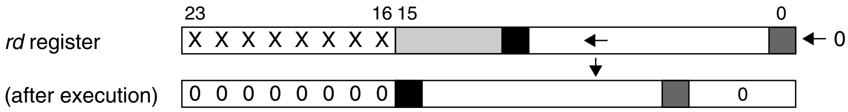
The number of bits to be shifted is specified by the *rs* register value as follows:

rs = 0–3: 0–3 bits

rs = 4–7: 4 bits

rs = 8 or more: 8 bits

Data “0” is placed in the least significant bit of the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.



(2) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit included.

sl *%rd, imm7*

Function

Logical shift to the left

Standard) Shift the content of *rd* to left as many bits as specified by *imm7* (0–3, 4, or 8 bits), LSB ← 0

Extension 1) *imm7* is extended to *imm20*

Extension 2) *imm7* is extended to *imm24*

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	1	0	*r d*			*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	↔	–	↔	↔

Mode

Src: Immediate (unsigned)

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

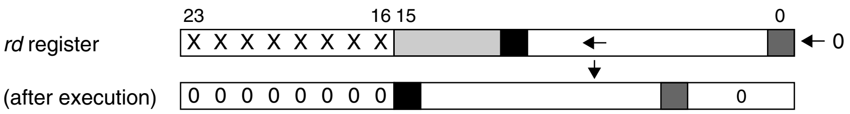
(1) Standard

The *rd* register is shifted as shown in the diagram below.

The number of bits to be shifted is specified by the 7-bit immediate *imm7* as follows:

imm7 = 0–3: 0–3 bits
imm7 = 4–7: 4 bits
imm7 = 8 or more: 8 bits

Data "0" is placed in the least significant bit of the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.



(2) Extension

Using the `ext` instruction extends the 7-bit immediate *imm7* to 20-bit immediate *imm20* or 24-bit immediate *imm24*. However, there is no difference in operation from the standard instruction without extension.

(3) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit included. In this case, extension of the immediate by the `ext` instruction cannot be performed.

slp

Function

SLEEP

Standard) Place the processor in SLEEP mode
Extension 1) Unusable
Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	0	0	0	0	0	0	1	0	0	0	0

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode –

CLK Six cycles

Description

Places the processor in SLEEP mode for power saving.
Program execution is halted at the same time that the S1C17 Core executes the `slp` instruction, and the processor enters SLEEP mode.
SLEEP mode commonly turns off the S1C17 Core and on-chip peripheral circuit operations, thereby it significantly reduces the current consumption in comparison to HALT mode.

Initial reset is one cause that can bring the processor out of SLEEP mode. Other causes depend on the implementation of the clock control circuit outside the S1C17 Core.
Initial reset, maskable external interrupts, NMI, and debug interrupts are commonly used for canceling SLEEP mode.
The interrupt enable/disable status set in the processor does not affect the cancellation of SLEEP mode even if an interrupt signal is used as the cancellation. In other words, interrupt signals are able to cancel SLEEP mode even if the IE flag in PSR or the interrupt enable bits in the interrupt controller (depending on the implementation) are set to disable interrupts.
When the processor is taken out of SLEEP mode using an interrupt that has been enabled (by the interrupt controller and IE flag), the corresponding interrupt handler routine is executed. Therefore, when the interrupt handler routine is terminated by the `reti` instruction, the processor returns to the instruction next to `slp`.
When the interrupt has been disabled, the processor restarts the program from the instruction next to `slp` after the processor is taken out of SLEEP mode.

Refer to the technical manual of each model for details of SLEEP mode.

Example

```
slp       ; The processor is placed in SLEEP mode.
```

sr %rd, %rs

Function

Logical shift to the right

Standard) Shift the content of *rd* to right as many bits as specified by *rs* (0–3, 4, or 8 bits), MSB ← 0

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	1	1		*rd*		1	1	0	0		*rs*	

Flag

IL	IE	C	V	Z	N
–	–	↔	–	↔	↔

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

The *rd* register is shifted as shown in the diagram below.

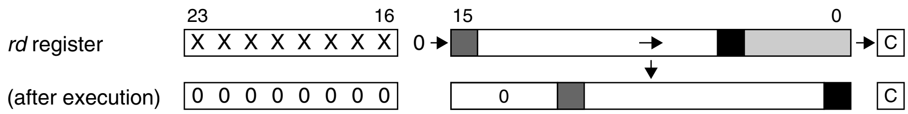
The number of bits to be shifted is specified by the *rs* register value as follows:

rs = 0–3: 0–3 bits

rs = 4–7: 4 bits

rs = 8 or more: 8 bits

Data “0” is placed in the bit 15 of the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.



(2) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit included.

sr %rd, imm7

Function

Logical shift to the right

Standard) Shift the content of *rd* to right as many bits as specified by *imm7* (0–3, 4, or 8 bits), MSB ← 0

Extension 1) *imm7* is extended to *imm20*

Extension 2) *imm7* is extended to *imm24*

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	1	0	0	*r d*			*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	↔	–	↔	↔

Mode

Src: Immediate (unsigned)

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

The *rd* register is shifted as shown in the diagram below.

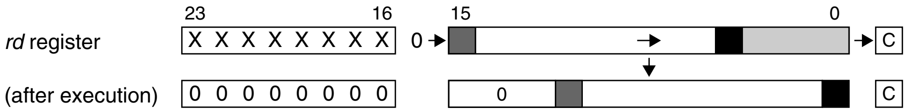
The number of bits to be shifted is specified by the 7-bit immediate *imm7* as follows:

imm7 = 0–3: 0–3 bits

imm7 = 4–7: 4 bits

imm7 = 8 or more: 8 bits

Data “0” is placed in the bit 15 of the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.



(2) Extension

Using the `ext` instruction extends the 7-bit immediate *imm7* to 20-bit immediate *imm20* or 24-bit immediate *imm24*. However, there is no difference in operation from the standard instruction without extension.

(3) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit included. In this case, extension of the immediate by the `ext` instruction cannot be performed.

sub %rd, %rs
sub/c %rd, %rs
sub/nc %rd, %rs

Function

16-bit subtraction

Standard) *rd*(15:0) ← *rd*(15:0) - *rs*(15:0), *rd*(23:16) ← 0

Extension 1) *rd*(15:0) ← *rs*(15:0) - *imm13*(zero extended), *rd*(23:16) ← 0

Extension 2) *rd*(15:0) ← *rs*(15:0) - *imm16*, *rd*(23:16) ← 0

Code

15	14	13	12	11	10	9–7	6	5	4	3	2–0	
0	0	1	1	1	0	rd	1	0	1	0	rs	sub
0	0	1	1	1	0	rd	0	0	1	0	rs	sub/c
0	0	1	1	1	0	rd	0	1	1	0	rs	sub/nc

Flag

IL	IE	C	V	Z	N	
–	–	↔	↔	↔	↔	sub
–	–	–	↔	↔	↔	sub/c, sub/nc

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
sub   %rd,%rs          ; rd ← rd - rs
```

The content of the *rs* register is subtracted from the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(2) Extension 1

```
ext   imm13
sub   %rd,%rs          ; rd ← rs - imm13
```

The 13-bit immediate *imm13* is subtracted from the *rs* register after being zero-extended, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0. The content of the *rs* register is not altered.

(3) Extension 2

```
ext   imm3             ; imm3(2:0)  = imm16(15:13)
ext   imm13            ; = imm16(12:0)
sub   %rd,%rs          ; rd ← rs - imm16
```

The 16-bit immediate *imm16* is subtracted from the *rs* register, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0. The content of the *rs* register is not altered.

(4) Conditional execution

The `/c` or `/nc` suffix on the opcode specifies conditional execution.

`sub/c` Executed as `sub` when the C flag is 1 or executed as `nop` when the flag is 0

`sub/nc` Executed as `sub` when the C flag is 0 or executed as `nop` when the flag is 1

In this case, the `ext` instruction can be used to extend the operand.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
sub   %r0,%r0          ; r0 = r0 - r0
```

(2)
```
ext   0x1
ext   0x1fff
sub   %r1,%r2          ; r1 = r2 - 0x3fff
```

sub *%rd, imm7*

Function

16-bit subtraction

Standard) *rd*(15:0) ← *rd*(15:0) - *imm7*(zero extended), *rd*(23:16) ← 0

Extension 1) *rd*(15:0) ← *rd*(15:0) - *imm16*, *rd*(23:16) ← 0

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	1	0	*rd*			*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	↔	↔	↔	↔

Mode

Src: Immediate data (unsigned)

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
sub  %rd,imm7         ; rd ← rd - imm7
```

The 7-bit immediate *imm7* is subtracted from the *rd* register after being zero-extended. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(2) Extension 1

```
ext  imm9             ; imm9(8:0) = imm16(15:7)
sub  %rd,imm7         ; rd ← rd - imm16, imm7 = imm16(6:0)
```

The 16-bit immediate *imm16* is subtracted from the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(3) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
sub  %r0,0x3f         ; r0 = r0 - 0x3f
```

(2)
```
ext  0x1ff
sub  %r1,0x7f         ; r1 = r1 - 0xffff
```

sub.a *%rd, %rs*
sub.a/c *%rd, %rs*
sub.a/nc *%rd, %rs*

Function

24-bit subtraction

Standard) *rd*(23:0) ← *rd*(23:0) - *rs*(23:0)
Extension 1) *rd*(23:0) ← *rs*(23:0) - *imm13*(zero extended)
Extension 2) *rd*(23:0) ← *rs*(23:0) - *imm24*

Code

15	14	13	12	11	10	9–7	6	5	4	3	2–0	
0	0	1	1	0	0	rd	1	0	1	0	rs	sub.a
0	0	1	1	0	0	rd	0	0	1	0	rs	sub.a/c
0	0	1	1	0	0	rd	0	1	1	0	rs	sub.a/nc

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
sub.a   %rd,%rs        ; rd ← rd - rs
```

The content of the *rs* register is subtracted from the *rd* register.

(2) Extension 1

```
ext     imm13
sub.a   %rd,%rs        ; rd ← rs - imm13
```

The 13-bit immediate *imm13* is subtracted from the content of the *rs* register after being zero-extended, and the result is loaded into the *rd* register. The content of the *rs* register is not altered.

(3) Extension 2

```
ext     imm11          ; imm11(10:0) = imm24(23:13)
ext     imm13          ; = imm24(12:0)
sub.a   %rd,%rs        ; rd ← rs - imm24
```

The 24-bit immediate *imm24* is subtracted from the content of the *rs* register, and the result is loaded into the *rd* register. The content of the *rs* register is not altered.

(4) Conditional execution

The `/c` or `/nc` suffix on the opcode specifies conditional execution.

`sub.a/c` Executed as `sub.a` when the C flag is 1 or executed as `nop` when the flag is 0
`sub.a/nc` Executed as `sub.a` when the C flag is 0 or executed as `nop` when the flag is 1

In this case, the `ext` instruction can be used to extend the operand.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
sub.a   %r0,%r0        ; r0 = r0 - r0
```

(2)
```
ext     0x7ff
ext     0x1fff
sub.a   %r1,%r2        ; r1 = r2 - 0xffffff
```

sub.a *%rd, imm7*

Function

24-bit subtraction

Standard) *rd*(23:0) ← *rd*(23:0) - *imm7*(zero extended)
Extension 1) *rd*(23:0) ← *rd*(23:0) - *imm20*(zero extended)
Extension 2) *rd*(23:0) ← *rd*(23:0) - *imm24*

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	1	0	*rd*			*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Immediate data (unsigned)
Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
sub.a   %rd,imm7       ; rd ← rd - imm7
```

The 7-bit immediate *imm7* is subtracted from the *rd* register after being zero-extended.

(2) Extension 1

```
ext     imm13          ; = imm20(19:7)
sub.a   %rd,imm7       ; rd ← rd - imm20, imm7 = imm20(6:0)
```

The 20-bit immediate *imm20* is subtracted from the *rd* register after being zero-extended.

(3) Extension 2

```
ext     imm4           ; imm4(3:0) = imm24(23:20)
ext     imm13          ; = imm24(19:7)
sub.a   %rd,imm7       ; rd ← rd - imm24, imm7 = imm24(6:0)
```

The 24-bit immediate *imm24* is subtracted from the *rs* register.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the “d” bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
sub.a   %r0,0x7f       ; r0 = r0 - 0x7f
```

(2)
```
ext     0xf
ext     0x1fff
sub.a   %r1,0x7f       ; r1 = r1 - 0xffffff
```

sub.a %sp, *%rs*

Function

24-bit subtraction

Standard) sp(23:0) ← sp(23:0) - *rs*(23:0)

Extension 1) sp(23:0) ← *rs*(23:0) - *imm13*(zero extended)

Extension 2) sp(23:0) ← *rs*(23:0) - *imm24*

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	1	0	0	0	0	0	0	0	1	1	*rs*		

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register direct `%sp`

CLK

One cycle

Description

(1) Standard

```
sub.a   %sp,%rs        ; sp ← sp - rs
```

The content of the *rs* register is subtracted from the stack pointer SP.

(2) Extension 1

```
ext     imm13
sub.a   %sp,%rs        ; sp ← rs - imm13
```

The 13-bit immediate *imm13* is subtracted from the content of the *rs* register after being zero-extended, and the result is loaded into the stack pointer SP. The content of the *rs* register is not altered.

(3) Extension 2

```
ext     imm11          ; imm11(10:0) = imm24(23:13)
ext     imm13          ; = imm24(12:0)
sub.a   %sp,%rs        ; sp ← rs - imm24
```

The 24-bit immediate *imm24* is subtracted from the content of the *rs* register, and the result is loaded into the stack pointer SP. The content of the *rs* register is not altered.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
sub.a   %sp,%r0        ; sp = sp - r0
```

(2)
```
ext     0x1
ext     0x1ffc
sub.a   %sp,%r2        ; sp = r2 - 0x3ffc
```

Caution

The 2 low-order bits of the subtruction results are always loaded to the SP as 0.

sub.a %sp, *imm7*

Function

24-bit subtraction

Standard) sp(23:0) ← sp(23:0) - *imm7*(zero extended)
Extension 1) sp(23:0) ← sp(23:0) - *imm20*(zero extended)
Extension 2) sp(23:0) ← sp(23:0) - *imm24*

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	1	1	0	1	1	0	0	0	*imm7*						

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Immediate data (unsigned)
Dst: Register direct `%sp`

CLK

One cycle

Description

(1) Standard

```
sub.a   %sp,imm7      ; sp ← sp - imm7
```

The 7-bit immediate *imm7* is subtracted from the stack pointer SP after being zero-extended.

(2) Extension 1

```
ext     imm13         ; = imm20(19:7)
sub.a   %sp,imm7      ; sp ← sp - imm20, imm7 = imm20(6:0)
```

The 20-bit immediate *imm20* is subtracted from the stack pointer SP after being zero-extended.

(3) Extension 2

```
ext     imm4          ; imm4(3:0) = imm24(23:20)
ext     imm13         ; = imm24(19:7)
sub.a   %sp,imm7      ; sp ← sp - imm24, imm7 = imm24(6:0)
```

The 24-bit immediate *imm24* is subtracted from the stack pointer SP.

(4) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
sub.a   %sp,0x7c      ; sp = sp - 0x7c
```

(2)
```
ext     0x1fff
sub.a   %sp,0x7c      ; sp = sp - 0xffffc
```

Caution

The 2 low-order bits of the subtruction results are always loaded to the SP as 0.

swap *%rd, %rs*

Function

Swap

Standard) *rd*(15:8) ← *rs*(7:0), *rd*(7:0) ← *rs*(15:8), *rd*(23:16) ← 0

Extension 1) Unusable

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	1	0	1	1		*rd*		1	1	1	1		*rs*	

Flag

IL	IE	C	V	Z	N
–	–	–	–	–	–

Mode

Src: Register direct `%rs` = `%r0` to `%r7`

Dst: Register direct `%rd` = `%r0` to `%r7`

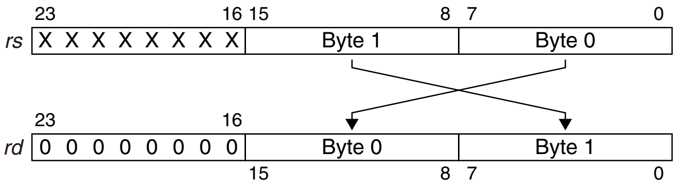
CLK

One cycle

Description

(1) Standard

Swaps the byte order of the 16 low-order bits of the *rs* register high and low and loads the results to the *rd* register.



(2) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit.

Example

When r1 = 0x123456

```
swap   %r2,%r1           ; 0x005634 → r2
```

xor %rd, %rs
xor/c %rd, %rs
xor/nc %rd, %rs

Function

16-bit exclusive OR

Standard) *rd*(15:0) ← *rd*(15:0) ^ *rs*(15:0), *rd*(23:16) ← 0

Extension 1) *rd*(15:0) ← *rs*(15:0) ^ *imm13*(zero extended), *rd*(23:16) ← 0

Extension 2) *rd*(15:0) ← *rs*(15:0) ^ *imm16*, *rd*(23:16) ← 0

Code

15	14	13	12	11	10	9–7	6	5	4	3	2–0	
0	0	1	0	1	1	rd	1	0	1	0	rs	xor
0	0	1	0	1	1	rd	0	0	1	0	rs	xor/c
0	0	1	0	1	1	rd	0	1	1	0	rs	xor/nc

Flag

IL	IE	C	V	Z	N
–	–	–	0	↔	↔

Mode

Src:Register direct `%rs` = `%r0` to `%r7`

Dst:Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
xor    %rd,%rs          ; rd ← rd ^ rs
```

The content of the *rs* register and that of the *rd* register are exclusively OR'ed, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(2) Extension 1

```
ext    imm13
xor    %rd,%rs          ; rd ← rs ^ imm13
```

The content of the *rs* register and the zero-extended 13-bit immediate *imm13* are exclusively OR'ed, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0. The content of the *rs* register is not altered.

(3) Extension 2

```
ext   imm3              ; imm3(2:0)  = imm16(15:13)
ext   imm13             ; = imm16(12:0)
xor   %rd,%rs           ; rd ← rs ^ imm16
```

The content of the *rs* register and the 16-bit immediate *imm16* are exclusively OR'ed, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0. The content of the *rs* register is not altered.

(4) Conditional execution

The `/c` or `/nc` suffix on the opcode specifies conditional execution.

`xor/c` Executed as `xor` when the C flag is 1 or executed as `nop` when the flag is 0

`xor/nc` Executed as `xor` when the C flag is 0 or executed as `nop` when the flag is 1

In this case, the `ext` instruction can be used to extend the operand.

(5) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
xor   %r0,%r0           ; r0 = r0 ^ r0
```

(2)
```
ext   0x1
ext   0x1fff
xor   %r1,%r2           ; r1 = r2 ^ 0x3fff
```

xor %rd, sign7

Function

16-bit exclusive OR

Standard) *rd*(15:0) ← *rd*(15:0) ^ *sign7*(sign extended), *rd*(23:16) ← 0

Extension 1) *rd*(15:0) ← *rd*(15:0) ^ *sign16*, *rd*(23:16) ← 0

Extension 2) Unusable

Code

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	1	0	*rd*			*sign7*						

Flag

IL	IE	C	V	Z	N
–	–	–	0	↔	↔

Mode

Src: Immediate data (signed)

Dst: Register direct `%rd` = `%r0` to `%r7`

CLK

One cycle

Description

(1) Standard

```
xor   %rd,sign7        ; rd ← rd ^ sign7
```

The content of the *rd* register and the sign-extended 7-bit immediate *sign7* are exclusively OR'ed, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(2) Extension 1

```
ext   imm9             ; imm9(8:0) = sign16(15:7)
xor   %rd,sign7        ; rd ← rd ^ sign16, sign7 = sign16(6:0)
```

The content of the *rd* register and the 16-bit immediate *sign16* are exclusively OR'ed, and the result is loaded into the *rd* register. The operation is performed in 16-bit size, and bits 23–16 of the *rd* register are set to 0.

(3) Delayed slot instruction

This instruction may be executed as a delayed slot instruction by writing it directly after a branch instruction with the "d" bit. In this case, extension of the immediate by the `ext` instruction cannot be performed.

Example

(1)
```
xor   %r0,0x7e         ; r0 = r0 ^ 0xfffe
```

(2)
```
ext   0x1ff
xor   %r1,0x7f         ; r1 = r1 ^ 0xffff
```

Appendix List of S1C17 Core Instructions

Symbols in the Instruction List

S1C17 Core Instruction Set

Registers/Register Data

%rd, rd:	A general-purpose register (R0–R7) used as the destination register or its contents
%rs, rs:	A general-purpose register (R0–R7) used as the source register or its contents
%rb, rb:	A general-purpose register (R0–R7) that has stored a base address to be accessed in the register indirect addressing mode or its contents
%sp, sp:	Stack pointer (SP) or its contents
%pc, pc:	Program counter (PC) or its contents

Memory/Addresses/Memory Data

[%rb], [%sp]:	Specification for register indirect addressing
[%rb]+, [%sp]+:	Specification for register indirect addressing with post-increment
[%rb]-, [%sp]-:	Specification for register indirect addressing with post-decrement
-[%rb], -[%sp]:	Specification for register indirect addressing with pre-decrement
[%sp+immX]:	Specification for register indirect addressing with a displacement
[imm7]:	Specification for a memory address with an immediate data
B[XXX]:	An address specified with XXX, or the byte data stored in the address
W[XXX]:	A 16-bit address specified with XXX, or the word data stored in the address
A[XXX]:	A 32-bit address specified with XXX, or the 24-bit or 32-bit data stored in the address

Immediate

immX:	A X-bit unsigned immediate data
signX:	A X-bit signed immediate data

Bit Field

(X):	Bit X of data
(X:Y):	A bit field from bit X to bit Y
{X, Y···}:	Indicates a bit (data) configuration.

Code

rd, rs, rb:	Register number (R0 = 0 ... R7 = 7)
d:	Delayed bit (0: Standard branch instruction, 1: Delayed branch instruction)

Functions

←:	Indicates that the right item is loaded or set to the left item.
+:	Addition
-:	Subtraction
&:	AND
|:	OR
^:	XOR
!:	NOT

Flags

IL:	Interrupt level
IE:	Interrupt enable flag
C:	Carry flag
V:	Overflow flag
Z:	Zero flag
N:	Negative flag
–:	Not changed
↔:	Set (1) or reset (0)
1:	Set (1)
0:	Reset (0)

EXT

*X:	Indicates that the operand can be extended (see the Remarks on each page for the extended operand).
–:	Indicates that the operand cannot be extended.

D

○:	Indicates that the instruction can be used as a delayed slot instruction.
–:	Indicates that the instruction cannot be used as a delayed slot instruction.

Data Transfer Instructions (1)

S1C17 Core Instruction Set

Mnemonic Opcode	Mnemonic Operand	Code MSB											Code LSB	Function	Cycle	IL	IE	C	V	Z	N	EXT	D
ld.b	%rd, %rs	0	0	1	0	1	0	rd	0	0	0	0	rs	rd(7:0)←rs(7:0), rd(15:8)←rs(7), rd(23:16)←0	1	–	–	–	–	–	–	–	○
	%rd, [%rb]	0	0	1	0	0	0	rd	0	0	0	0	rb	rd(7:0)←B[rb], rd(15:8)←B[rb](7), rd(23:16)←0	1, 2*7	–	–	–	–	–	–	*1	○
	%rd, [%rb]+	0	0	1	0	0	0	rd	0	1	0	0	rb	rd(7:0)←B[rb], rd(15:8)←B[rb](7), rd(23:16)←0, rb(23:0)←rb(23:0)+1	2	–	–	–	–	–	–	*6	○
	%rd, [%rb]-	0	0	1	0	0	0	rd	1	1	0	0	rb	rd(7:0)←B[rb], rd(15:8)←B[rb](7), rd(23:16)←0, rb(23:0)←rb(23:0)-1	2	–	–	–	–	–	–	*6	○
	%rd, -[%rb]	0	0	1	0	0	0	rd	1	0	0	0	rb	rb(23:0)←rb(23:0)-1, rd(7:0)←B[rb], rd(15:8)←B[rb](7), rd(23:16)←0	2	–	–	–	–	–	–	*6	○
	%rd, [%sp+imm7]	1	1	1	0	0	0	rd	imm7					rd(7:0)←B[sp+imm7], rd(15:8)←B[sp+imm7](7), rd(23:16)←0	2	–	–	–	–	–	–	*5	○
	%rd, [imm7]	1	1	0	0	0	0	rd	imm7					rd(7:0)←B[imm7], rd(15:8)←B[imm7](7), rd(23:16)←0	1	–	–	–	–	–	–	*4	○
	[%rb], %rs	0	0	1	0	0	1	rs	0	0	0	0	rb	B[rb]←rs(7:0)	1, 2*7	–	–	–	–	–	–	*1	○
	[%rb]+, %rs	0	0	1	0	0	1	rs	0	1	0	0	rb	B[rb]←rs(7:0), rb(23:0)←rb(23:0)+1	2	–	–	–	–	–	–	*6	○
	[%rb]-, %rs	0	0	1	0	0	1	rs	1	1	0	0	rb	B[rb]←rs(7:0), rb(23:0)←rb(23:0)-1	2	–	–	–	–	–	–	*6	○
	-[%rb], %rs	0	0	1	0	0	1	rs	1	0	0	0	rb	rb(23:0)←rb(23:0)-1, B[rb]←rs(7:0)	2	–	–	–	–	–	–	*6	○
	[%sp+imm7], %rs	1	1	1	1	0	0	rs	imm7					B[sp+imm7]←rs(7:0)	2	–	–	–	–	–	–	*5	○
	[imm7], %rs	1	1	0	1	0	0	rs	imm7					B[imm7]←rs(7:0)	1	–	–	–	–	–	–	*4	○
ld.ub	%rd, %rs	0	0	1	0	1	0	rd	0	0	0	1	rs	rd(7:0)←rs(7:0), rd(15:8)←0, rd(23:16)←0	1	–	–	–	–	–	–	–	○
	%rd, [%rb]	0	0	1	0	0	0	rd	0	0	0	1	rb	rd(7:0)←B[rb], rd(15:8)←0, rd(23:16)←0	1, 2*7	–	–	–	–	–	–	*1	○
	%rd, [%rb]+	0	0	1	0	0	0	rd	0	1	0	1	rb	rd(7:0)←B[rb], rd(15:8)←0, rd(23:16)←0, rb(23:0)←rb(23:0)+1	2	–	–	–	–	–	–	*6	○
	%rd, [%rb]-	0	0	1	0	0	0	rd	1	1	0	1	rb	rd(7:0)←B[rb], rd(15:8)←0, rd(23:16)←0, rb(23:0)←rb(23:0)-1	2	–	–	–	–	–	–	*6	○
	%rd, -[%rb]	0	0	1	0	0	0	rd	1	0	0	1	rb	rb(23:0)←rb(23:0)-1, rd(7:0)←B[rb], rd(15:8)←0, rd(23:16)←0	2	–	–	–	–	–	–	*6	○
	%rd, [%sp+imm7]	1	1	1	0	0	1	rd	imm7					rd(7:0)←B[sp+imm7], rd(15:8)←0, rd(23:16)←0	2	–	–	–	–	–	–	*5	○
	%rd, [imm7]	1	1	0	0	0	1	rd	imm7					rd(7:0)←B[imm7], rd(15:8)←0, rd(23:16)←0	1	–	–	–	–	–	–	*4	○
ld	%rd, %rs	0	0	1	0	1	0	rd	0	0	1	0	rs	rd(15:0)←rs(15:0), rd(23:16)←0	1	–	–	–	–	–	–	–	○
	%rd, sign7	1	0	0	1	1	0	rd	sign7					rd(6:0)←sign7(6:0), rd(15:7)←sign7(6), rd(23:16)←0	1	–	–	–	–	–	–	*2	○
	%rd, [%rb]	0	0	1	0	0	0	rd	0	0	1	0	rb	rd(15:0)←W[rb], rd(23:16)←0	1, 2*7	–	–	–	–	–	–	*1	○
	%rd, [%rb]+	0	0	1	0	0	0	rd	0	1	1	0	rb	rd(15:0)←W[rb], rd(23:16)←0, rb(23:0)←rb(23:0)+2	2	–	–	–	–	–	–	*6	○
	%rd, [%rb]-	0	0	1	0	0	0	rd	1	1	1	0	rb	rd(15:0)←W[rb], rd(23:16)←0, rb(23:0)←rb(23:0)-2	2	–	–	–	–	–	–	*6	○
	%rd, -[%rb]	0	0	1	0	0	0	rd	1	0	1	0	rb	rb(23:0)←rb(23:0)-2, rd(15:0)←W[rb], rd(23:16)←0	2	–	–	–	–	–	–	*6	○
	%rd, [%sp+imm7]	1	1	1	0	1	0	rd	imm7					rd(15:0)←W[sp+imm7], rd(23:16)←0	2	–	–	–	–	–	–	*5	○
	%rd, [imm7]	1	1	0	0	1	0	rd	imm7					rd(15:0)←W[imm7], rd(23:16)←0	1	–	–	–	–	–	–	*4	○
	[%rb], %rs	0	0	1	0	0	1	rs	0	0	1	0	rb	W[rb]←rs(15:0)	1, 2*7	–	–	–	–	–	–	*1	○
	[%rb]+, %rs	0	0	1	0	0	1	rs	0	1	1	0	rb	W[rb]←rs(15:0), rb(23:0)←rb(23:0)+2	2	–	–	–	–	–	–	*6	○
	[%rb]-, %rs	0	0	1	0	0	1	rs	1	1	1	0	rb	W[rb]←rs(15:0), rb(23:0)←rb(23:0)-2	2	–	–	–	–	–	–	*6	○
	-[%rb], %rs	0	0	1	0	0	1	rs	1	0	1	0	rb	rb(23:0)←rb(23:0)-2, W[rb]←rs(15:0)	2	–	–	–	–	–	–	*6	○

Remarks

*1) With one EXT: base address = rb+imm13, With two EXT: base address = rb+imm24
*2) With one EXT: data = sign16
*3) With one EXT: data = imm20, With two EXT: data = imm24
*4) With one EXT: base address = imm20, With two EXT: base address = imm24
*5) With one EXT: base address = sp+imm20, With two EXT: base address = sp+imm24
*6) With one EXT: base address = rb, address increment/decrement rb/sp ← rb/sp±imm13, With two EXT: base address = rb, address increment/decrement rb/sp ← rb/sp±imm24
*7) With no EXT: 1 cycle, With EXT: 2 cycles

Data Transfer Instructions (2)

S1C17 Core Instruction Set

Mnemonic: Opcode	Mnemonic: Operand	Code (MSB … LSB)	Function	Cycle	Flags: IL	IE	C	V	Z	N	EXT	D
ld	[%sp+imm7], %rs	1 1 1 1 1 0 rs imm7	W[sp+imm7]←rs(15:0)	2	–	–	–	–	–	–	*5	○
	[imm7], %rs	1 1 0 1 1 0 rs imm7	W[imm7]←rs(15:0)	1	–	–	–	–	–	–	*4	○
ld.a	%rd, %rs	0 0 1 0 1 0 rd 0 0 1 1 rs	rd(23:0)←rs(23:0)	1	–	–	–	–	–	–	–	○
	%rd, imm7	1 0 0 1 1 1 rd imm7	rd(6:0)←imm7(6:0), rd(23:7)←0	1	–	–	–	–	–	–	*3	○
	%rd, [%rb]	0 0 1 0 0 0 rd 0 0 1 1 rb	rd(23:0)←A[rb](23:0), ignored←A[rb](31:24)	1, 2*8	–	–	–	–	–	–	*1	○
	%rd, [%rb]+	0 0 1 0 0 0 rd 0 1 1 1 rb	rd(23:0)←A[rb](23:0), ignored←A[rb](31:24), rb(23:0)←rb(23:0)+4	2	–	–	–	–	–	–	*6	○
	%rd, [%rb]-	0 0 1 0 0 0 rd 1 1 1 1 rb	rd(23:0)←A[rb](23:0), ignored←A[rb](31:24), rb(23:0)←rb(23:0)-4	2	–	–	–	–	–	–	*6	○
	%rd, -[%rb]	0 0 1 0 0 0 rd 1 0 1 1 rb	rb(23:0)←rb(23:0)-4, rd(23:0)←A[rb](23:0), ignored←A[rb](31:24)	2	–	–	–	–	–	–	*6	○
	%rd, [%sp+imm7]	1 1 1 0 1 1 rd imm7	rd(23:0)←A[sp+imm7](23:0), ignored←A[sp+imm7](31:24)	2	–	–	–	–	–	–	*5	○
	%rd, [imm7]	1 1 0 0 1 1 rd imm7	rd(23:0)←A[imm7](23:0), ignored←A[imm7](31:24)	1	–	–	–	–	–	–	*4	○
	[%rb], %rs	0 0 1 0 0 1 rs 0 0 1 1 rb	A[rb](23:0)←rs(23:0), A[rb](31:24)←0	1, 2*8	–	–	–	–	–	–	*1	○
	[%rb]+, %rs	0 0 1 0 0 1 rs 0 1 1 1 rb	A[rb](23:0)←rs(23:0), A[rb](31:24)←0, rb(23:0)←rb(23:0)+4	2	–	–	–	–	–	–	*6	○
	[%rb]-, %rs	0 0 1 0 0 1 rs 1 1 1 1 rb	A[rb](23:0)←rs(23:0), A[rb](31:24)←0, rb(23:0)←rb(23:0)-4	2	–	–	–	–	–	–	*6	○
	-[%rb], %rs	0 0 1 0 0 1 rs 1 0 1 1 rb	rb(23:0)←rb(23:0)-4, A[rb](23:0)←rs(23:0), A[rb](31:24)←0	2	–	–	–	–	–	–	*6	○
	[%sp+imm7], %rs	1 1 1 1 1 1 rs imm7	A[sp+imm7](23:0)←rs(23:0), A[sp+imm7](31:24)←0	2	–	–	–	–	–	–	*5	○
	[imm7], %rs	1 1 0 1 1 1 rs imm7	A[imm7](23:0)←rs(23:0), A[imm7](31:24)←0	1	–	–	–	–	–	–	*4	○
	%rd, %sp	0 0 1 1 1 1 rd 0 0 1 0 0 0 0	rd(23:2)←sp(23:2), rd(1:0)←0	1	–	–	–	–	–	–	–	○
	%rd, %pc (*7)	0 0 1 1 1 1 rd 0 1 1 0 0 0 0	rd(23:0)←pc(23:0)+2	1	–	–	–	–	–	–	–	○
	%rd, [%sp]	0 0 1 1 1 1 rd 0 0 1 1 0 0 0	rd(23:0)←A[sp](23:0), ignored←A[sp](31:24)	1, 2*8	–	–	–	–	–	–	*1	○
	%rd, [%sp]+	0 0 1 1 1 1 rd 0 1 1 1 0 0 0	rd(23:0)←A[sp](23:0), ignored←A[sp](31:24), sp(23:0)←sp(23:0)+4	2	–	–	–	–	–	–	*6	○
	%rd, [%sp]-	0 0 1 1 1 1 rd 1 1 1 1 0 0 0	rd(23:0)←A[sp](23:0), ignored←A[sp](31:24), sp(23:0)←sp(23:0)-4	2	–	–	–	–	–	–	*6	○
	%rd, -[%sp]	0 0 1 1 1 1 rd 1 0 1 1 0 0 0	sp(23:0)←sp(23:0)-4, rd(23:0)←A[sp](23:0), ignored←A[sp](31:24)	2	–	–	–	–	–	–	*6	○
	[%sp], %rs	0 0 1 1 1 1 rs 0 0 1 1 1 0 0	A[sp](23:0)←rs(23:0), A[sp](31:24)←0	1, 2*8	–	–	–	–	–	–	*1	○
	[%sp]+, %rs	0 0 1 1 1 1 rs 0 1 1 1 1 0 0	A[sp](23:0)←rs(23:0), A[sp](31:24)←0, sp(23:0)←sp(23:0)+4	2	–	–	–	–	–	–	*6	○
	[%sp]-, %rs	0 0 1 1 1 1 rs 1 1 1 1 1 0 0	A[sp](23:0)←rs(23:0), A[sp](31:24)←0, sp(23:0)←sp(23:0)-4	2	–	–	–	–	–	–	*6	○
	-[%sp], %rs	0 0 1 1 1 1 rs 1 0 1 1 1 0 0	sp(23:0)←sp(23:0)-4, A[sp](23:0)←rs(23:0), A[sp](31:24)←0	2	–	–	–	–	–	–	*6	○
	%sp, %rs	0 0 1 1 1 1 rs 1 0 1 0 0 0 0	sp(23:2)←rs(23:2), sp(1:0)←0	1	–	–	–	–	–	–	–	○
	%sp, imm7	1 0 1 1 1 1 0 0 0 imm7	sp(6:2)←imm7(6:2), sp(23:7)←0, sp(1:0)←0	1	–	–	–	–	–	–	*3	○

Remarks

*1) With one EXT: base address = rb+imm13, With two EXT: base address = rb+imm24
*2) With one EXT: data = sign16
*3) With one EXT: data = imm20, With two EXT: data = imm24
*4) With one EXT: base address = imm20, With two EXT: base address = imm24
*5) With one EXT: base address = sp+imm20, With two EXT: base address = sp+imm24
*6) With one EXT: base address = rb, address increment/decrement rb/sp ← rb/sp±imm13, With two EXT: base address = rb, address increment/decrement rb/sp ← rb/sp±imm24
7) The "ld.a %rd, %pc" instruction should be used as a delayed slot instruction for the jr.d, jpr.d, or jpa.d delayed branch instruction.
*8) With no EXT: 1 cycle, With EXT: 2 cycles

Integer Arithmetic Operation Instructions (1)

S1C17 Core Instruction Set

Mnemonic		Code	Function	Cycle	Flags						EXT	D
Opcode	Operand	MSB … LSB			IL	IE	C	V	Z	N		
add	%rd, %rs	0 0 1 1 1 0 rd 1 0 0 0 rs	rd(15:0)←rd(15:0)+rs(15:0), rd(23:16)←0	1	–	–	↔	↔	↔	↔	*1	○
add/c	%rd, %rs	0 0 1 1 1 0 rd 0 0 0 0 rs	rd(15:0)←rd(15:0)+rs(15:0), rd(23:16)←0 if C = 1 (nop if C = 0)	1	–	–	–	↔	↔	↔	*1	○
add/nc	%rd, %rs	0 0 1 1 1 0 rd 0 1 0 0 rs	rd(15:0)←rd(15:0)+rs(15:0), rd(23:16)←0 if C = 0 (nop if C = 1)	1	–	–	–	↔	↔	↔	*1	○
add	%rd, imm7	1 0 0 0 0 0 rd imm7	rd(15:0)←rd(15:0)+imm7(zero extended), rd(23:16)←0	1	–	–	↔	↔	↔	↔	*3	○
add.a	%rd, %rs	0 0 1 1 0 0 rd 1 0 0 0 rs	rd(23:0)←rd(23:0)+rs(23:0)	1	–	–	–	–	–	–	*2	○
add.a/c	%rd, %rs	0 0 1 1 0 0 rd 0 0 0 0 rs	rd(23:0)←rd(23:0)+rs(23:0) if C = 1 (nop if C = 0)	1	–	–	–	–	–	–	*2	○
add.a/nc	%rd, %rs	0 0 1 1 0 0 rd 0 1 0 0 rs	rd(23:0)←rd(23:0)+rs(23:0) if C = 0 (nop if C = 1)	1	–	–	–	–	–	–	*2	○
add.a	%sp, %rs	0 0 1 1 0 0 0 0 0 0 0 0 1 rs	sp(23:0)←sp(23:0)+rs(23:0)	1	–	–	–	–	–	–	*2	○
	%rd, imm7	0 1 1 0 0 0 rd imm7	rd(23:0)←rd(23:0)+imm7(zero extended)	1	–	–	–	–	–	–	*4	○
	%sp, imm7	0 1 1 0 0 1 0 0 0 imm7	sp(23:0)←sp(23:0)+imm7(zero extended)	1	–	–	–	–	–	–	*4	○
adc	%rd, %rs	0 0 1 1 1 0 rd 1 0 0 1 rs	rd(15:0)←rd(15:0)+rs(15:0)+C, rd(23:16)←0	1	–	–	↔	↔	↔	↔	*1	○
adc/c	%rd, %rs	0 0 1 1 1 0 rd 0 0 0 1 rs	rd(15:0)←rd(15:0)+rs(15:0)+C, rd(23:16)←0 if C = 1 (nop if C = 0)	1	–	–	–	↔	↔	↔	*1	○
adc/nc	%rd, %rs	0 0 1 1 1 0 rd 0 1 0 1 rs	rd(15:0)←rd(15:0)+rs(15:0)+C, rd(23:16)←0 if C = 0 (nop if C = 1)	1	–	–	–	↔	↔	↔	*1	○
adc	%rd, imm7	1 0 0 0 0 1 rd imm7	rd(15:0)←rd(15:0)+imm7(zero extended)+C, rd(23:16)←0	1	–	–	↔	↔	↔	↔	*3	○
sub	%rd, %rs	0 0 1 1 1 0 rd 1 0 1 0 rs	rd(15:0)←rd(15:0)-rs(15:0), rd(23:16)←0	1	–	–	↔	↔	↔	↔	*1	○
sub/c	%rd, %rs	0 0 1 1 1 0 rd 0 0 1 0 rs	rd(15:0)←rd(15:0)-rs(15:0), rd(23:16)←0 if C = 1 (nop if C = 0)	1	–	–	–	↔	↔	↔	*1	○
sub/nc	%rd, %rs	0 0 1 1 1 0 rd 0 1 1 0 rs	rd(15:0)←rd(15:0)-rs(15:0), rd(23:16)←0 if C = 0 (nop if C = 1)	1	–	–	–	↔	↔	↔	*1	○
sub	%rd, imm7	1 0 0 0 1 0 rd imm7	rd(15:0)←rd(15:0)-imm7(zero extended), rd(23:16)←0	1	–	–	↔	↔	↔	↔	*3	○
sub.a	%rd, %rs	0 0 1 1 0 0 rd 1 0 1 0 rs	rd(23:0)←rd(23:0)-rs(23:0)	1	–	–	–	–	–	–	*2	○
sub.a/c	%rd, %rs	0 0 1 1 0 0 rd 0 0 1 0 rs	rd(23:0)←rd(23:0)-rs(23:0) if C = 1 (nop if C = 0)	1	–	–	–	–	–	–	*2	○
sub.a/nc	%rd, %rs	0 0 1 1 0 0 rd 0 1 1 0 rs	rd(23:0)←rd(23:0)-rs(23:0) if C = 0 (nop if C = 1)	1	–	–	–	–	–	–	*2	○
sub.a	%sp, %rs	0 0 1 1 0 0 0 0 0 0 0 1 1 rs	sp(23:0)←sp(23:0)-rs(23:0)	1	–	–	–	–	–	–	*2	○
	%rd, imm7	0 1 1 0 1 0 rd imm7	rd(23:0)←rd(23:0)-imm7(zero extended)	1	–	–	–	–	–	–	*4	○
	%sp, imm7	0 1 1 0 1 1 0 0 0 imm7	sp(23:0)←sp(23:0)-imm7(zero extended)	1	–	–	–	–	–	–	*4	○
sbc	%rd, %rs	0 0 1 1 1 0 rd 1 0 1 1 rs	rd(15:0)←rd(15:0)-rs(15:0)-C, rd(23:16)←0	1	–	–	↔	↔	↔	↔	*1	○
sbc/c	%rd, %rs	0 0 1 1 1 0 rd 0 0 1 1 rs	rd(15:0)←rd(15:0)-rs(15:0)-C, rd(23:16)←0 if C = 1 (nop if C = 0)	1	–	–	–	↔	↔	↔	*1	○
sbc/nc	%rd, %rs	0 0 1 1 1 0 rd 0 1 1 1 rs	rd(15:0)←rd(15:0)-rs(15:0)-C, rd(23:16)←0 if C = 0 (nop if C = 1)	1	–	–	–	↔	↔	↔	*1	○
sbc	%rd, imm7	1 0 0 0 1 1 rd imm7	rd(15:0)←rd(15:0)-imm7(zero extended)-C, rd(23:16)←0	1	–	–	↔	↔	↔	↔	*3	○
cmp	%rd, %rs	0 0 1 1 1 1 rd 1 0 0 0 rs	rd(15:0)-rs(15:0)	1	–	–	↔	↔	↔	↔	*1	○
cmp/c	%rd, %rs	0 0 1 1 1 1 rd 0 0 0 0 rs	rd(15:0)-rs(15:0) if C = 1 (nop if C = 0)	1	–	–	–	↔	↔	↔	*1	○
cmp/nc	%rd, %rs	0 0 1 1 1 1 rd 0 1 0 0 rs	rd(15:0)-rs(15:0) if C = 0 (nop if C = 1)	1	–	–	–	↔	↔	↔	*1	○
cmp	%rd, sign7	1 0 0 1 0 0 rd sign7	rd(15:0)-sign7(sign extended)	1	–	–	↔	↔	↔	↔	*3	○

Remarks

*1) With one EXT: rd ← rs <op> imm13, With two EXT: rd ← rs <op> imm16
*2) With one EXT: rd ← rs <op> imm13, With two EXT: rd ← rs <op> imm24
*3) With one EXT: data = imm16/sign16
*4) With one EXT: data = imm20, With two EXT: data = imm24

Integer Arithmetic Operation Instructions (2)

S1C17 Core Instruction Set

Mnemonic		Code												Function	Cycle	Flags						EXT	D
Opcode	Operand	MSB											LSB			IL	IE	C	V	Z	N		
cmp.a	%rd, %rs	0	0	1	1	0	1	rd	1	0	0	0	rs	rd(23:0)-rs(23:0)	1	–	–	↔	–	↔	–	*2	○
cmp.a/c	%rd, %rs	0	0	1	1	0	1	rd	0	0	0	0	rs	rd(23:0)-rs(23:0) if C = 1 (nop if C = 0)	1	–	–	–	–	↔	–	*2	○
cmp.a/nc	%rd, %rs	0	0	1	1	0	1	rd	0	1	0	0	rs	rd(23:0)-rs(23:0) if C = 0 (nop if C = 1)	1	–	–	–	–	↔	–	*2	○
cmp.a	%rd, imm7	0	1	1	1	0	0	rd	imm7					rd(23:0)-imm7(zero extended)	1	–	–	↔	–	↔	–	*4	○
cmc	%rd, %rs	0	0	1	1	1	1	rd	1	0	0	1	rs	rd(15:0)-rs(15:0)-C	1	–	–	↔	↔	↔	↔	*1	○
cmc/c	%rd, %rs	0	0	1	1	1	1	rd	0	0	0	1	rs	rd(15:0)-rs(15:0)-C if C = 1 (nop if C = 0)	1	–	–	–	↔	↔	↔	*1	○
cmc/nc	%rd, %rs	0	0	1	1	1	1	rd	0	1	0	1	rs	rd(15:0)-rs(15:0)-C if C = 0 (nop if C = 1)	1	–	–	–	↔	↔	↔	*1	○
cmc	%rd, sign7	1	0	0	1	0	1	rd	sign7					rd(15:0)-sign7(sign extended)-C	1	–	–	↔	↔	↔	↔	*3	○

Remarks

*1) With one EXT: rd ← rs <op> imm13, With two EXT: rd ← rs <op> imm16
*2) With one EXT: rd ← rs <op> imm13, With two EXT: rd ← rs <op> imm24
*3) With one EXT: data = imm16/sign16
*4) With one EXT: data = imm20, With two EXT: data = imm24

Logic Operation Instructions

S1C17 Core Instruction Set

Mnemonic		Code												Function	Cycle	Flags						EXT	D
Opcode	Operand	MSB											LSB			IL	IE	C	V	Z	N		
and	%rd, %rs	0	0	1	0	1	1	rd	1	0	0	0	rs	rd(15:0)←rd(15:0)&rs(15:0), rd(23:16)←0	1	–	–	–	0	↔	↔	*1	○
and/c	%rd, %rs	0	0	1	0	1	1	rd	0	0	0	0	rs	rd(15:0)←rd(15:0)&rs(15:0), rd(23:16)←0 if C = 1 (nop if C = 0)	1	–	–	–	0	↔	↔	*1	○
and/nc	%rd, %rs	0	0	1	0	1	1	rd	0	1	0	0	rs	rd(15:0)←rd(15:0)&rs(15:0), rd(23:16)←0 if C = 0 (nop if C = 1)	1	–	–	–	0	↔	↔	*1	○
and	%rd, sign7	1	0	1	0	0	0	rd	sign7					rd(15:0)←rd(15:0)&sign7(sign extended), rd(23:16)←0	1	–	–	–	0	↔	↔	*2	○
or	%rd, %rs	0	0	1	0	1	1	rd	1	0	0	1	rs	rd(15:0)←rd(15:0) | rs(15:0), rd(23:16)←0	1	–	–	–	0	↔	↔	*1	○
or/c	%rd, %rs	0	0	1	0	1	1	rd	0	0	0	1	rs	rd(15:0)←rd(15:0) | rs(15:0), rd(23:16)←0 if C = 1 (nop if C = 0)	1	–	–	–	0	↔	↔	*1	○
or/nc	%rd, %rs	0	0	1	0	1	1	rd	0	1	0	1	rs	rd(15:0)←rd(15:0) | rs(15:0), rd(23:16)←0 if C = 0 (nop if C = 1)	1	–	–	–	0	↔	↔	*1	○
or	%rd, sign7	1	0	1	0	0	1	rd	sign7					rd(15:0)←rd(15:0) | sign7(sign extended), rd(23:16)←0	1	–	–	–	0	↔	↔	*2	○
xor	%rd, %rs	0	0	1	0	1	1	rd	1	0	1	0	rs	rd(15:0)←rd(15:0)^rs(15:0), rd(23:16)←0	1	–	–	–	0	↔	↔	*1	○
xor/c	%rd, %rs	0	0	1	0	1	1	rd	0	0	1	0	rs	rd(15:0)←rd(15:0)^rs(15:0), rd(23:16)←0 if C = 1 (nop if C = 0)	1	–	–	–	0	↔	↔	*1	○
xor/nc	%rd, %rs	0	0	1	0	1	1	rd	0	1	1	0	rs	rd(15:0)←rd(15:0)^rs(15:0), rd(23:16)←0 if C = 0 (nop if C = 1)	1	–	–	–	0	↔	↔	*1	○
xor	%rd, sign7	1	0	1	0	1	0	rd	sign7					rd(15:0)←rd(15:0)^sign7(sign extended), rd(23:16)←0	1	–	–	–	0	↔	↔	*2	○
not	%rd, %rs	0	0	1	0	1	1	rd	1	0	1	1	rs	rd(15:0)←!rs(15:0), rd(23:16)←0	1	–	–	–	0	↔	↔	*3	○
not/c	%rd, %rs	0	0	1	0	1	1	rd	0	0	1	1	rs	rd(15:0)←!rs(15:0), rd(23:16)←0 if C = 1 (nop if C = 0)	1	–	–	–	0	↔	↔	*3	○
not/nc	%rd, %rs	0	0	1	0	1	1	rd	0	1	1	1	rs	rd(15:0)←!rs(15:0), rd(23:16)←0 if C = 0 (nop if C = 1)	1	–	–	–	0	↔	↔	*3	○
not	%rd, sign7	1	0	1	0	1	1	rd	sign7					rd(15:0)←!sign7(sign extended), rd(23:16)←0	1	–	–	–	0	↔	↔	*2	○

Remarks

*1) With one EXT: rd ← rs <op> imm13, With two EXT: rd ← rs <op> imm16
*2) With one EXT: data = sign16
*3) With one EXT: rd ← !imm13, With two EXT: rd ← !imm16

Branch Instructions

S1C17 Core Instruction Set

Mnemonic Opcode	Mnemonic Operand	Code (MSB … LSB)	Function	Cycle *6	IL	IE	C	V	Z	N	EXT	D
jpr / jpr.d	sign10	0 0 0 1 0 d sign10	pc←pc+2+sign11; sign11={sign10,0} (*3)		–	–	–	–	–	–	*4	–
	%rb	0 0 0 0 0 0 0 1 d 1 0 0 0 rb	pc←pc+2+rb (*3)	3	–	–	–	–	–	–	–	–
jpa / jpa.d	imm7	0 0 0 0 0 0 1 1 d imm7	pc←imm7 (*3)	2(.d)	–	–	–	–	–	–	*2	–
	%rb	0 0 0 0 0 0 0 1 d 1 0 0 1 rb	pc←rb (*3)	3	–	–	–	–	–	–	–	–
jrgt / jrgt.d	sign7	0 0 0 0 0 1 1 0 d sign7	pc←pc+2+sign8 if !Z&!(N^V) is true; sign8={sign7,0} (*3)	2(.d)	–	–	–	–	–	–	*1	–
jrge / jrge.d	sign7	0 0 0 0 0 1 1 1 d sign7	pc←pc+2+sign8 if !(N^V) is true; sign8={sign7,0} (*3)		–	–	–	–	–	–	*1	–
jrlt / jrlt.d	sign7	0 0 0 0 1 0 0 0 d sign7	pc←pc+2+sign8 if N^V is true; sign8={sign7,0} (*3)	2	–	–	–	–	–	–	*1	–
jrle / jrle.d	sign7	0 0 0 0 1 0 0 1 d sign7	pc←pc+2+sign8 if Z | (N^V) is true; sign8={sign7,0} (*3)	(false)	–	–	–	–	–	–	*1	–
jrugt / jrugt.d	sign7	0 0 0 0 1 0 1 0 d sign7	pc←pc+2+sign8 if !Z&!C is true; sign8={sign7,0} (*3)	or	–	–	–	–	–	–	*1	–
jruge / jruge.d	sign7	0 0 0 0 1 0 1 1 d sign7	pc←pc+2+sign8 if !C is true; sign8={sign7,0} (*3)	3	–	–	–	–	–	–	*1	–
jrult / jrult.d	sign7	0 0 0 0 1 1 0 0 d sign7	pc←pc+2+sign8 if C is true; sign8={sign7,0} (*3)	(true)	–	–	–	–	–	–	*1	–
jrule / jrule.d	sign7	0 0 0 0 1 1 0 1 d sign7	pc←pc+2+sign8 if Z | C is true; sign8={sign7,0} (*3)	*5	–	–	–	–	–	–	*1	–
jreq / jreq.d	sign7	0 0 0 0 1 1 1 0 d sign7	pc←pc+2+sign8 if Z is true; sign8={sign7,0} (*3)		–	–	–	–	–	–	*1	–
jrne / jrne.d	sign7	0 0 0 0 1 1 1 1 d sign7	pc←pc+2+sign8 if !Z is true; sign8={sign7,0} (*3)	2(.d)	–	–	–	–	–	–	*1	–
call / call.d	sign10	0 0 0 1 1 d sign10	sp←sp-4, A[sp]←pc+2(d=0)/4(d=1), pc←pc+2+sign11; sign11={sign10,0} (*3)		–	–	–	–	–	–	*4	–
	%rb	0 0 0 0 0 0 0 1 d 0 0 0 0 rb	sp←sp-4, A[sp]←pc+2(d=0)/4(d=1), pc←pc+2+rb (*3)	4	–	–	–	–	–	–	–	–
calla / calla.d	imm7	0 0 0 0 0 1 0 1 d imm7	sp←sp-4, A[sp]←pc+2(d=0)/4(d=1), pc←imm7 (*3)	3(.d)	–	–	–	–	–	–	*2	–
	%rb	0 0 0 0 0 0 0 1 d 0 0 0 1 rb	sp←sp-4, A[sp]←pc+2(d=0)/4(d=1), pc←rb (*3)	4	–	–	–	–	–	–	–	–
ret / ret.d		0 0 0 0 0 0 0 1 d 0 1 0 0 0 0 0	pc←A[sp](23:0), sp←sp+4 (*3)	3(.d)	–	–	–	–	–	–	–	–
int	imm5	0 1 1 1 0 1 0 0 0 imm5 0 1	sp←sp-4, A[sp]←{psr, pc+2}, pc←vector(TTBR+imm5×4)	3, 2(.d)	–	0	–	–	–	–	–	–
intl	imm5, imm3	0 1 1 1 0 1 imm3 imm5 1 1	sp←sp-4, A[sp]←{psr, pc+2}, pc←vector(TTBR+imm5×4), psr(IL)←imm3	3	↔	0	–	–	–	–	–	–
reti / reti.d		0 0 0 0 0 0 0 1 d 0 1 0 1 0 0 0	{psr, pc}←A[sp], sp←sp+4	3	↔	↔	↔	↔	↔	↔	–	–
brk		0 0 0 0 0 0 0 1 0 1 1 0 0 0 0 0	A[DBRAM]←{psr, pc+2}, A[DBRAM+4]←r0, pc←0xfffc00	3, 2(.d)	–	0	–	–	–	–	–	–
retd		0 0 0 0 0 0 0 1 0 1 1 0 1 0 0 0	r0←A[DBRAM+4](23:0), {psr, pc}←A[DBRAM]	4	↔	↔	↔	↔	↔	↔	–	–

Remarks

*1) With one EXT: displacement = sign21 (= {imm13, sign7, 0}), With two EXT: displacement = sign24 (= {1st imm13(2:0), 2nd imm13, sign7, 0})
*2) With one EXT: absolute address= sign20 (= {imm13, imm7}), With two EXT: absolute address = sign24 (= {1st imm13(3:0), 2nd imm13, imm7})
*3) These instructions become a delayed branch instruction when the d bit in the code is set to 1 by suffixing ".d" to the opcode (jrgt.d, call.d, etc.).
*4) With one EXT: displacement = sign24 (= {imm13, sign10, 0})
*5) The conditional branch instructions other than delayed slot instructions (without ".d") are executed in two cycles when the program flow does not branch or three cycles when the program flow branches.
*6) The value with (.d) attached shows the number of cycles when 1-cycle delayed slot instruction follows. Otherwise, the same number of cycles as that shown without (.d) is required.

Immediate Extension Instruction

S1C17 Core Instruction Set

Mnemonic Opcode	Mnemonic Operand	Code (MSB … LSB)	Function	Cycle	IL	IE	C	V	Z	N	EXT	D
ext	imm13	0 1 0 imm13	Extends the immediate or operand of the following instruction.	1	–	–	–	–	–	–	*1	–

Remarks

*1) One or two ext instruction can be placed prior to the instructions that can be extended.

Shift and Swap Instructions

S1C17 Core Instruction Set

Mnemonic		Code	Function	Cycle	Flags						EXT	D
Opcode	Operand	MSB … LSB			IL	IE	C	V	Z	N		
sr	%rd, %rs	0 0 1 0 1 1 rd 1 1 0 0 rs	Logical shift to right; rd(15:0)←rd(15:0)>>rs(15:0), rd(23:16)←0, zero enters to MSB (∗1)	1	–	–	↔	–	↔	↔	–	○
	%rd, imm7	1 0 1 1 0 0 rd imm7	Logical shift to right; rd(15:0)←rd(15:0)>>imm7, rd(23:16)←0, zero enters to MSB (∗1)	1	–	–	↔	–	↔	↔	∗2	○
sa	%rd, %rs	0 0 1 0 1 1 rd 1 1 0 1 rs	Arithmetical shift to right; rd(15:0)←rd(15:0)>>rs(15:0), rd(23:16)←0, sign copied to MSB (∗1)	1	–	–	↔	–	↔	↔	–	○
	%rd, imm7	1 0 1 1 0 1 rd imm7	Arithmetical shift to right; rd(15:0)←rd(15:0)>>imm7, rd(23:16)←0, sign copied to MSB (∗1)	1	–	–	↔	–	↔	↔	∗2	○
sl	%rd, %rs	0 0 1 0 1 1 rd 1 1 1 0 rs	Logical shift to left; rd(15:0)←rd(15:0)<<rs(15:0), rd(23:16)←0, zero enters to LSB (∗1)	1	–	–	↔	–	↔	↔	–	○
	%rd, imm7	1 0 1 1 1 0 rd imm7	Logical shift to left; rd(15:0)←rd(15:0)<<imm7, rd(23:16)←0, zero enters to LSB (∗1)	1	–	–	↔	–	↔	↔	∗2	○
swap	%rd, %rs	0 0 1 0 1 1 rd 1 1 1 1 rs	rd(15:8)←rs(7:0), rd(7:0)←rs(15:8), rd(23:16)←0	1	–	–	–	–	–	–	–	○

Remarks

∗1) Number of bits to be shifted: Zero to three bits when rs/imm7 = 0–3, four bits when rs/imm7 = 4–7, eight bits when rs/imm7 ≥ 8

∗2) With one EXT: immediate = imm20, With two EXT: immediate = imm24

Conversion Instructions

S1C17 Core Instruction Set

Mnemonic		Code	Function	Cycle	Flags						EXT	D
Opcode	Operand	MSB … LSB			IL	IE	C	V	Z	N		
cv.ab	%rd, %rs	0 0 1 0 1 0 rd 0 1 1 1 rs	rd(23:8)←rs(7), rd(7:0)←rs(7:0)	1	–	–	–	–	–	–	–	○
cv.as	%rd, %rs	0 0 1 0 1 0 rd 1 0 1 1 rs	rd(23:16)←rs(15), rd(15:0)←rs(15:0)	1	–	–	–	–	–	–	–	○
cv.al	%rd, %rs	0 0 1 0 1 0 rd 1 1 1 1 rs	rd(23:16)←rs(7:0), rd(15:0)←rd(15:0)	1	–	–	–	–	–	–	–	○
cv.la	%rd, %rs	0 0 1 0 1 0 rd 0 1 1 0 rs	rd(23:8)←0, rd(7:0)←rs(23:16)	1	–	–	–	–	–	–	–	○
cv.ls	%rd, %rs	0 0 1 0 1 0 rd 1 0 1 0 rs	rd(23:16)←0, rd(15:0)←rs(15)	1	–	–	–	–	–	–	–	○

System Control Instructions

S1C17 Core Instruction Set

Mnemonic		Code	Function	Cycle	Flags						EXT	D
Opcode	Operand	MSB … LSB			IL	IE	C	V	Z	N		
nop		0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	No operation	1	–	–	–	–	–	–	–	○
halt		0 0 0 0 0 0 0 0 0 0 0 0 1 0 0 0	HALT mode	6	–	–	–	–	–	–	–	–
slp		0 0 0 0 0 0 0 0 0 0 0 1 0 0 0 0	SLEEP mode	6	–	–	–	–	–	–	–	–
ei		0 0 0 0 0 0 0 0 0 0 1 0 0 0 0 0	psr(IE)←1	1	–	1	–	–	–	–	–	○
di		0 0 0 0 0 0 0 0 0 1 0 0 0 0 0 0	psr(IE)←0	1	–	0	–	–	–	–	–	○

Coprocessor Interface Instructions

S1C17 Core Instruction Set

Mnemonic		Code	Function	Cycle	Flags						EXT	D
Opcode	Operand	MSB … LSB			IL	IE	C	V	Z	N		
ld.cw	%rd, %rs	0 0 1 1 0 1 rd 0 0 1 0 rs	co_dout0←rd, co_dout1←rs	1	–	–	–	–	–	–	–	○
	%rd, imm7	0 1 1 1 1 0 rd imm7	co_dout0←rd, co_dout1←imm7	1	–	–	–	–	–	–	*1	○
ld.ca	%rd, %rs	0 0 1 1 0 1 rd 0 0 1 1 rs	co_dout0←rd, co_dout1←rs, rd←co_din, psr(C, V, Z, N)←co_cvzn	1	–	–	↔	↔	↔	↔	–	○
	%rd, imm7	0 1 1 1 1 1 rd imm7	co_dout0←rd, co_dout1←imm7, rd←co_din, psr(C, V, Z, N)←co_cvzn	1	–	–	↔	↔	↔	↔	*1	○
ld.cf	%rd, %rs	0 0 1 1 0 1 rd 0 0 0 1 rs	co_dout0←rd, co_dout1←rs, psr(C, V, Z, N)←co_cvzn	1	–	–	↔	↔	↔	↔	–	○
	%rd, imm7	1 1 0 1 0 1 rd imm7	co_dout0←rd, co_dout1←imm7, psr(C, V, Z, N)←co_cvzn	1	–	–	↔	↔	↔	↔	*1	○

Remarks

*1) With one EXT: co_dout1 output = imm20, With two EXT: co_dout1 output = imm24

Revision History

Code No.	Page	Contents
410905900	All	New establishment
410905901	All	Made an overall revision.
410905902	6-10	Corrected the description in “Canceling HALT or SLEEP mode.”
	7-32, 7-33	Added “Caution” to the di and ei instruction pages.

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